

Hierarchical Petri nets

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Where innovation starts

Extensions of the classical Petri net

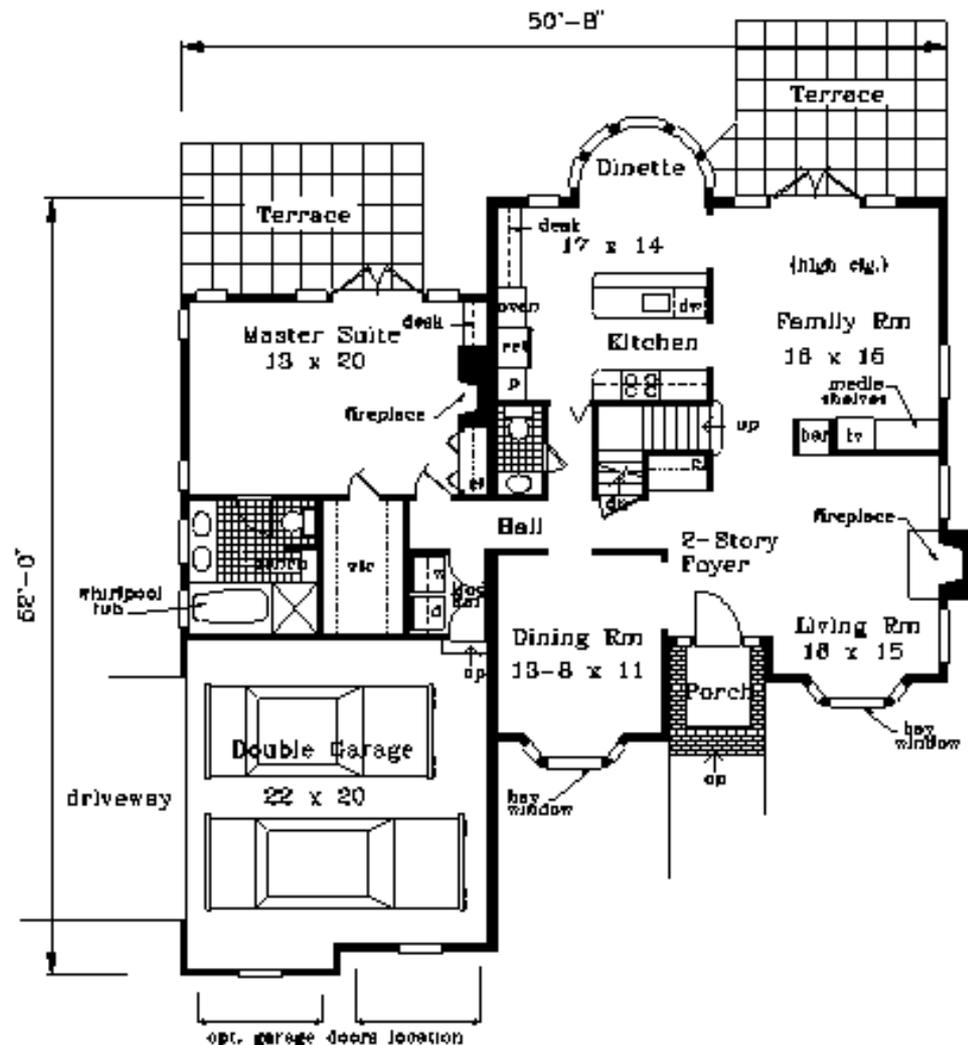
- **Color** (i.e., tokens have values, places have color sets, arcs have inscriptions, transitions have guards, etc.) was introduced in Chapter 5 and formalized in Chapter 6.
- **Time** (i.e., tokens have timestamps, color set may be timed, arcs have delays, etc.) was introduced in Chapter 5 and formalized in Chapter 6.
- **Hierarchy** is introduced and formalized in Chapter 7.
- A colored, timed, hierarchical Petri net is called a **high-level Petri net**.
- **Hierarchical CPNs** (=CPN + hierarchy) provide a concrete language for high-level Petri nets.

The need for hierarchy: Compare with construction drawings of a house

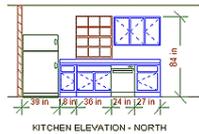
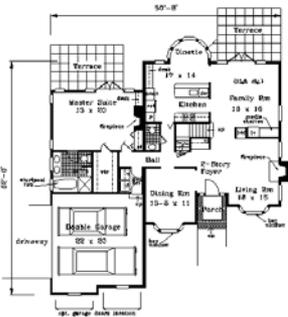


- **The top-level: The house as a whole**

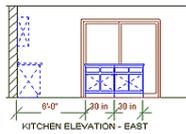
The second level: The first floor of the house.



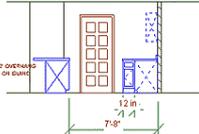
Top-down versus bottom-up



KITCHEN ELEVATION - NORTH



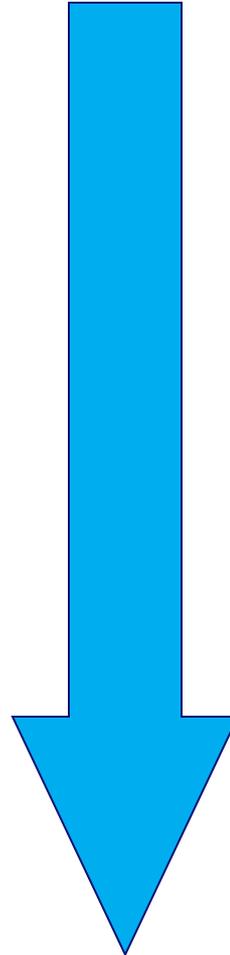
KITCHEN ELEVATION - EAST



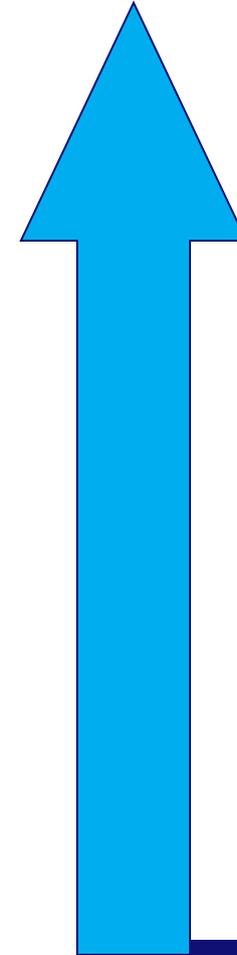
KITCHEN ELEVATION - SOUTH



KITCHEN ELEVATION - WEST



top-down



bottom-up

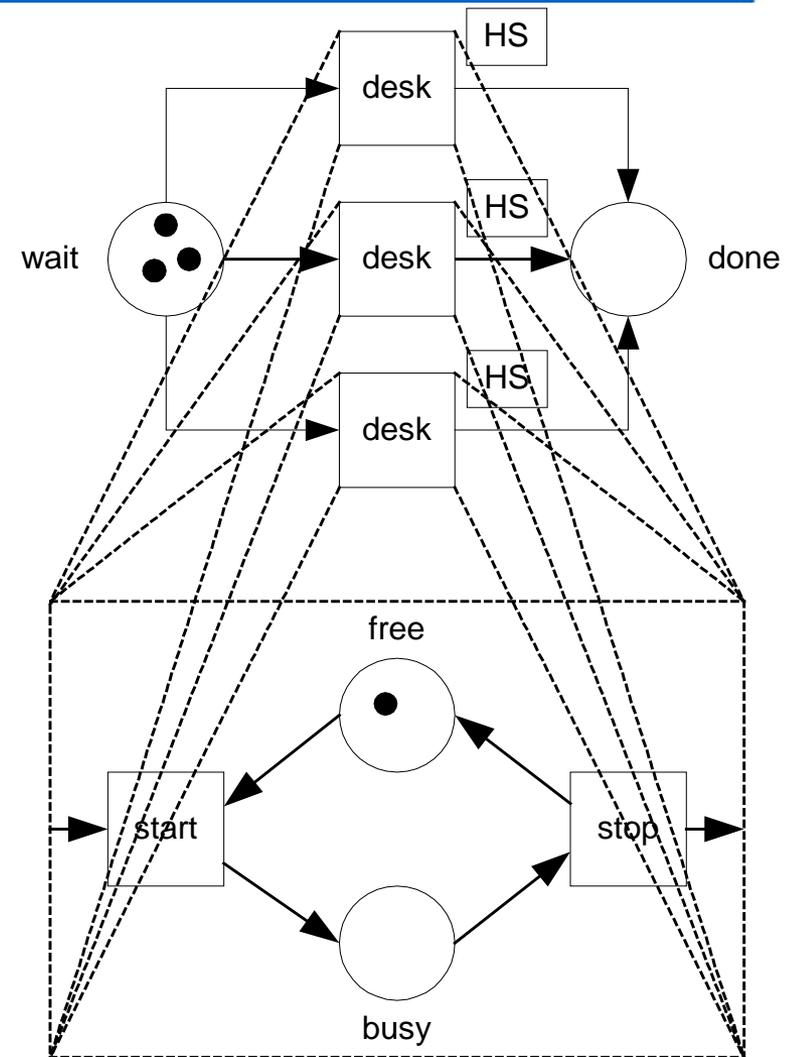
Recall

- Three good reasons for making a process model:
 - **gain insight**
for a better understanding of the system
 - **analysis**
validation and verification
 - **specification**
a blue print for construction
- Like the construction drawing of an architect!
- However, despite the addition of color and time, a hierarchy concept is still missing thus far!

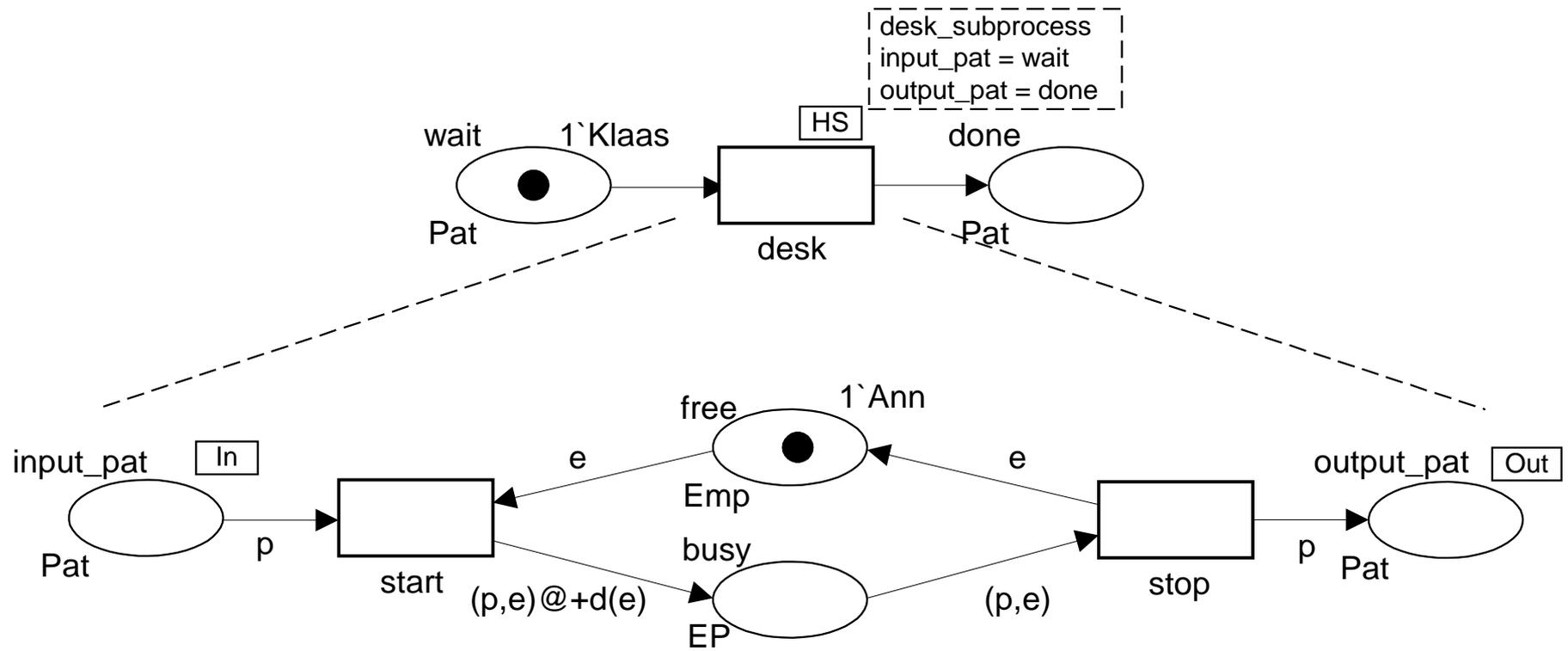
Basic idea

- Transitions correspond to subsystems/subprocesses
- Divide and conquer
- Reuse

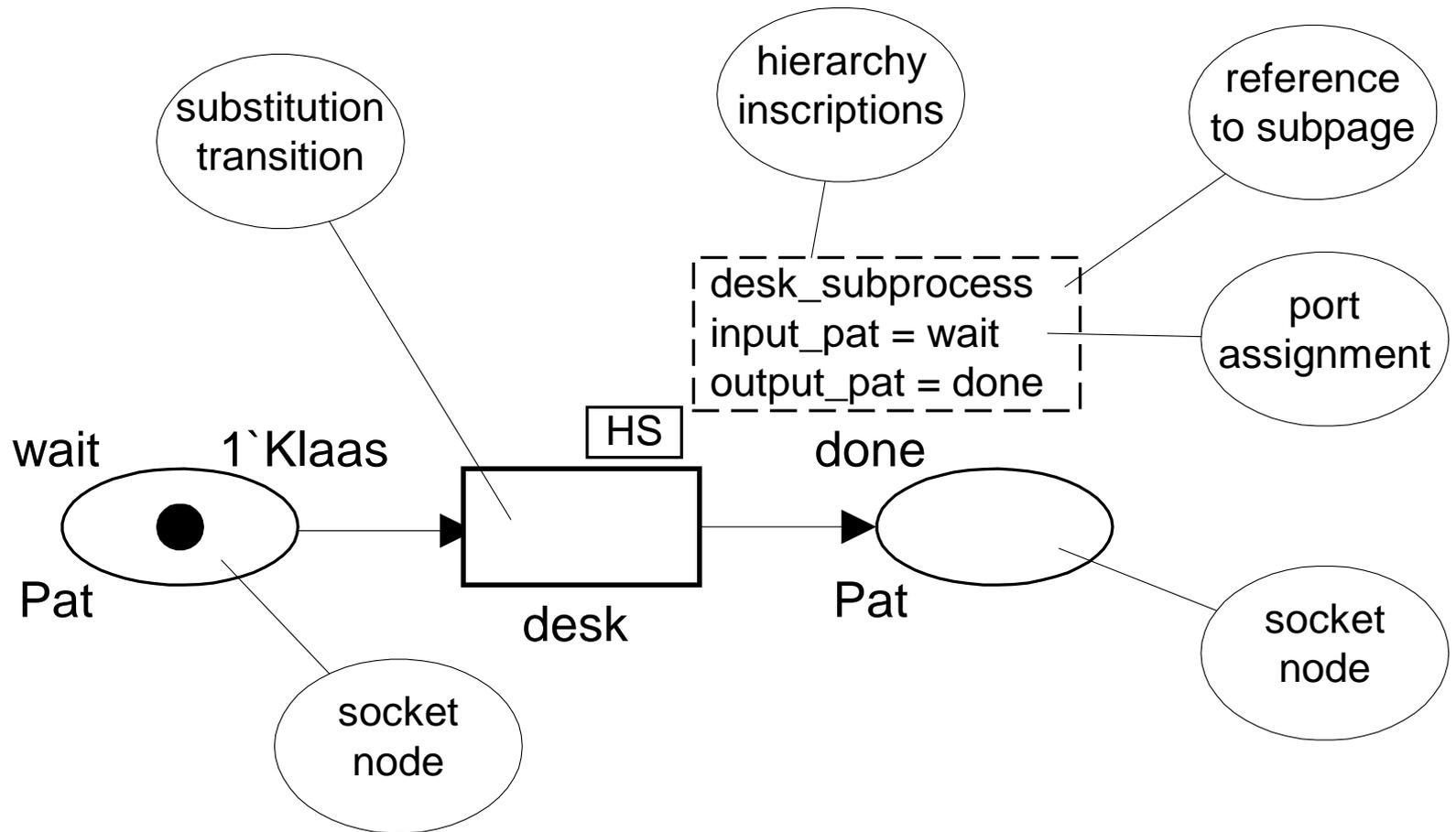
Let us formalize this in CPN...



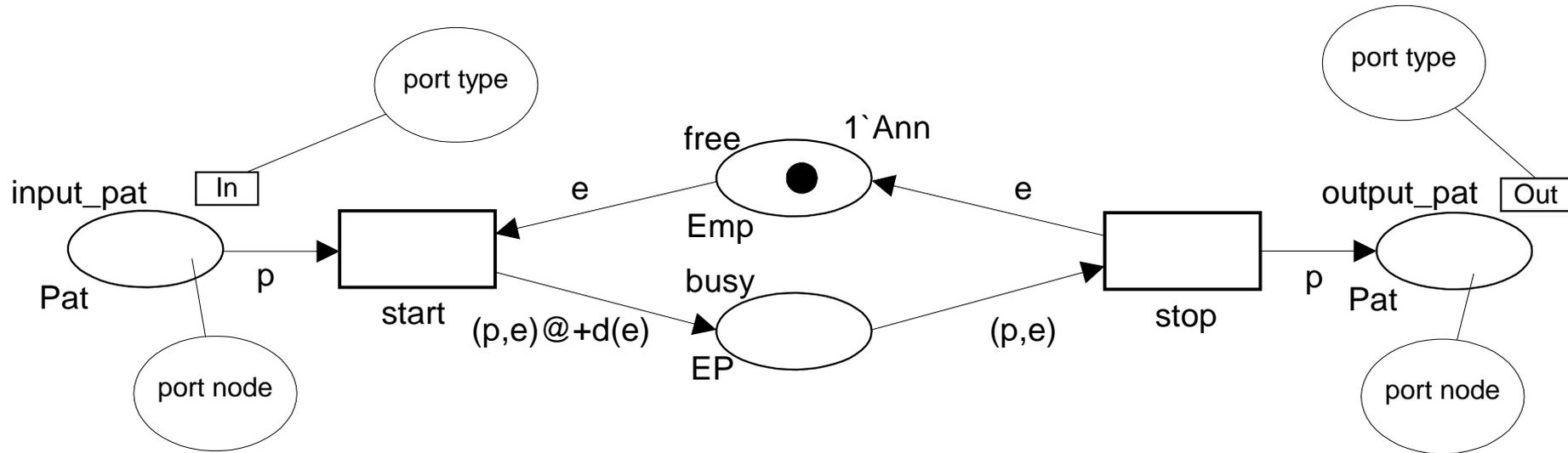
Hierarchical CPN



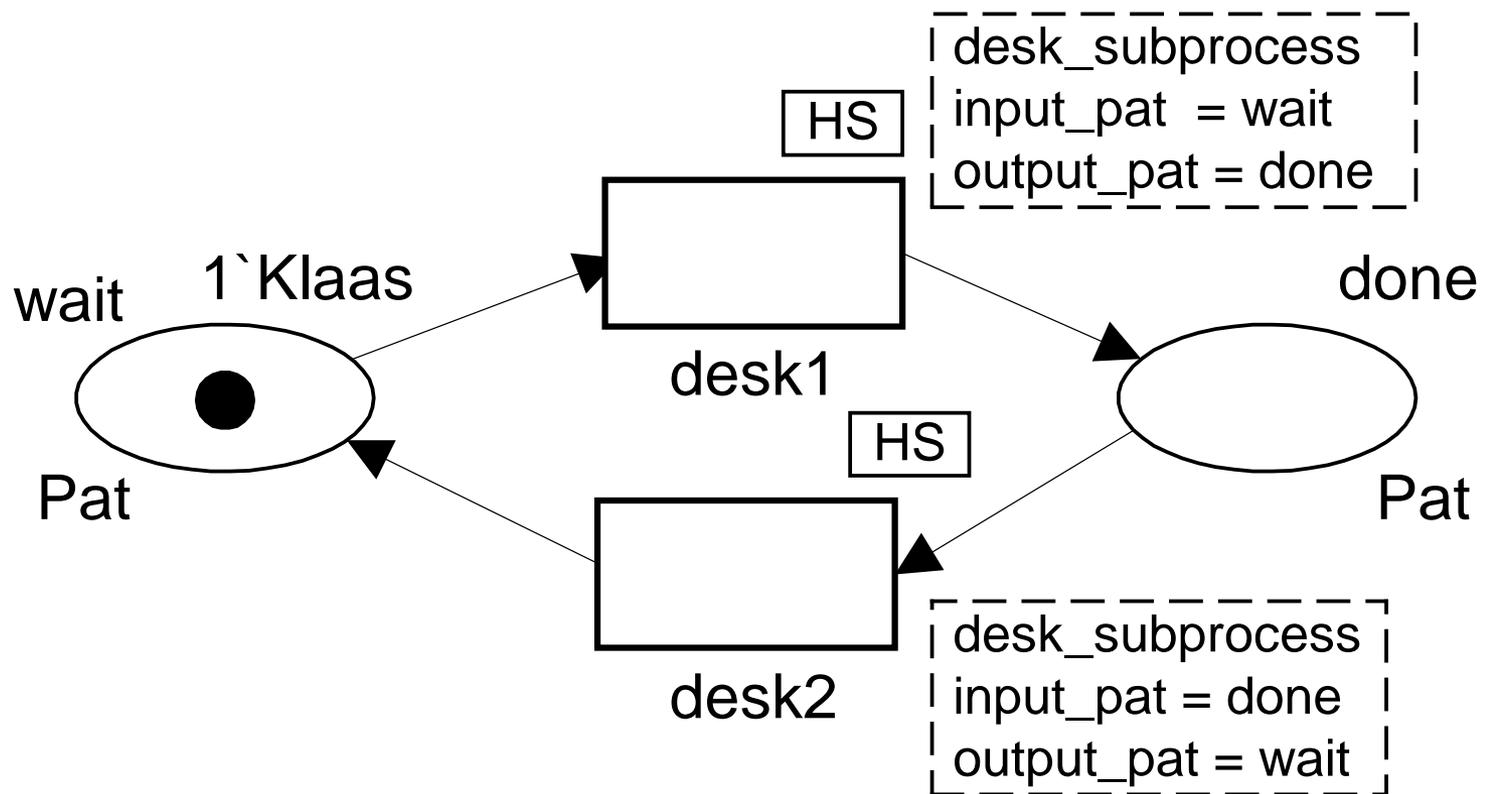
Superpage: *main*



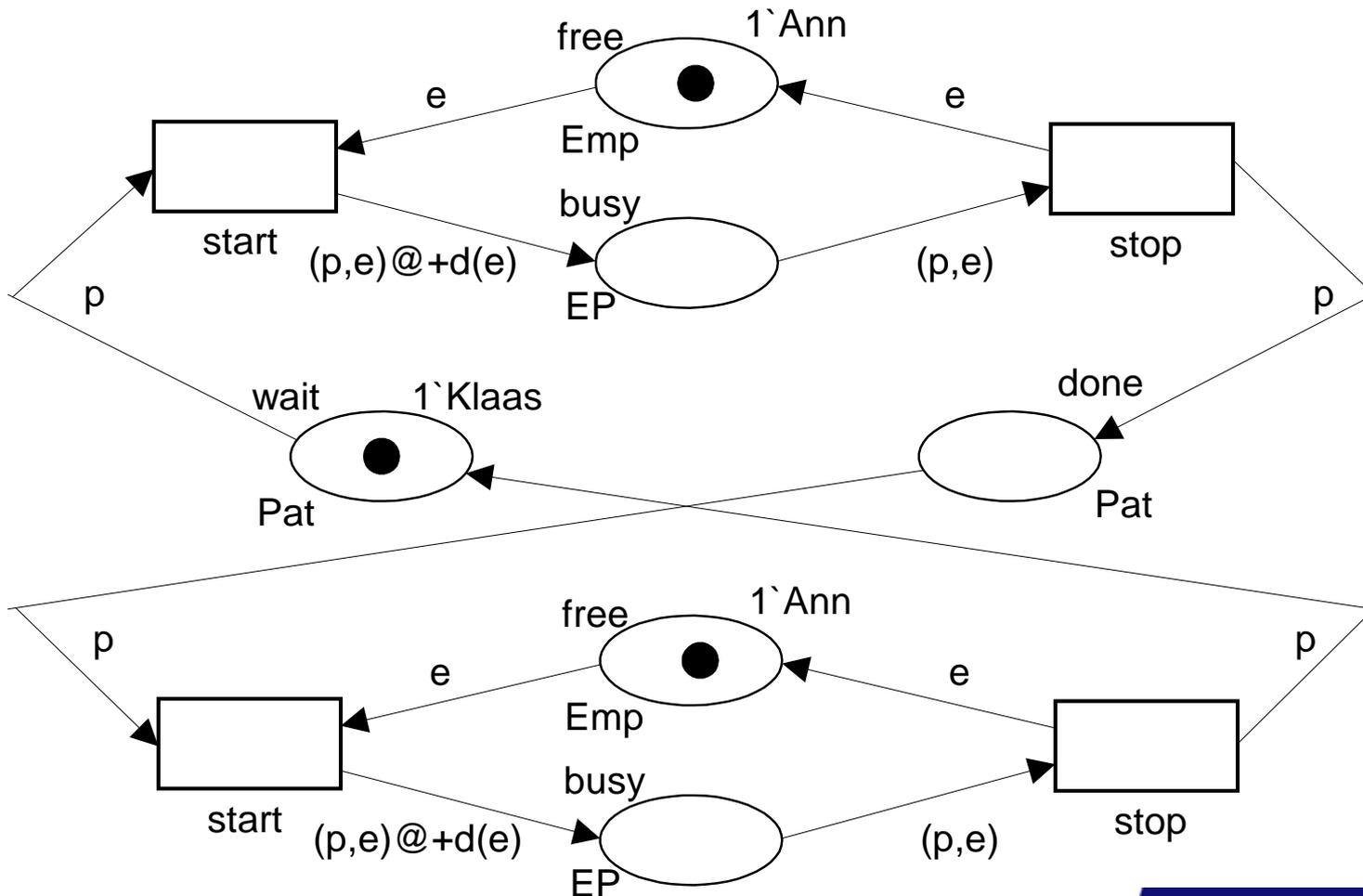
Subpage: *desk_subprocess*



One page can have multiple instances (i.e., reuse)



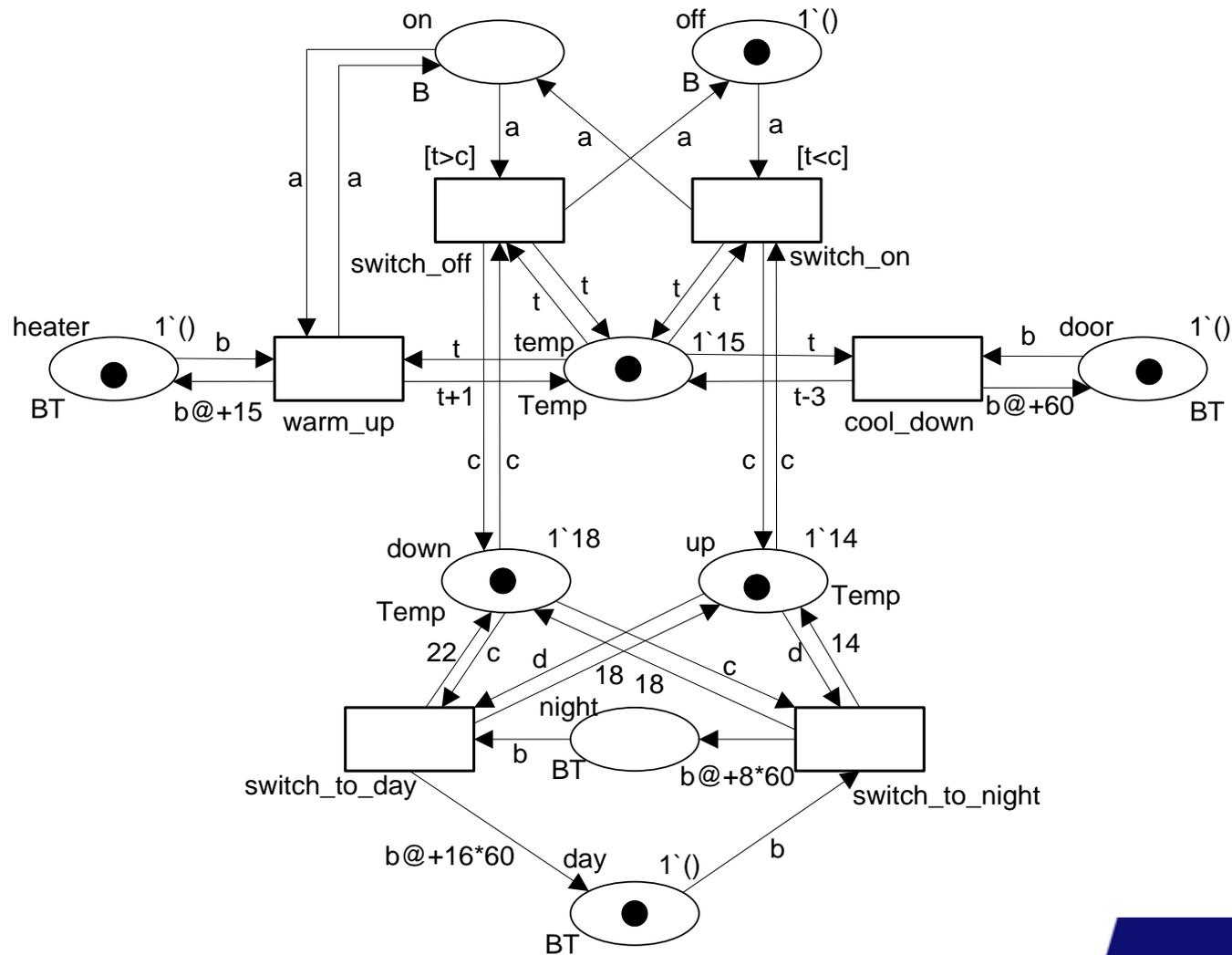
Semantics: Replace each substitution transition by a copy of the corresponding page



Thermostat system: Add hierarchy

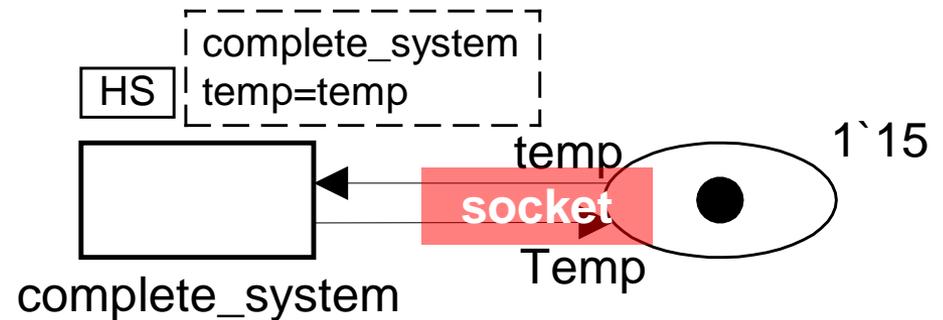
```

color Temp = string;
color B = unit;
color BT = B timed;
var t:Temp;
var a:B;
var b:BT;
var c,d:Temp;
    
```

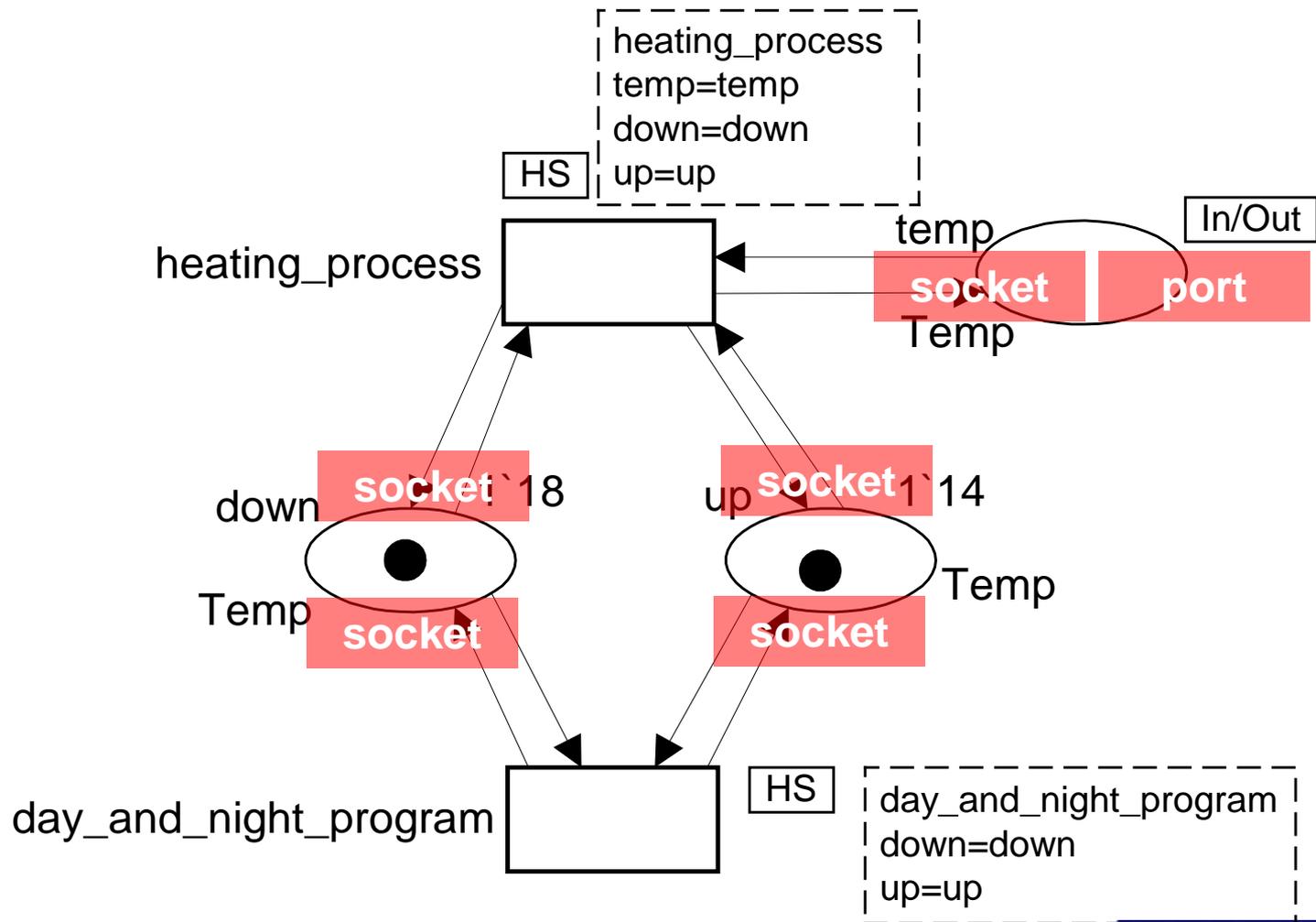


Top-level page: *main*

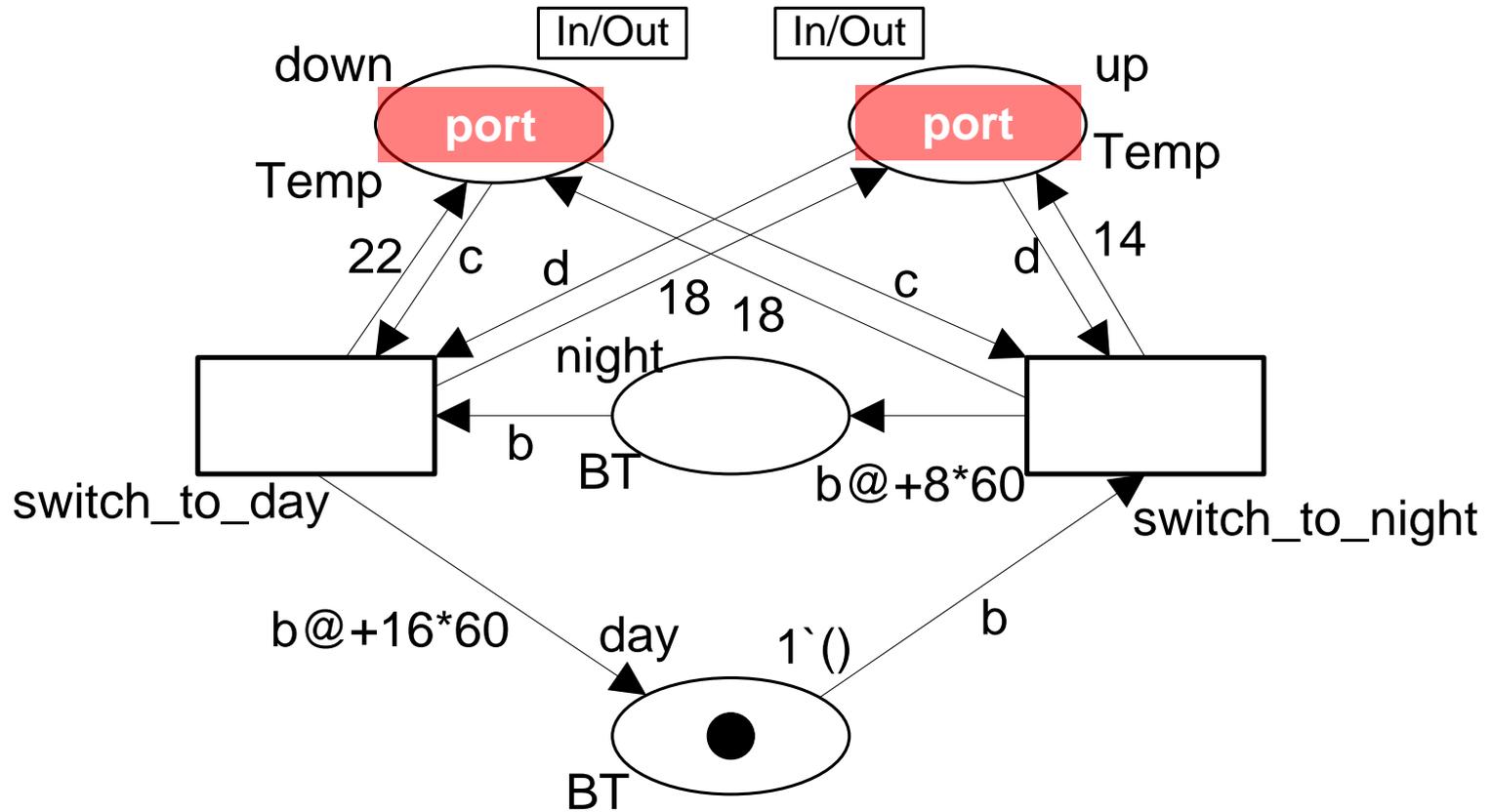
```
color Temp = string;  
color B = unit;  
color BT = B timed;  
var t:Temp;  
var a:B;  
var b:BT;  
var c:Temp;
```



Subpage: *complete_system*



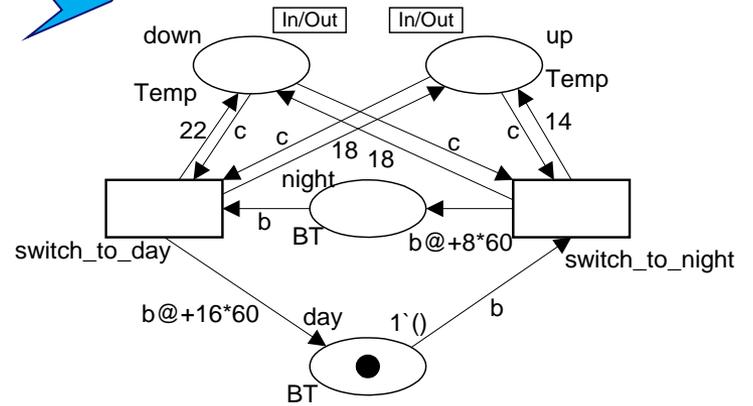
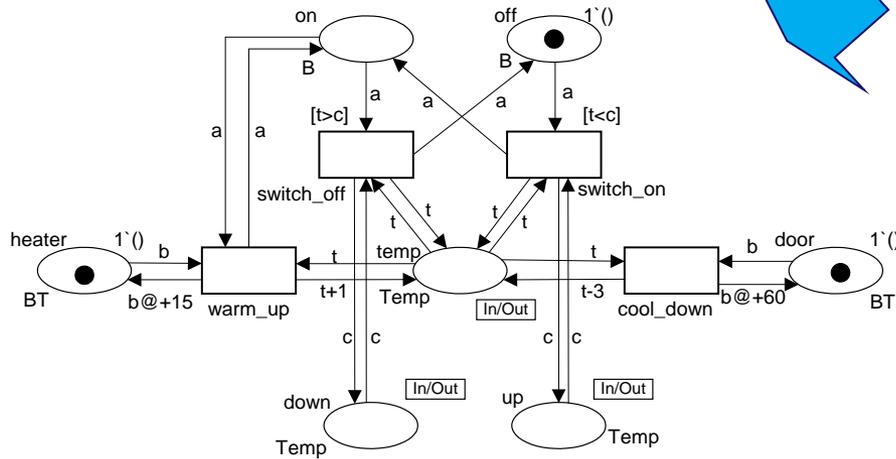
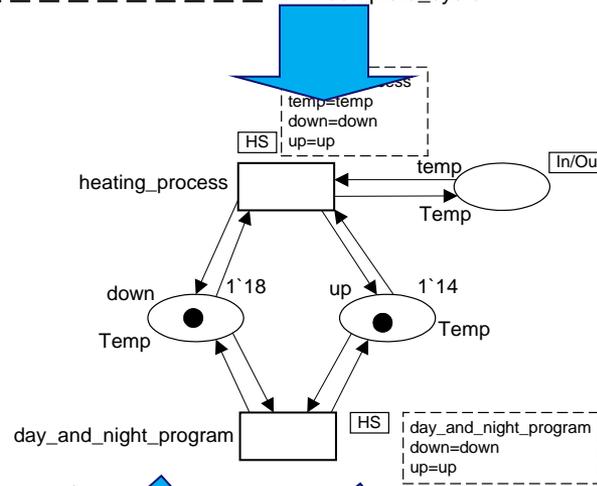
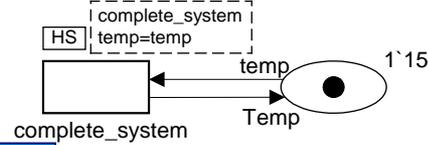
Subpage: *day_and_night_program*



Overview

```

color Temp = string;
color B = unit;
color BT = B timed;
var t:Temp;
var a:B;
var b:BT;
var c:Temp;
    
```



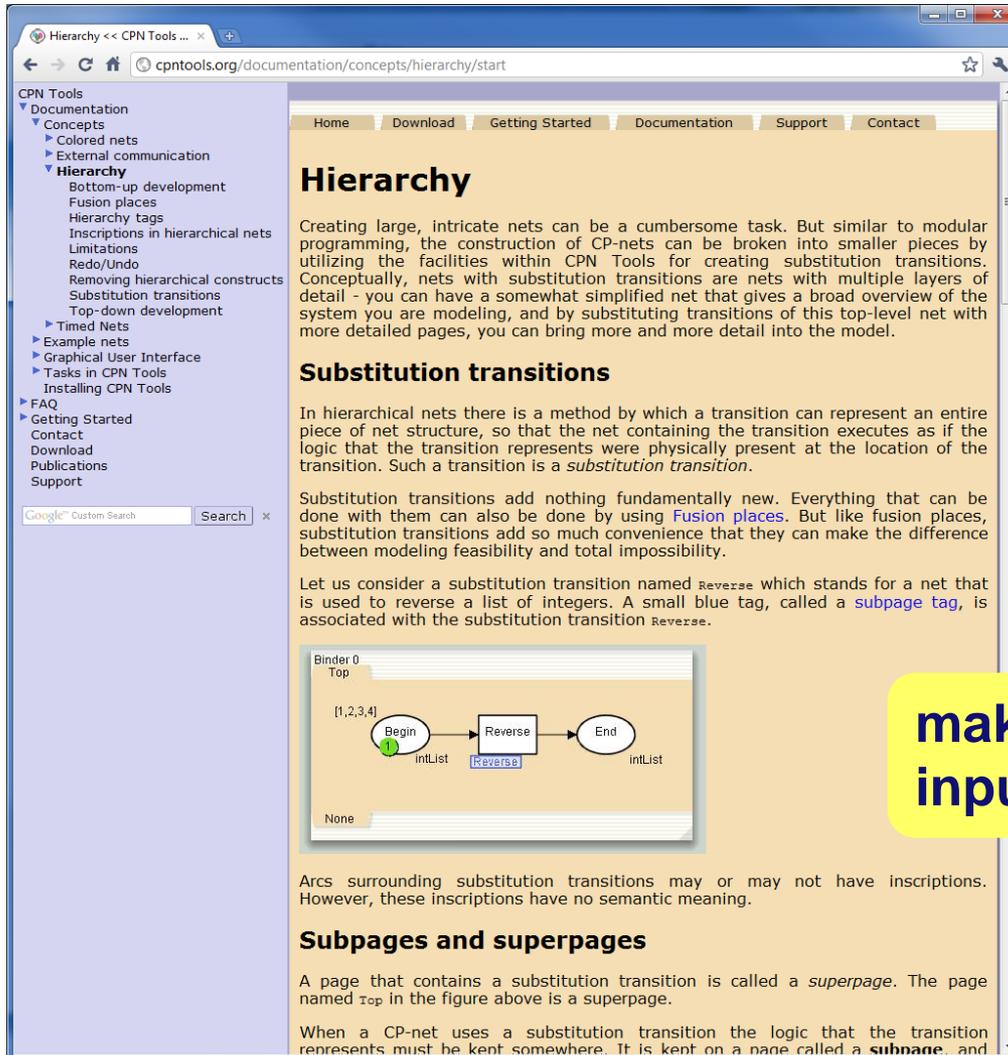
CPN Tools

move to subpage

assign subpage

unfold

connect



The screenshot shows the CPN Tools website documentation page for Hierarchy. The page title is "Hierarchy" and it is part of the "Documentation" section. The main content area contains the following text:

Hierarchy

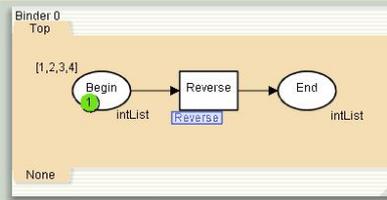
Creating large, intricate nets can be a cumbersome task. But similar to modular programming, the construction of CP-nets can be broken into smaller pieces by utilizing the facilities within CPN Tools for creating substitution transitions. Conceptually, nets with substitution transitions are nets with multiple layers of detail - you can have a somewhat simplified net that gives a broad overview of the system you are modeling, and by substituting transitions of this top-level net with more detailed pages, you can bring more and more detail into the model.

Substitution transitions

In hierarchical nets there is a method by which a transition can represent an entire piece of net structure, so that the net containing the transition executes as if the logic that the transition represents were physically present at the location of the transition. Such a transition is a *substitution transition*.

Substitution transitions add nothing fundamentally new. Everything that can be done with them can also be done by using [Fusion places](#). But like fusion places, substitution transitions add so much convenience that they can make the difference between modeling feasibility and total impossibility.

Let us consider a substitution transition named `Reverse` which stands for a net that is used to reverse a list of integers. A small blue tag, called a [subpage tag](#), is associated with the substitution transition `Reverse`.



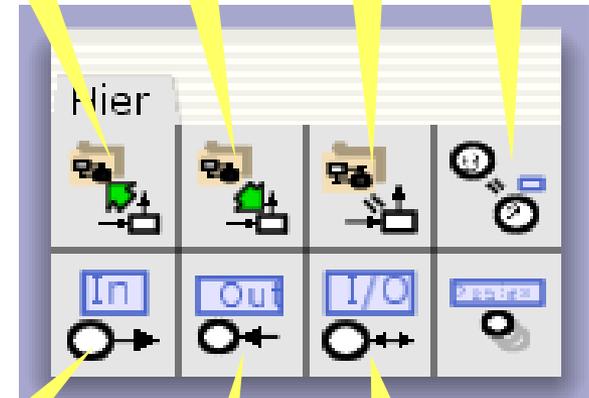
The diagram shows a substitution transition named "Reverse" with a subpage tag "[1,2,3,4]". It has an input place "Begin" and an output place "End", both with "intList" inscriptions. The transition is represented by a box with the name "Reverse" and a subpage tag "[Reverse]".

Arcs surrounding substitution transitions may or may not have inscriptions. However, these inscriptions have no semantic meaning.

Subpages and superpages

A page that contains a substitution transition is called a *superpage*. The page named `top` in the figure above is a superpage.

When a CP-net uses a substitution transition the logic that the transition represents must be kept somewhere. It is kept on a page called a **subpage**, and

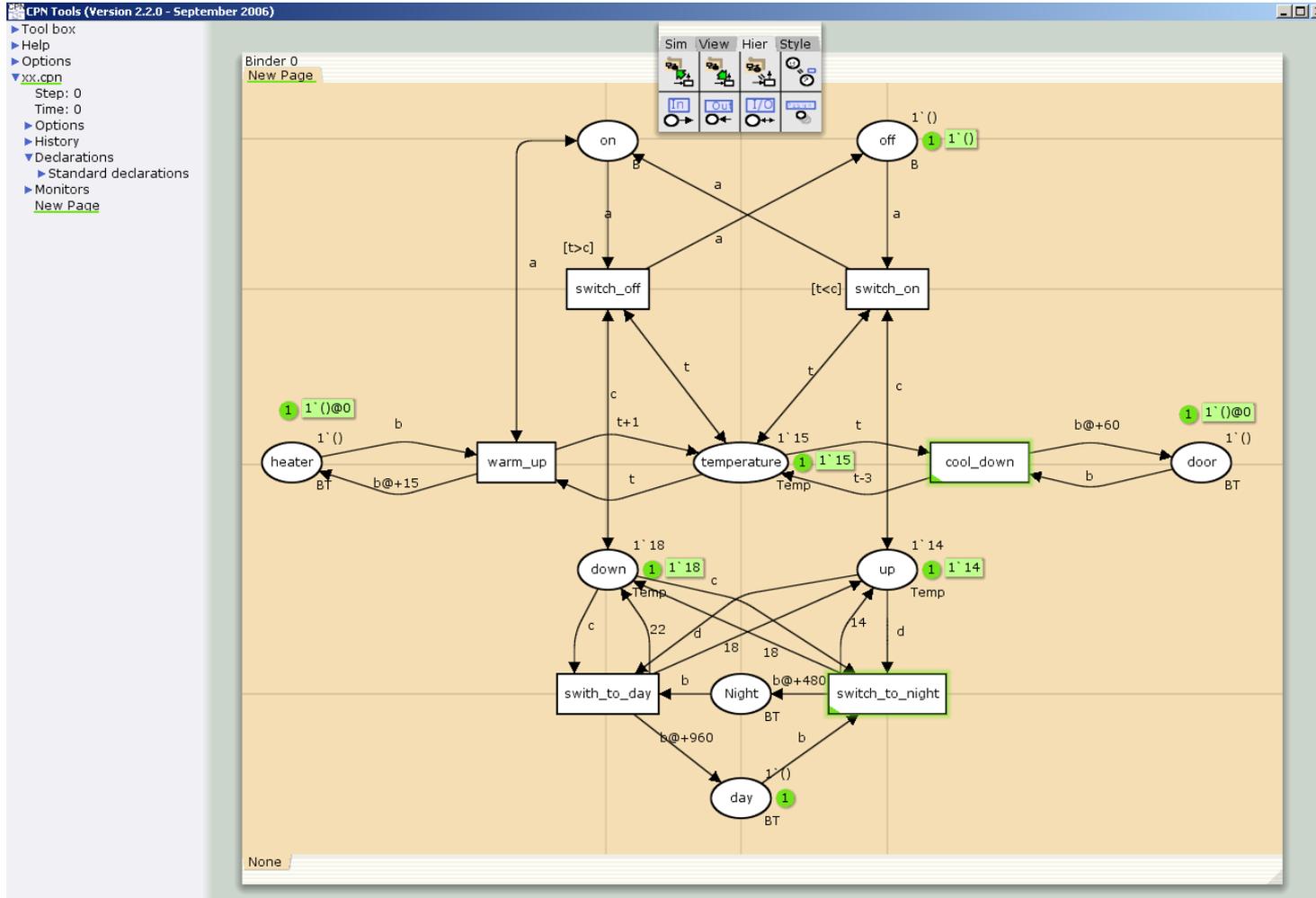


make input port

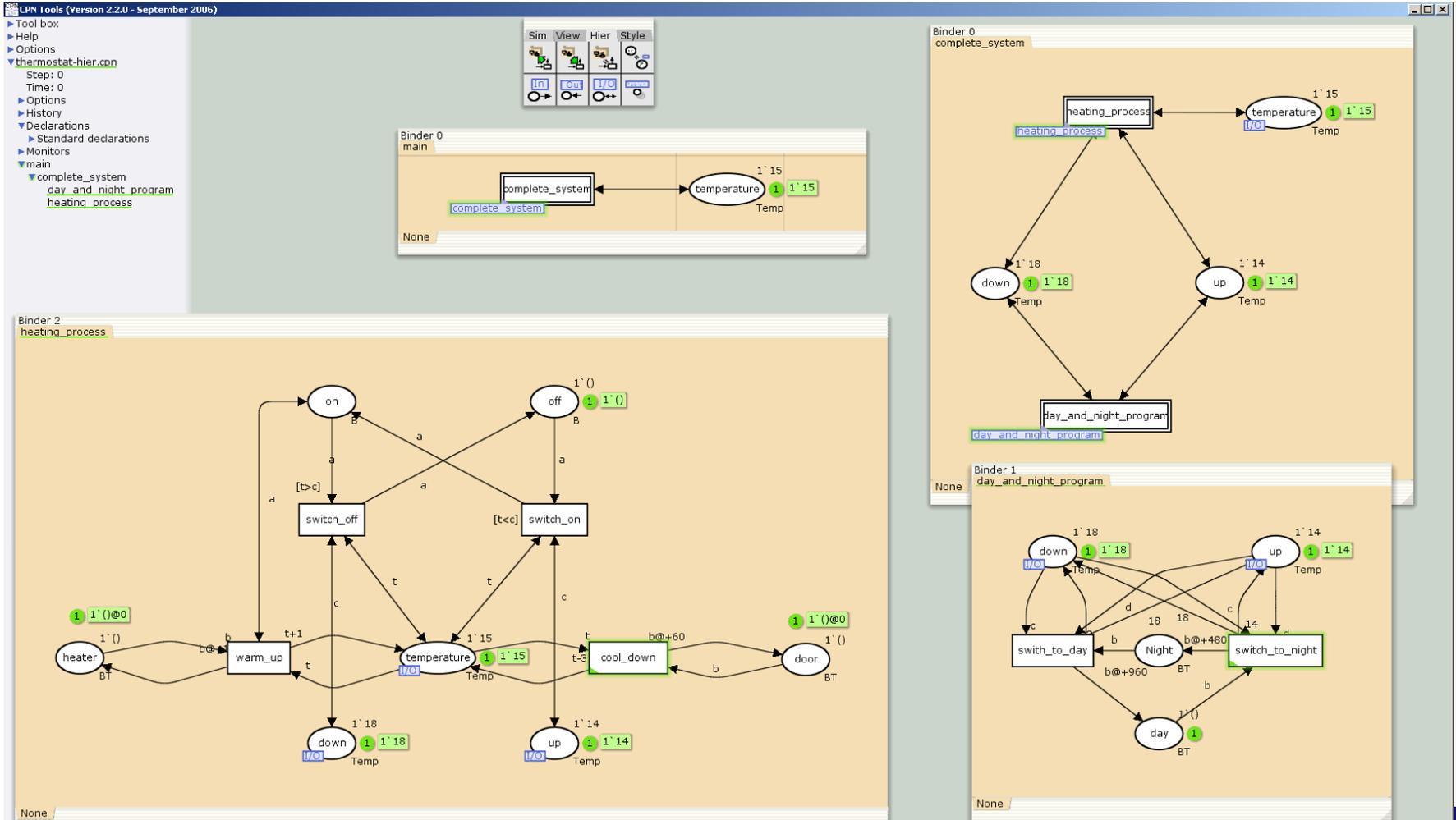
make I/O port

make output port

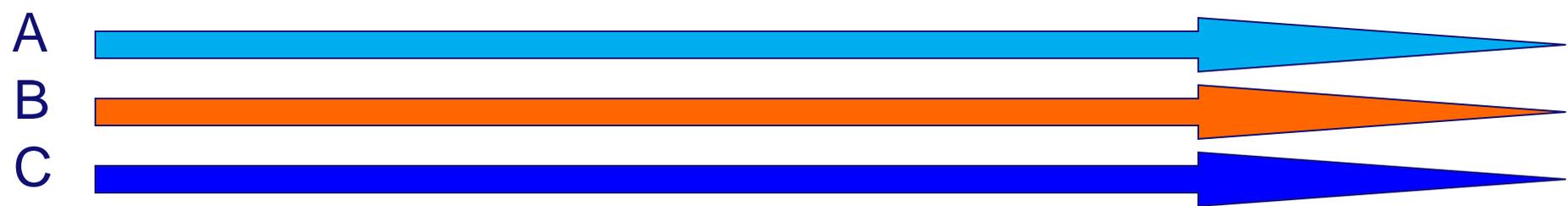
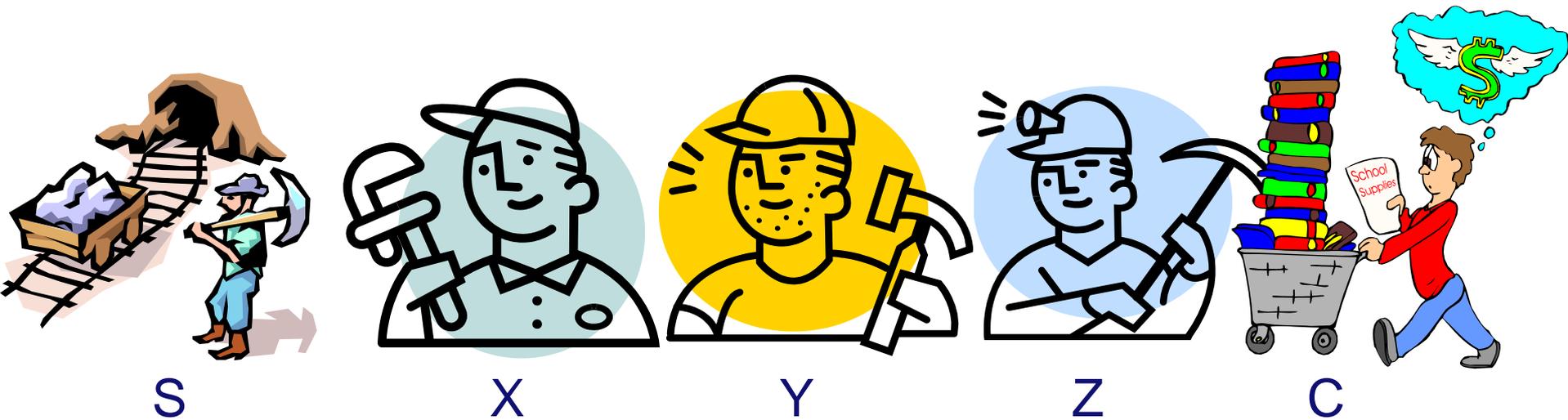
CPN Tools (flat model)



With hierarchy



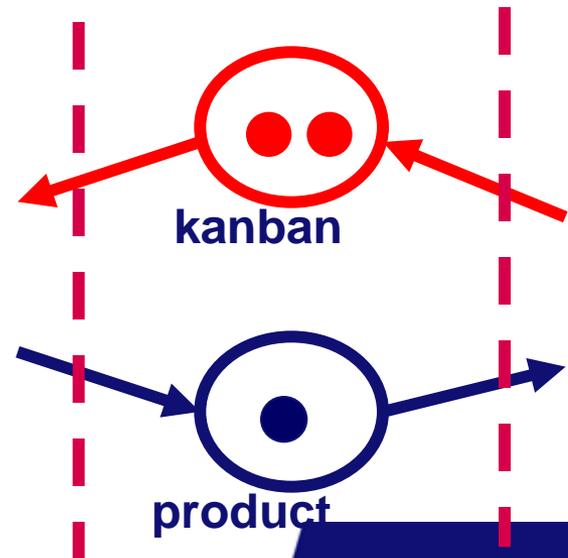
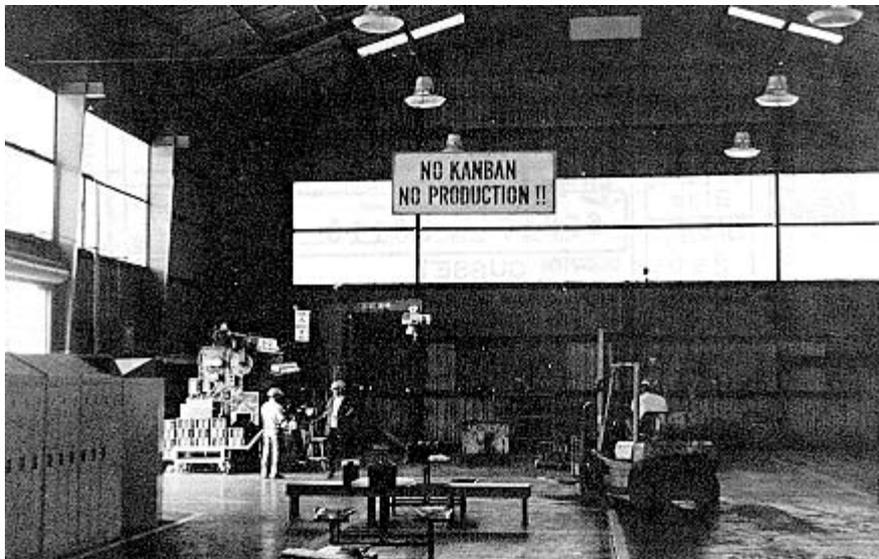
Example: Production system



Kanban



- Kanbans are a means to achieve JIT (Just-In-Time).
- Japanese word for card.
- Attributed to Taiichi Ohno (the father of the Toyota Production System)



Two beer kanban



Data

	X	Y	Z
A	2	5	8
B	3	6	9
C	4	7	1

Processing times

X	2
Y	3
Z	4

*Resources per
work center*

A	2
B	1
C	2

*Replenishment
lead times*

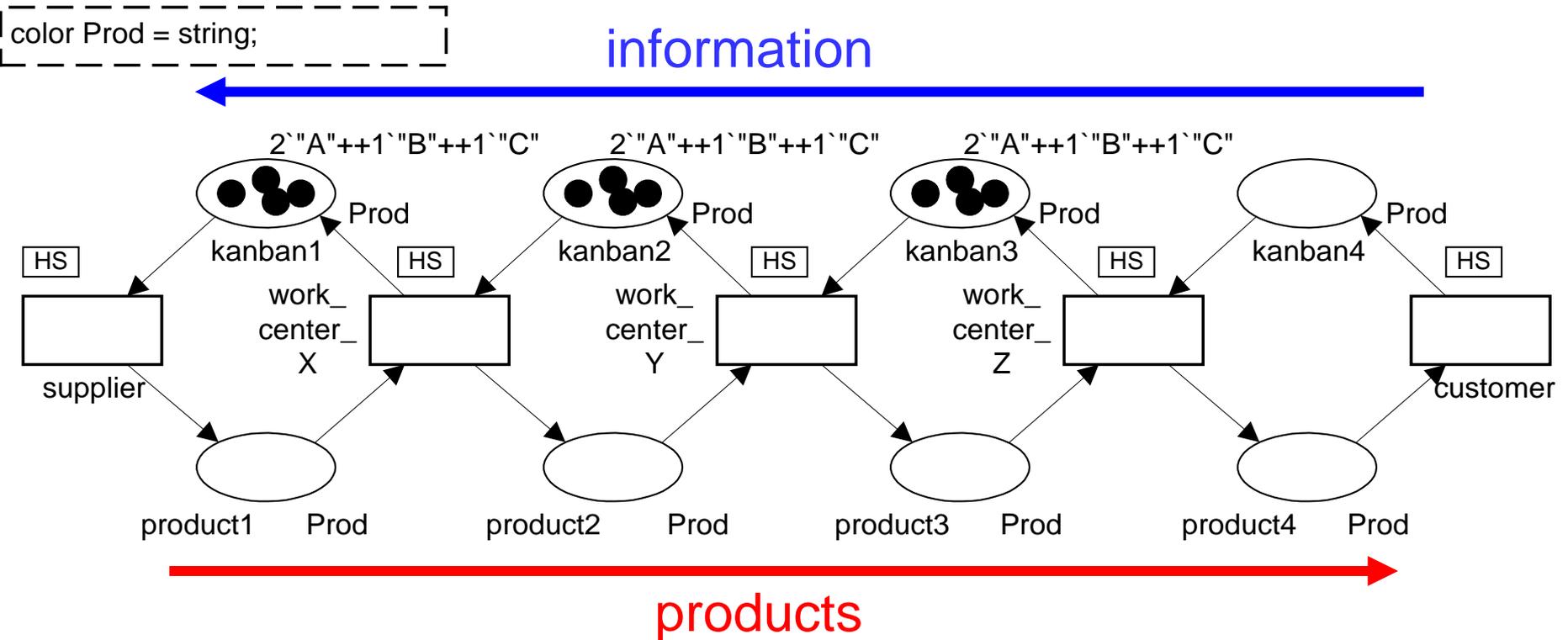
A	2
B	1
C	1

*Kanbans in-
between work
centers*

A	7
B	9
C	8

*Time in-
between
subsequent
orders*

Initial model



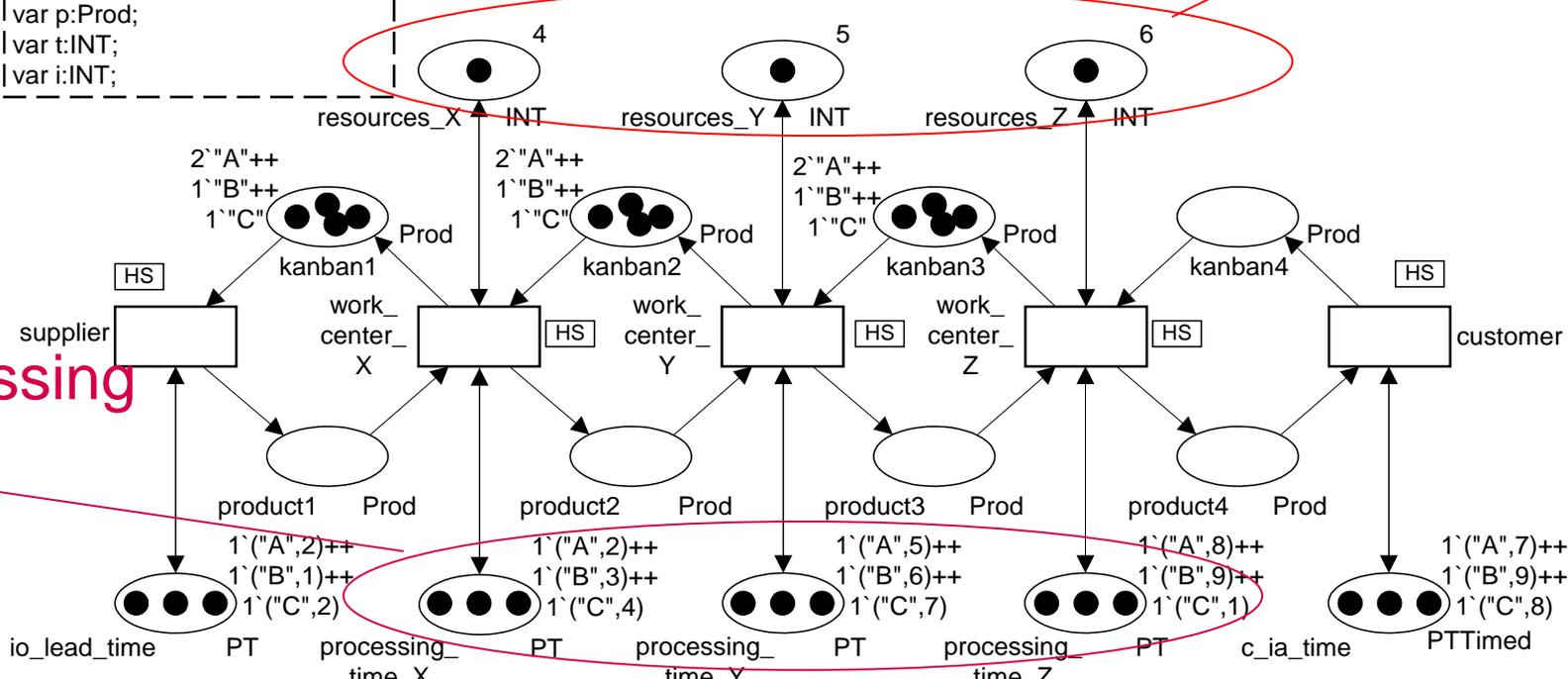
How to configure each work center if they share the same page definition???

Top level page: *main*

```

color INT = int;
color Prod = string;
color PT = product Prod * INT;
color PTimed = Prod timed;
color PTTimed = PT timed;
var p:Prod;
var t:INT;
var i:INT;
    
```

resources



processing times

```

supplier
io_lead_time = io_lead_time
kanban_in = kanban1
product_out = product1
    
```

```

work_center
processing_time =
processing_time_X
resources = resources_X
kanban_in = kanban1
product_in = product1
kanban_out = kanban1
product_out = product2
    
```

```

work_center
processing_time =
processing_time_Y
resources = resources_Y
kanban_in = kanban3
product_in = product2
kanban_out = kanban2
product_out = product3
    
```

```

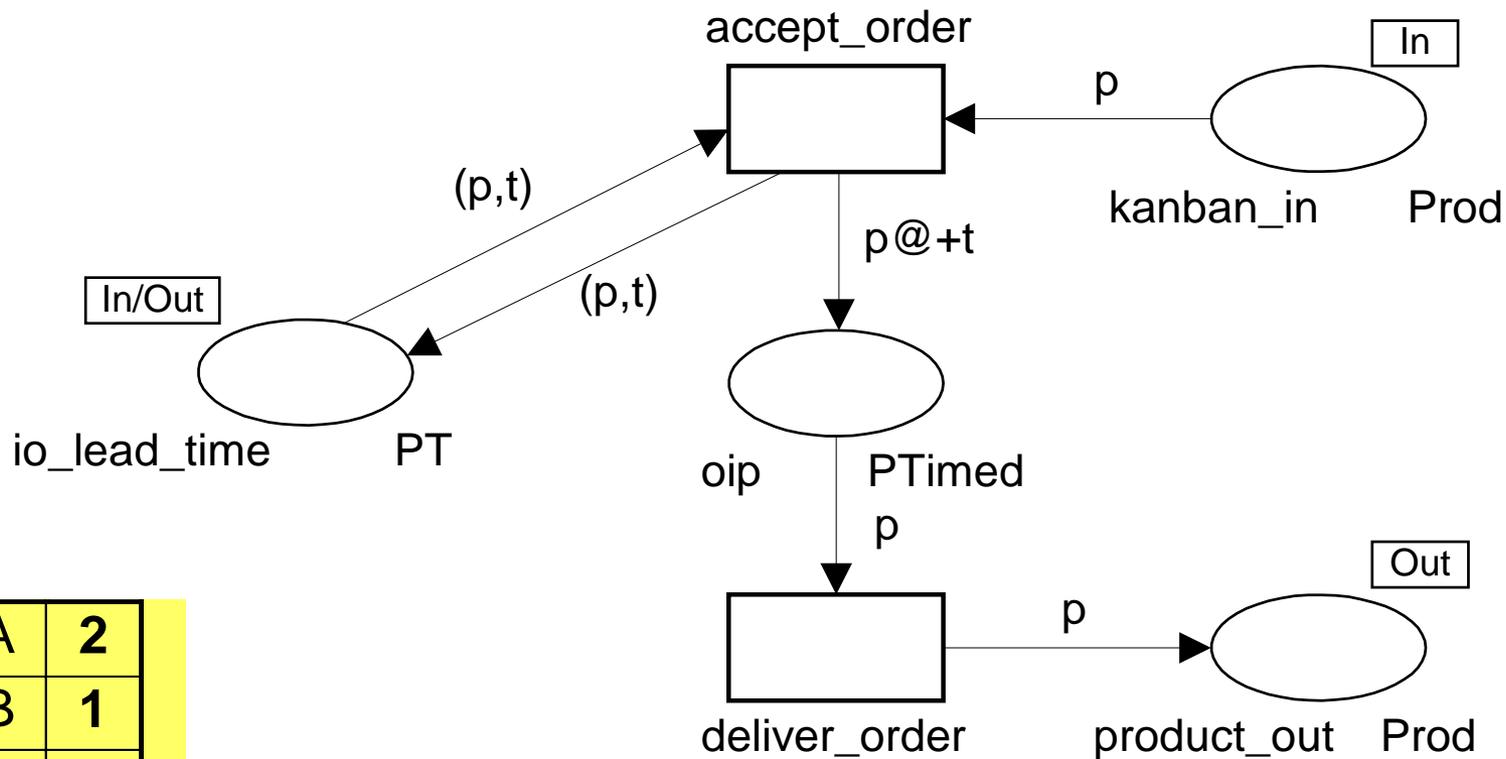
work_center
processing_time =
processing_time_Z
resources = resources_Z
kanban_in = kanban4
product_in = product3
kanban_out = kanban3
product_out = product4
    
```

```

customer
c_ia_time = c_ia_time
kanban_out = kanban4
product_in = product4
    
```

port socket assignments

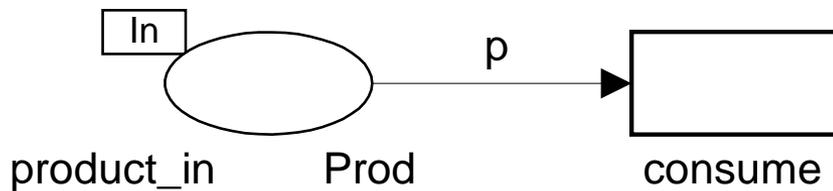
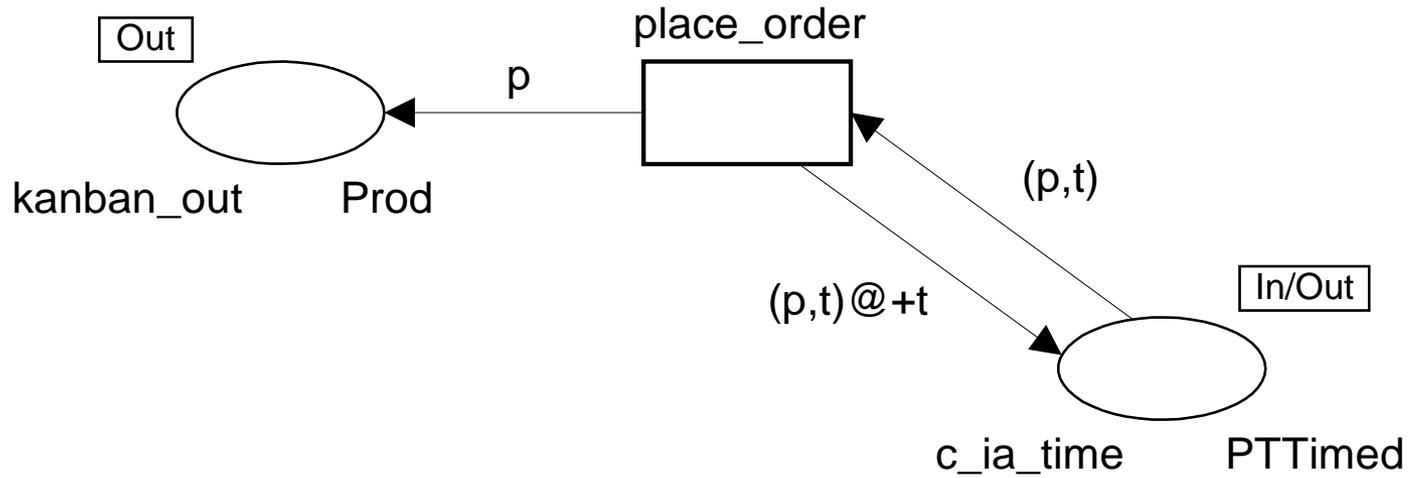
Sub page: *supplier*



A	2
B	1
C	2

*Replenishment
lead times*

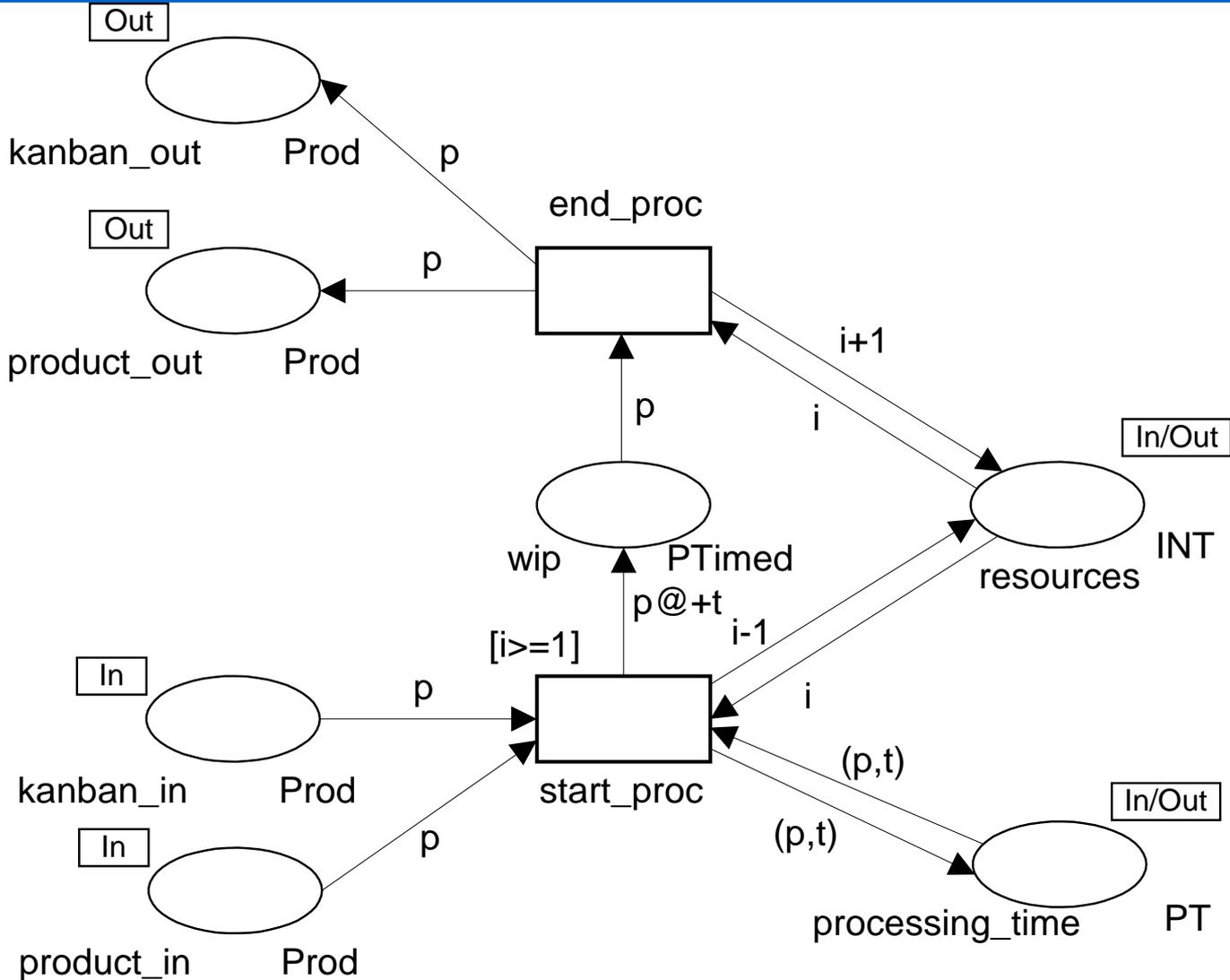
Sub page: *customer*



A	7
B	9
C	8

Time in-between subsequent orders

Sub page: *work_center*



X	2
Y	3
Z	4

Resources per work center

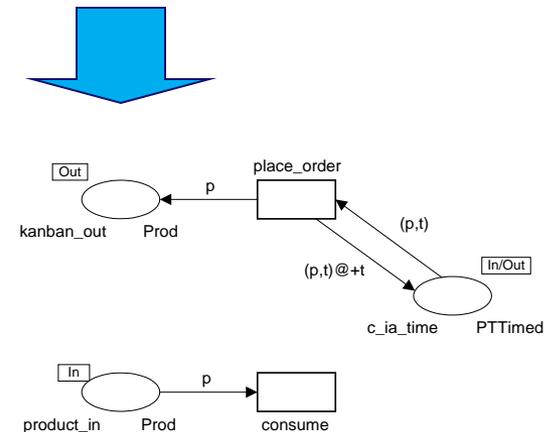
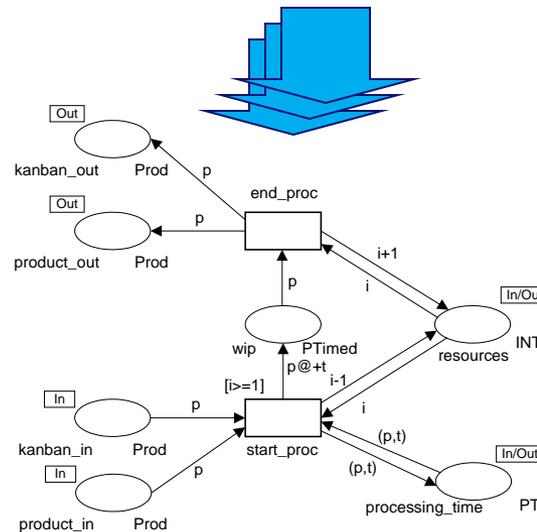
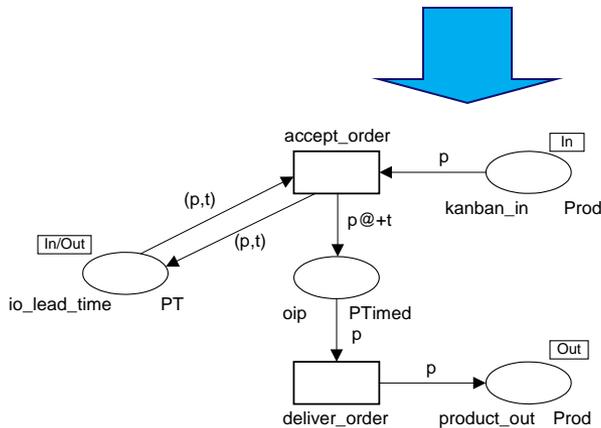
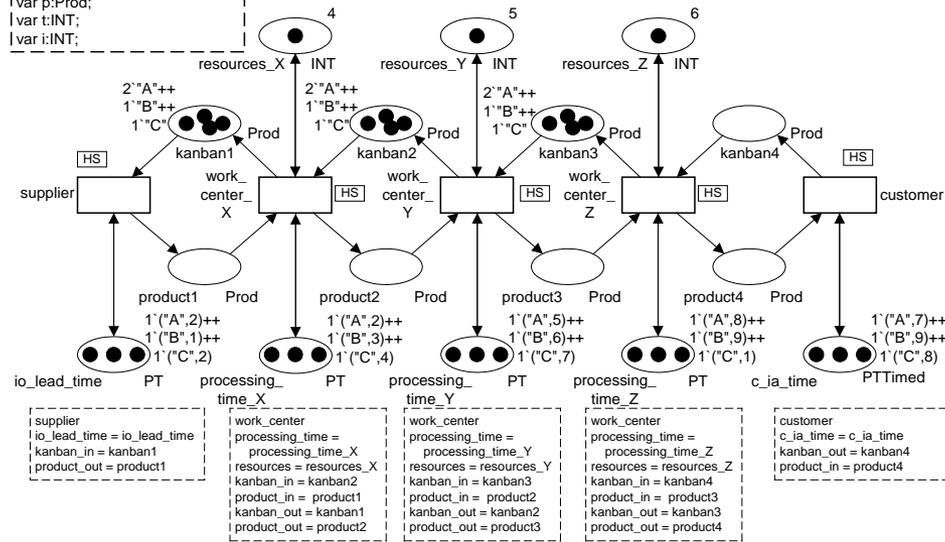
	X	Y	Z
A	2	5	8
B	3	6	9
C	4	7	1

Processing times

Overview

```

color INT = int;
color Prod = string;
color PT = product Prod * INT;
color PTTimed = Prod timed;
color PTTimed = PT timed;
var p:Prod;
var t:INT;
var i:INT;
    
```



*One page definition,
three page instances!!*

CPN tools

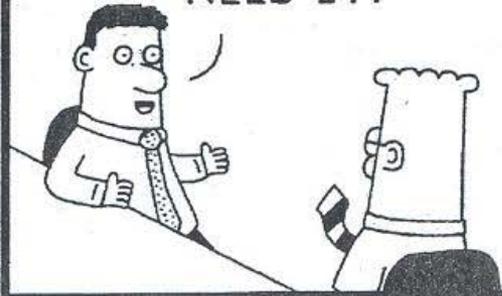
CPN Tools

- History
- Tool box
 - Auxiliary
 - Create
 - View
 - Hierarchy
 - Net
 - Simulation
 - Statespace
 - Style
- Options
- ps5.cpn
 - Step: 0
 - Time: 0
- Declarations
 - color INT
 - color INT = int;
 - color Product
 - color Product = string;
 - color Prod
 - color Prod = string;
 - color PT
 - color PT = product Prod * INT;
 - color PTimed
 - color PTimed = Prod timed;
 - color PTTimed
 - color PTTimed = PT timed;
 - var p
 - var p: Prod;
 - var t
 - var t: INT;
 - var i
 - var i: INT;
 - var pt
 - var pt: PT;
- main
 - supplier
 - customer
 - work_center (1)
 - work_center (2)
 - work_center (3)

The diagram shows a Petri net for a work center. It includes places like resources, kanban_in, kanban_out, start_proc, end_proc, product_in, product_out, processing_time, and wip. Transitions are labeled with actions like p, p@+t, and p, f-1. The interface also shows a menu on the left and a list of declarations and main components.

Inventory strategies: from (s,Q)-b to (R,s,S)+b

MY COMPANY IS
MOVING TO A "JUST
IN TIME" INVENTORY
STRATEGY. YOU'LL
DELIVER WHEN WE
NEED IT.



www.dilbert.com
scottadams@aol.com

SO...YOUR SUCCESS
DEPENDS ON MY
COMPANY DOING
WHAT IT PROMISES?
YOU HAVE MY DEEPEST
SYMPATHY.



I FEEL
A SHARP,
STABBING
PAIN IN
MY CHEST.

AND
SO IT
BEGINS...



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Application: Modeling logistic processes

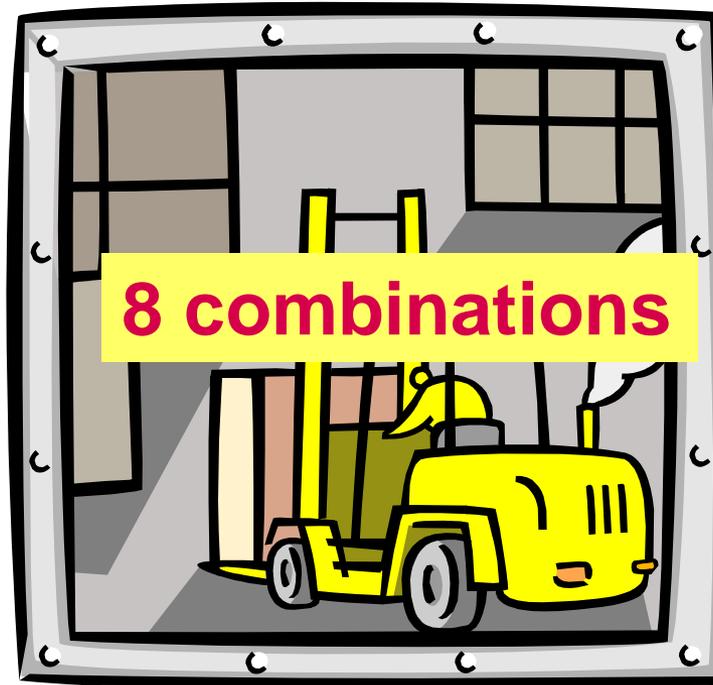
periodic or instant: R,s or s

$(s,S)-b$

$(R,s,S)-b$

$(s,Q)+b$

$(R,S)-b$



$(s,Q)-b$

$(s,S)+b$

$(R,s,S)+b$

$(R,S)+b$

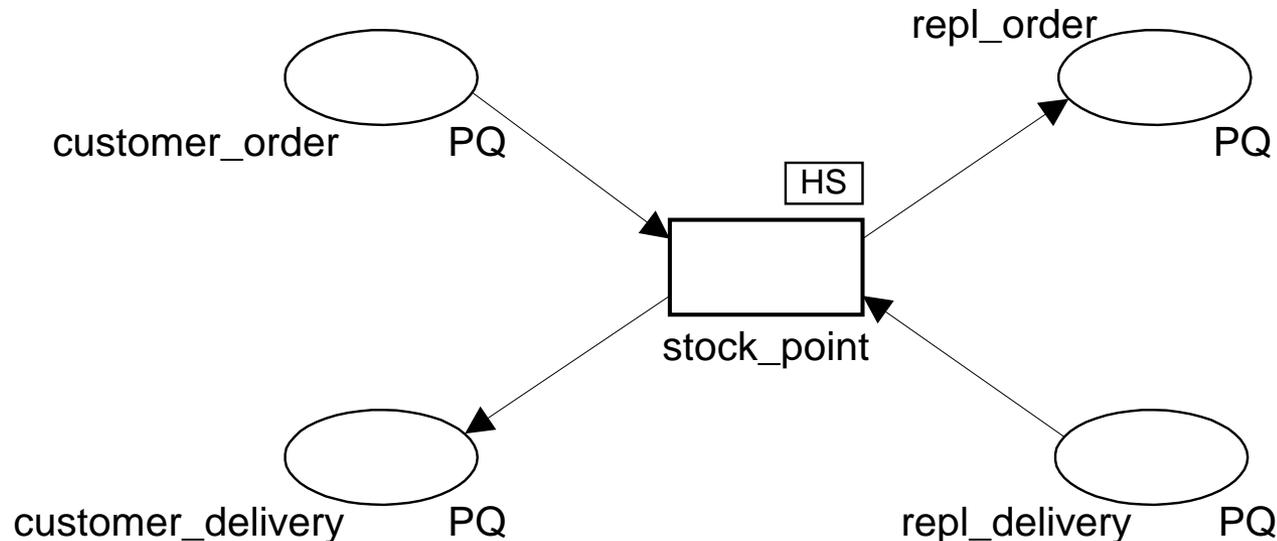
fixed or order-up-to-level: Q or S

backorders or not: -b or +b

External interface of a stock point

orders

```
color Product = string;  
color Quantity = int;  
color PQ = product Product * Quantity;
```



goods

(For simplicity we do not add external configuration places.)

Inventory policy (s,Q)-b

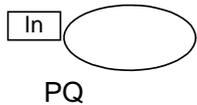
- Complete lost sales, i.e., **no backorders**.
- **Continuous review**, i.e., $R=0$.
- **Order point s** .
- **Fixed order quantity Q** .
- s and Q determined for each product individually.

Order point is compared with **inventory position**, i.e.,
on-hand stock + ordered – back orders.
(Note that back orders = 0)

Step 1: Port nodes

```
| color Product = string;  
| color Quantity = int;  
| color PQ = product Product * Quantity;
```

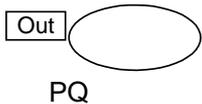
customer_order



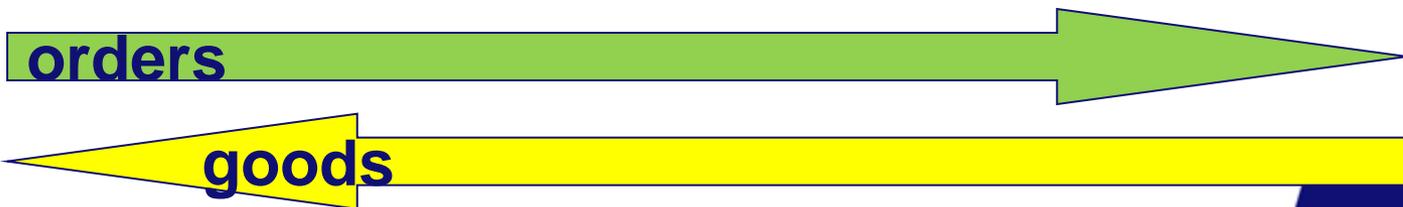
repl_order



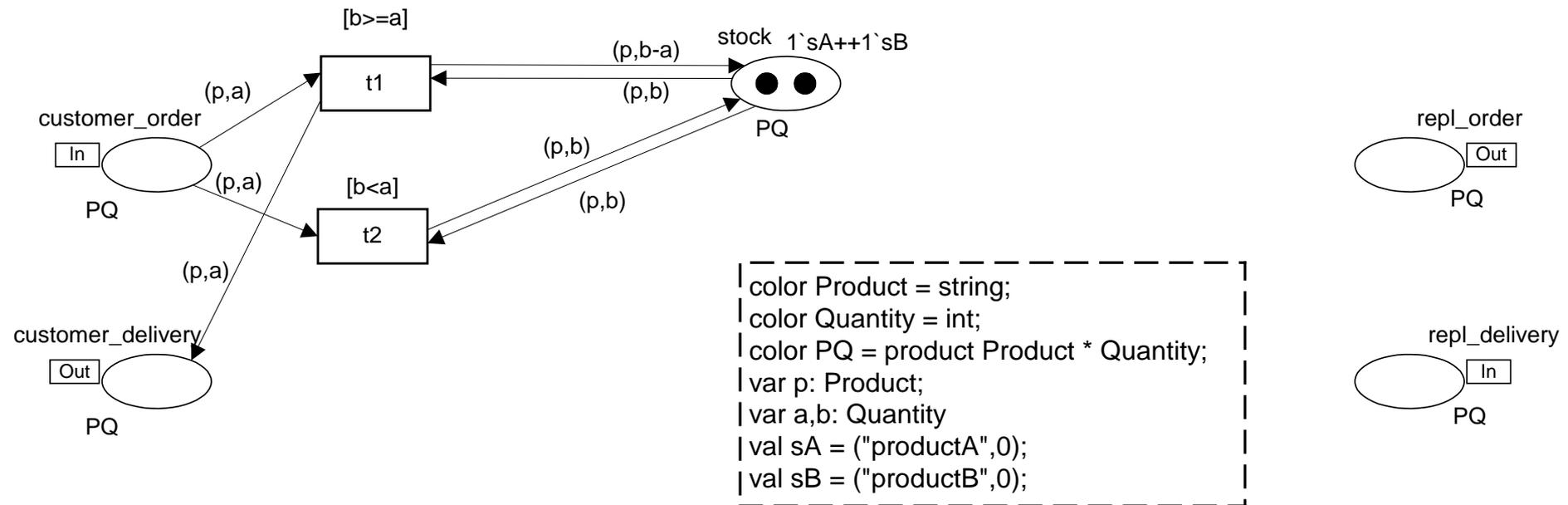
customer_delivery



repl_delivery

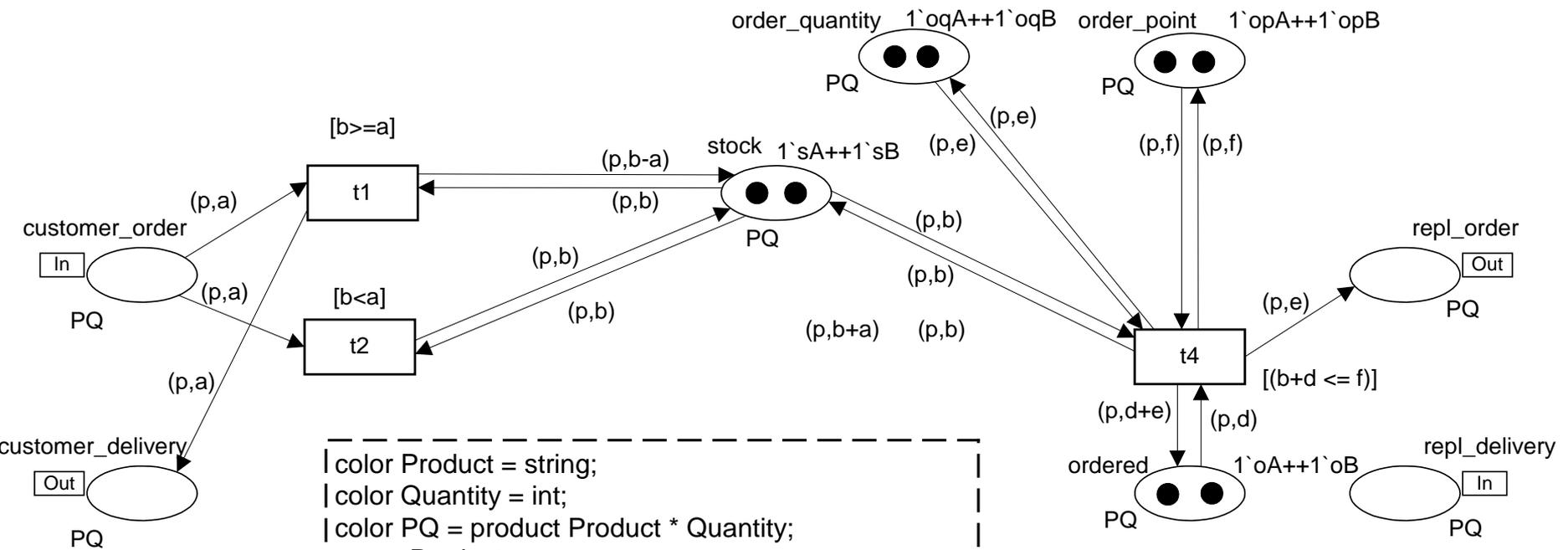


Step 2: Dealing with customer orders



In example, we assume two products: "productA" and "productB". Of course the structure of the model stays the same for any set of products.

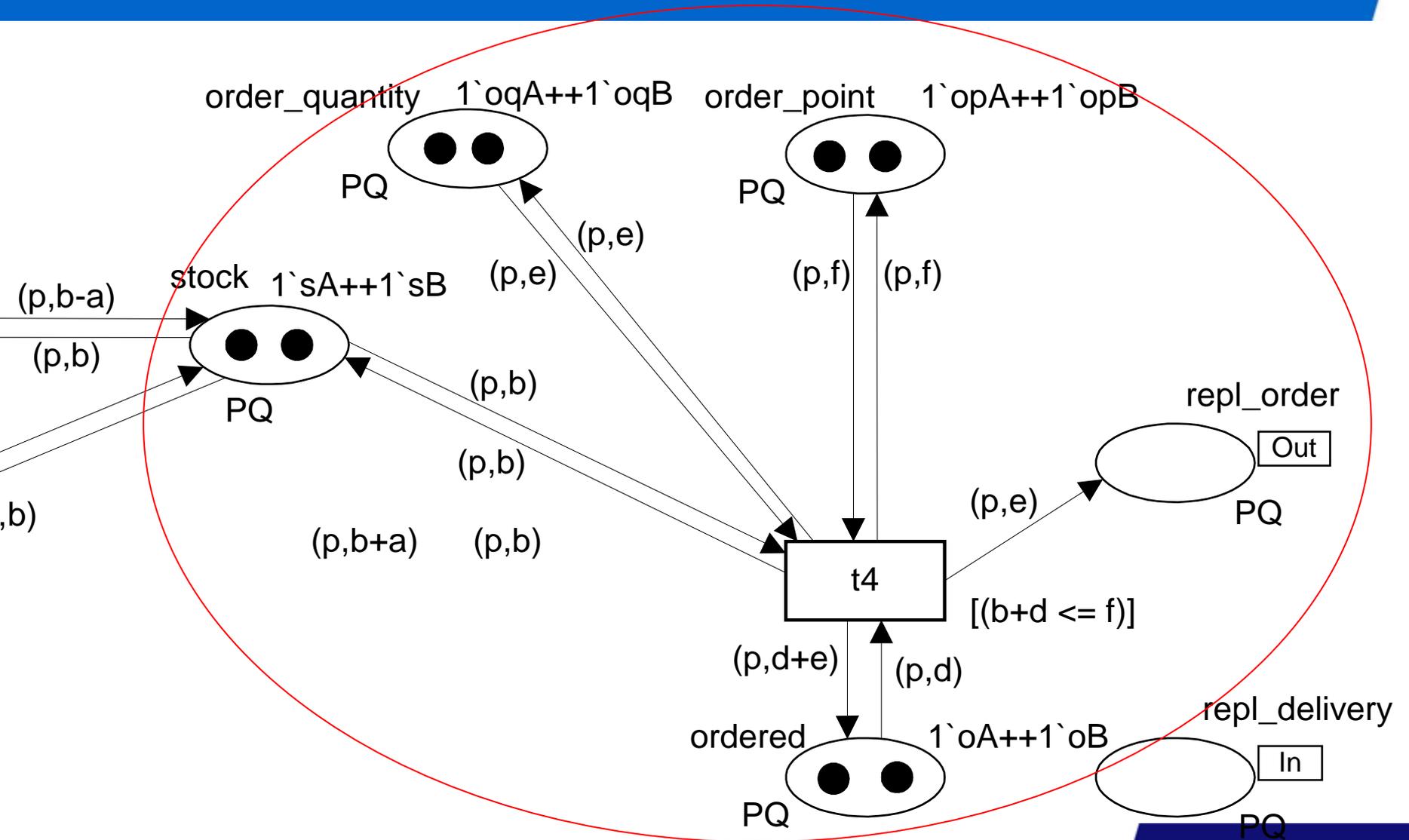
Step 3: Placing orders



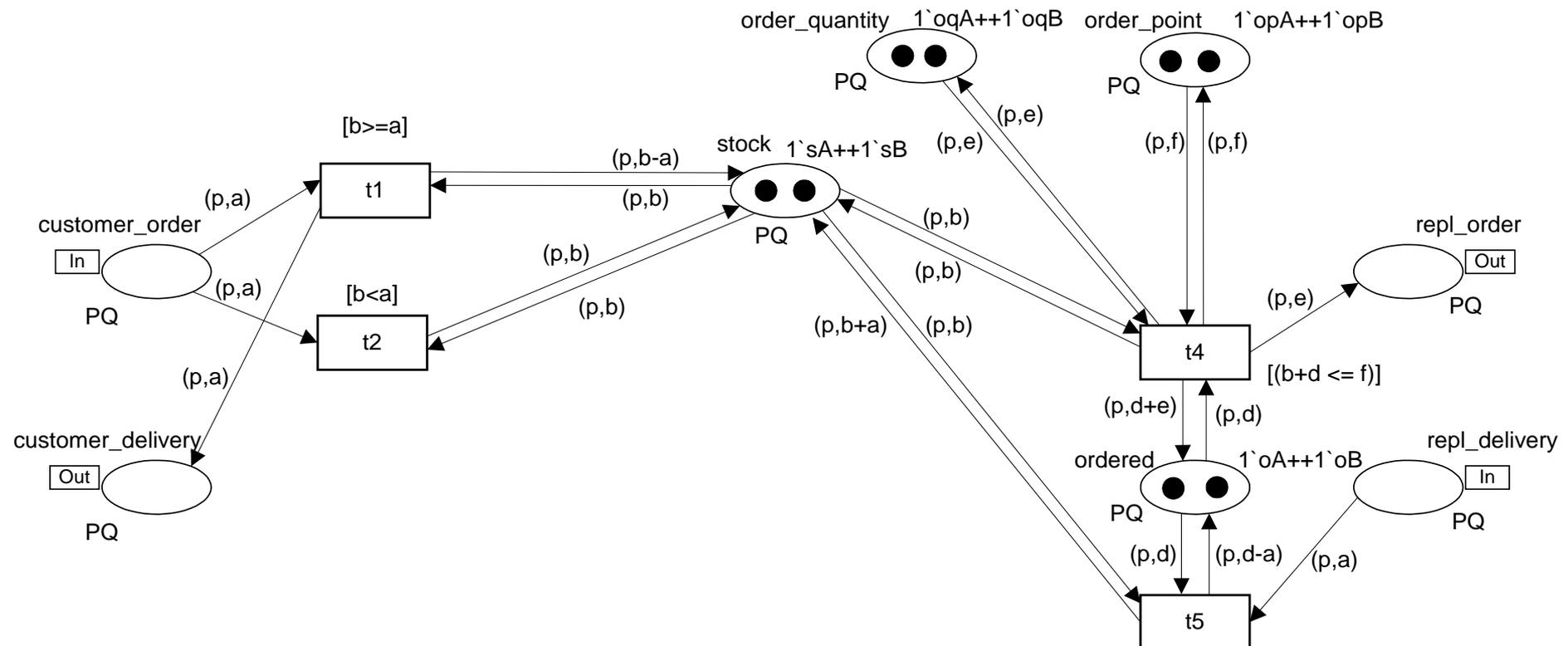
```

| color Product = string;
| color Quantity = int;
| color PQ = product Product * Quantity;
| var p: Product;
| var a,b,c,d,e,f: Quantity
| val sA = ("productA",0);
| val sB = ("productB",0);
| val oA = ("productA",0);
| val oB = ("productB",0);
| val oqA = ("productA",150);
| val oqB = ("productB",100);
| val opA = ("productA",50);
| val opB = ("productA",60);
    
```

Zoom in



Step 4: Dealing with deliveries



Inventory policy $(R,s,S)+b$

- **Complete backordering**, i.e., no lost sales.
- **Periodic review**, i.e., R is review period.
- **Order point s** .
- **Variable order quantity S** , i.e., S is the order-up-to-level.

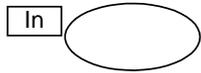
Every R time units, the inventory position is compared with s . If it is below s , the inventory position is raised to S by placing an order.

Recall: inventory position = on-hand stock + ordered – back orders.

Step 1: Port nodes (as before)

```
| color Product = string;  
| color Quantity = int;  
| color PQ = product Product * Quantity;
```

customer_order



PQ

repl_order



PQ

customer_delivery



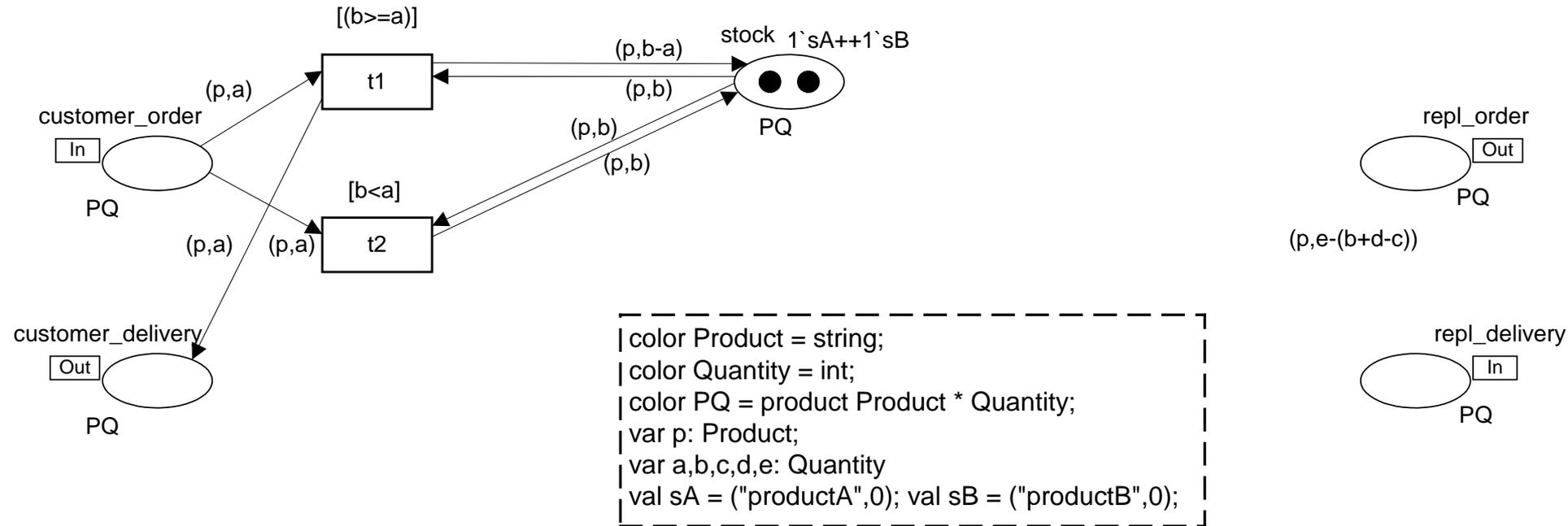
PQ

repl_delivery



PQ

Step 2: Handle customer order (as before)



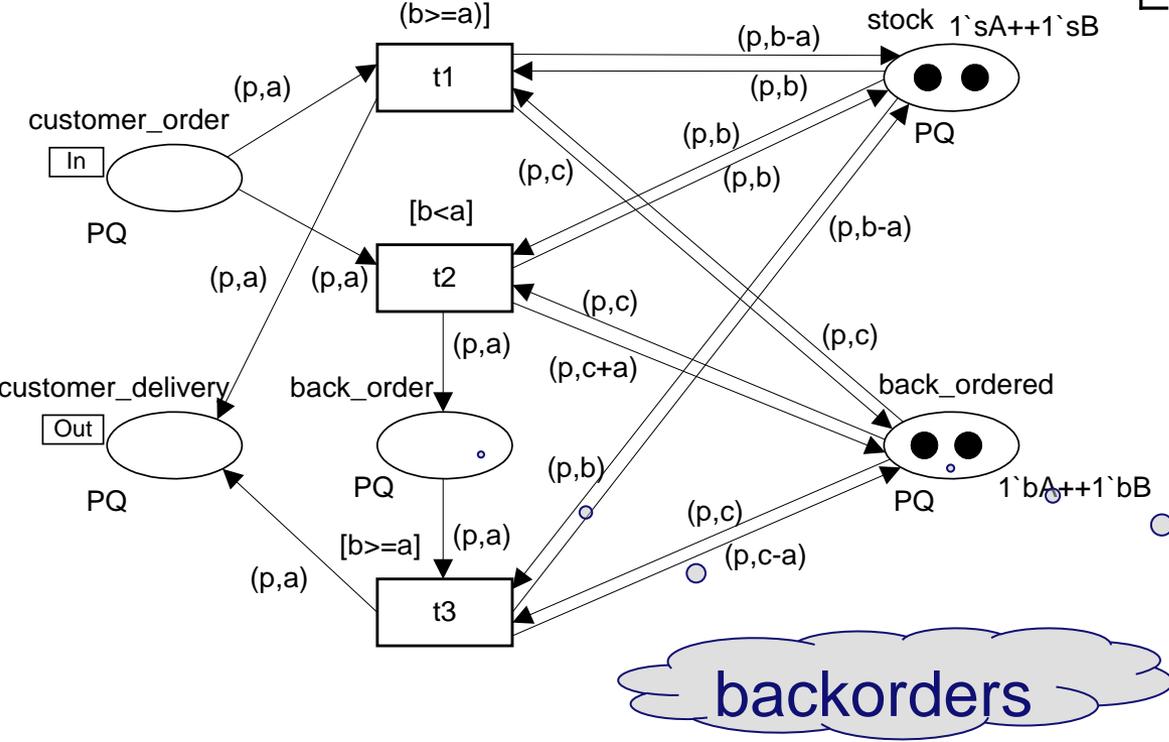
Step 3: Adding back orders

```

color Product = string;
color Quantity = int;
color PQ = product Product * Quantity;
var p: Product;
var a,b,c,d,e: Quantity
val sA = ("productA",0); val sB = ("productB",0);
val bA = ("productA",0); val bB = ("productB",0);
    
```

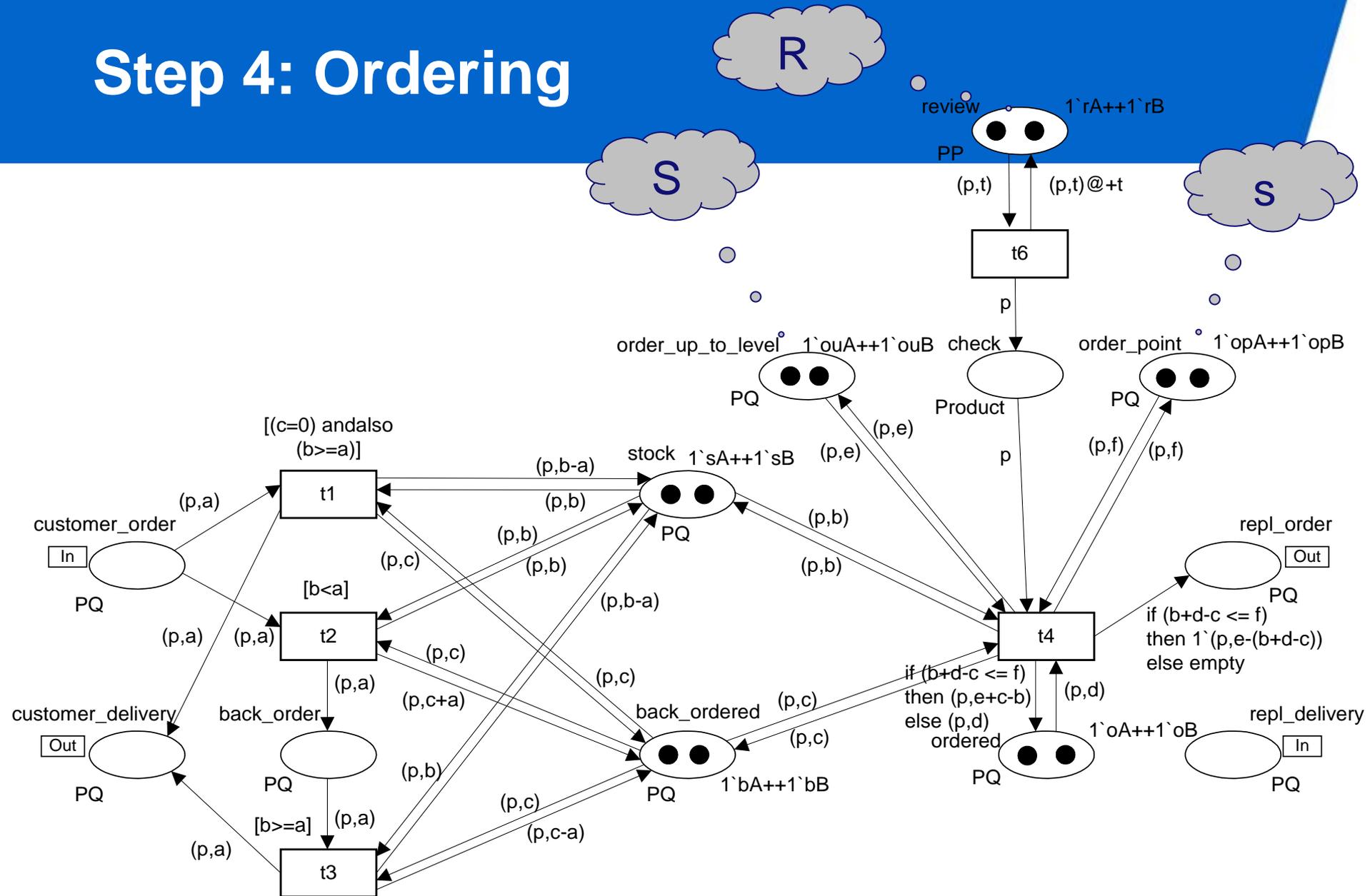


[(c=0) andalso (b>=a)]



sum

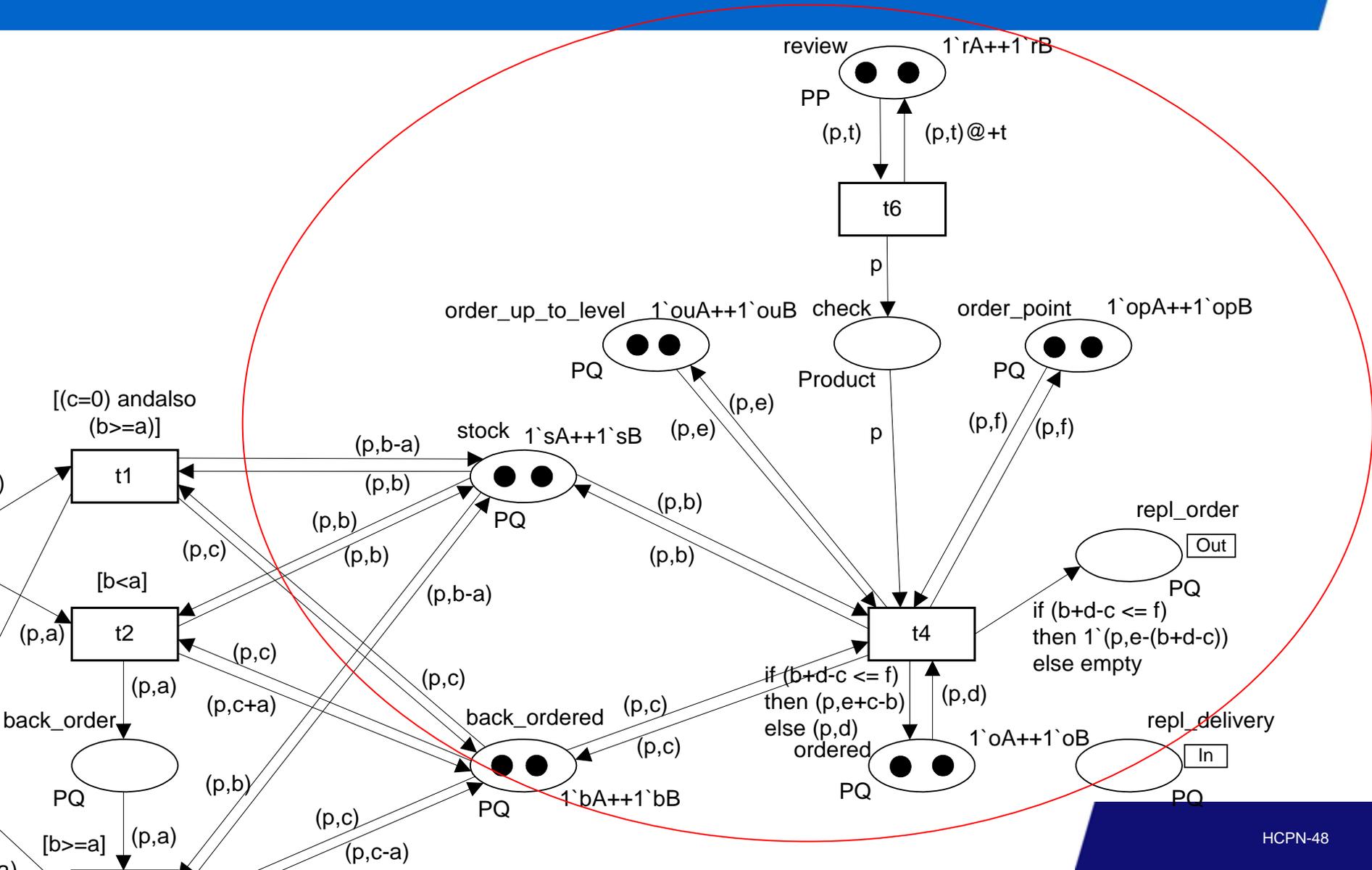
Step 4: Ordering



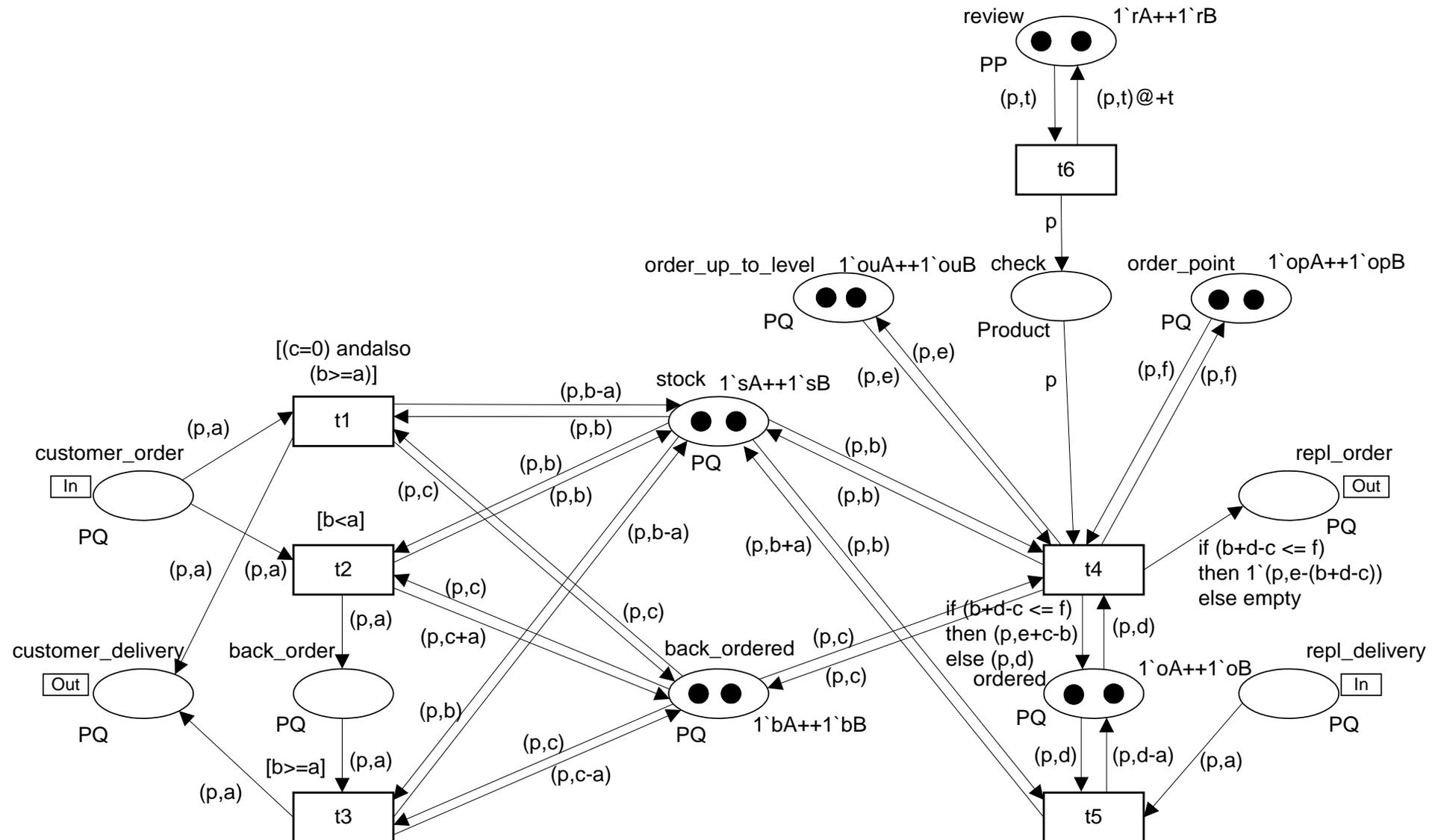
Declarations

```
| color Product = string;  
| color Quantity = int;  
| color PQ = product Product * Quantity;  
| color Period = int;  
| color PP = product Product * Period timed;  
| var p: Product;  
| var t: Period;  
| var a,b,c,d,e,f: Quantity  
| val sA = ("productA",0); val sB = ("productB",0);  
| val oA = ("productA",0); val oB = ("productB",0);  
| val bA = ("productA",0); val bB = ("productB",0);  
| val ouA = ("productA",150); val ouB = ("productB",100);  
| val opA = ("productA",50); val opB = ("productA",60);  
| val rA = ("productA",3);  
| val rB = ("productB",4);
```

Zoom in



Step 6: Dealing with deliveries



Six remaining inventory policies

$(s,S)-b$

$(R,s,S)-b$

$(s,Q)+b$

$(R,S)-b$



$(s,Q)-b$



$(s,S)+b$

$(R,s,S)+b$

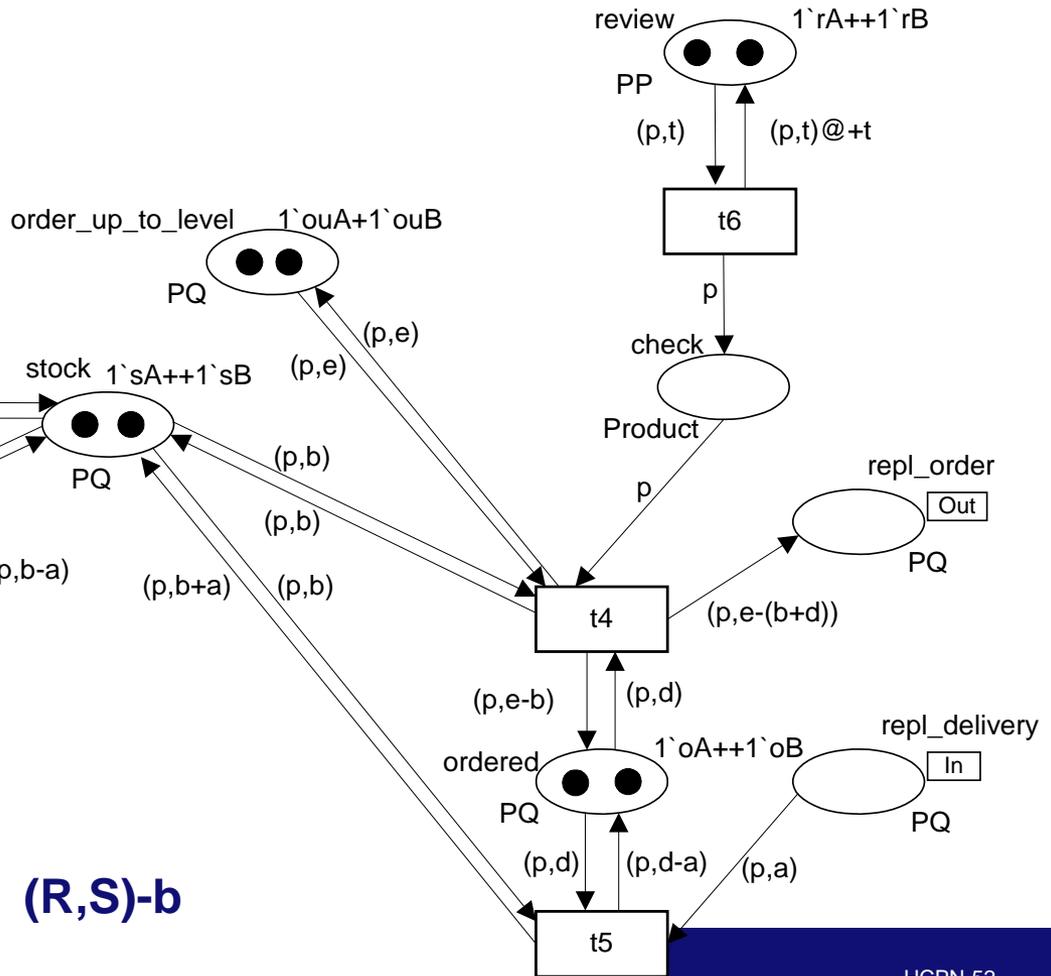


$(R,S)+b$

Question: Which inventory policy?

```

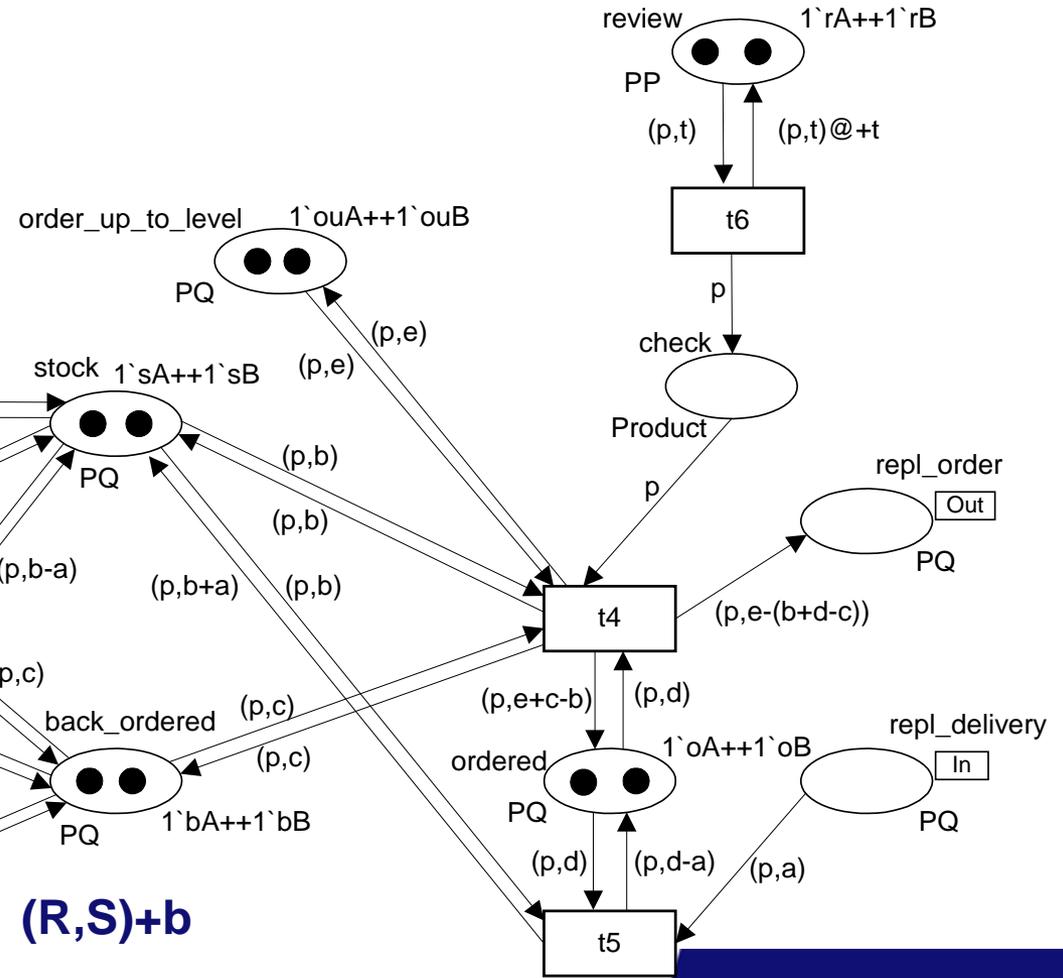
color Product = string;
color Quantity = int;
color PQ = product Product * Quantity;
color Period = int;
color PP = product Product * Period timed;
var p: Product;
var t: Period;
var a,b,c,d,e: Quantity
val sA = ("productA",0); val sB = ("productB",0);
val oA = ("productA",0); val oB = ("productB",0);
val ouA = ("productA",150); val ouB = ("productB",100);
val rA = ("productA",3); val rB = ("productB",4);
    
```



Question: Which inventory policy?

```

color Product = string;
color Quantity = int;
color PQ = product Product * Quantity;
color Period = int;
color PP = product Product * Period timed;
var p: Product;
var t: Period;
var a,b,c,d,e: Quantity
val sA = ("productA",0); val sB = ("productB",0);
val oA = ("productA",0); val oB = ("productB",0);
val bA = ("productA",0); val bB = ("productB",0);
val ouA = ("productA",150); val ouB = ("productB",100);
val rA = ("productA",3);
val rB = ("productB",4);
    
```

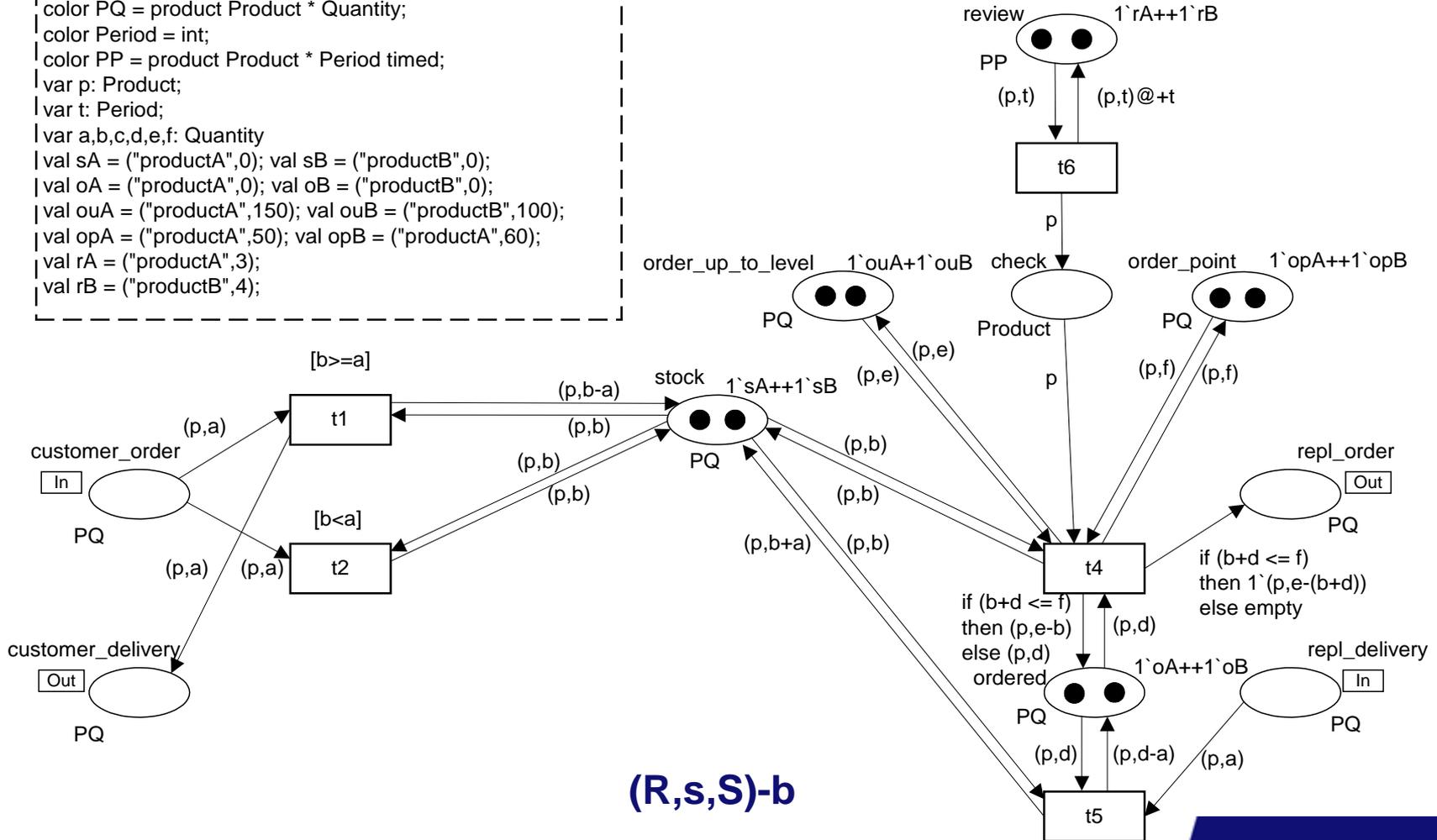


(R,S)+b

Question: Which inventory policy?

```

color Product = string;
color Quantity = int;
color PQ = product Product * Quantity;
color Period = int;
color PP = product Product * Period timed;
var p: Product;
var t: Period;
var a,b,c,d,e,f: Quantity
val sA = ("productA",0); val sB = ("productB",0);
val oA = ("productA",0); val oB = ("productB",0);
val ouA = ("productA",150); val ouB = ("productB",100);
val opA = ("productA",50); val opB = ("productA",60);
val rA = ("productA",3);
val rB = ("productB",4);
    
```

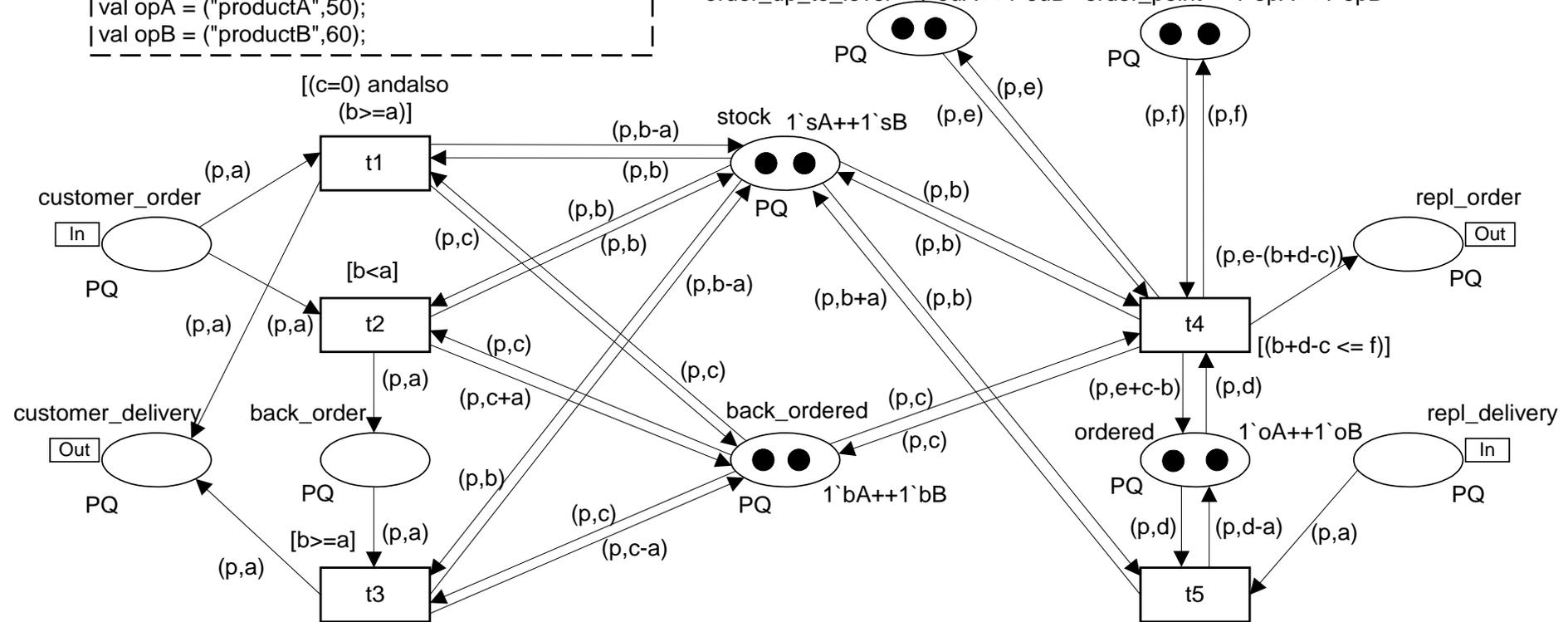


Question: Which inventory policy?

```

color Product = string;
color Quantity = int;
color PQ = product Product * Quantity;
var p: Product;
var a,b,c,d,e,f: Quantity
val sA = ("productA",0);
val sB = ("productB",0);
val oA = ("productA",0);
val oB = ("productB",0);
val bA = ("productA",0);
val bB = ("productB",0);
val ouA = ("productA",150);
val ouB = ("productB",100);
val opA = ("productA",50);
val opB = ("productB",60);
    
```

order_up_to_level 1`ouA++1`ouB order_point 1`opA++1`opB



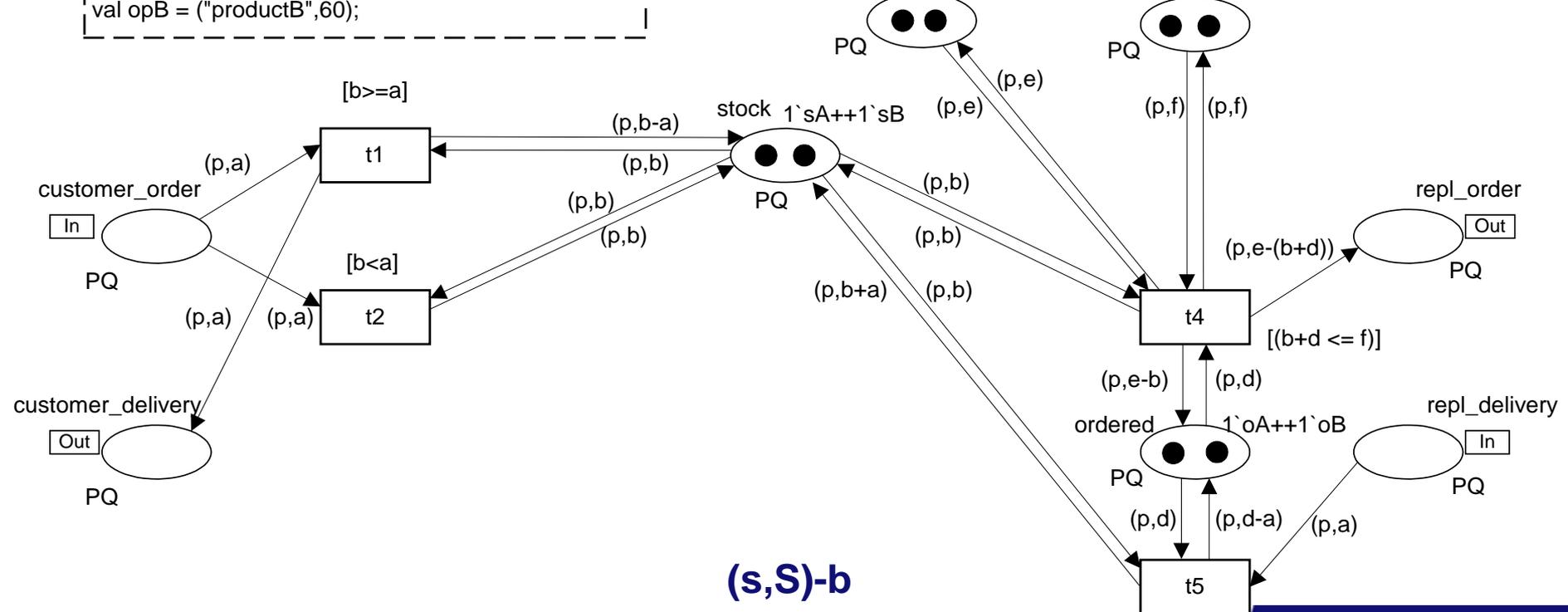
$(s,S)+b$

Question: Which inventory policy?

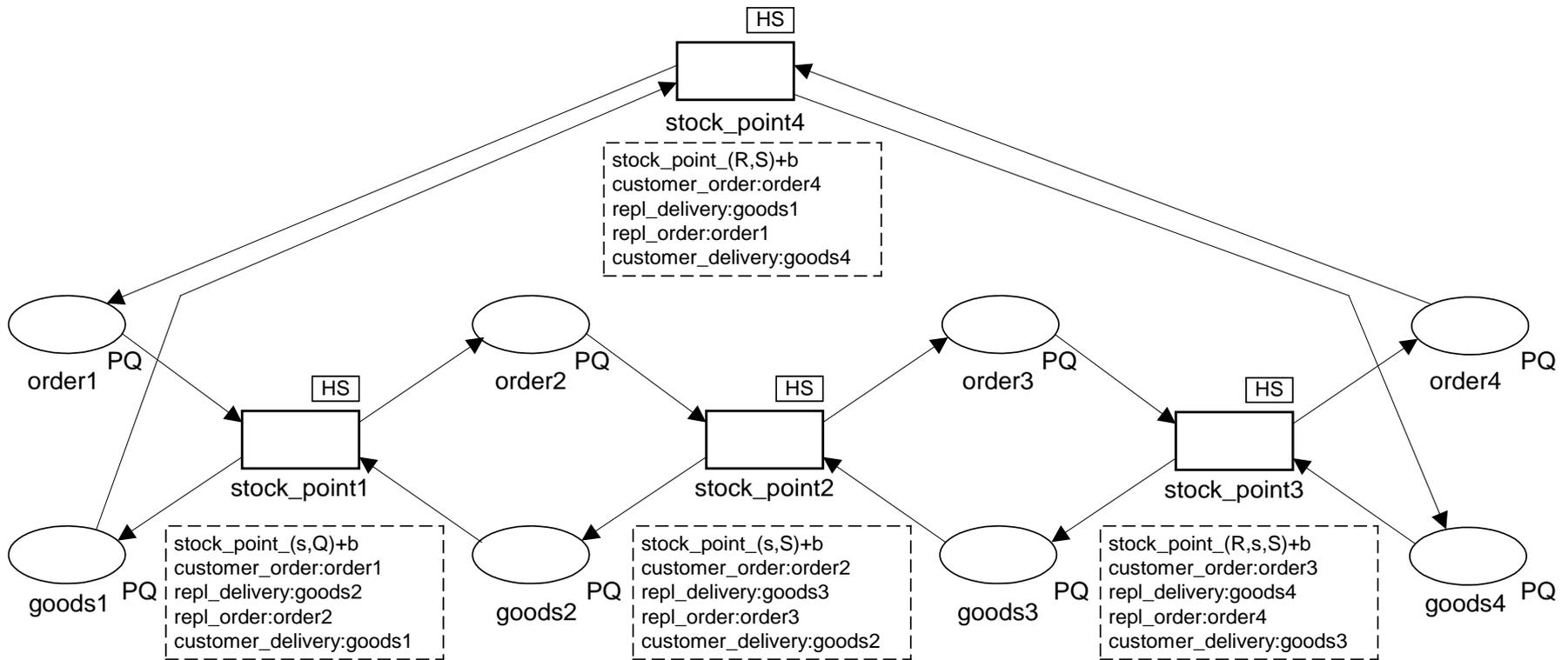
```

color Product = string;
color Quantity = int;
color PQ = product Product * Quantity;
var p: Product;
var a,b,c,d,e,f: Quantity
val sA = ("productA",0);
val sB = ("productB",0);
val oA = ("productA",0);
val oB = ("productB",0);
val ouA = ("productA",150);
val ouB = ("productB",100);
val opA = ("productA",50);
val opB = ("productB",60);
    
```

order_up_to_level 1`ouA++1`ouB order_point 1`opA++1`opB

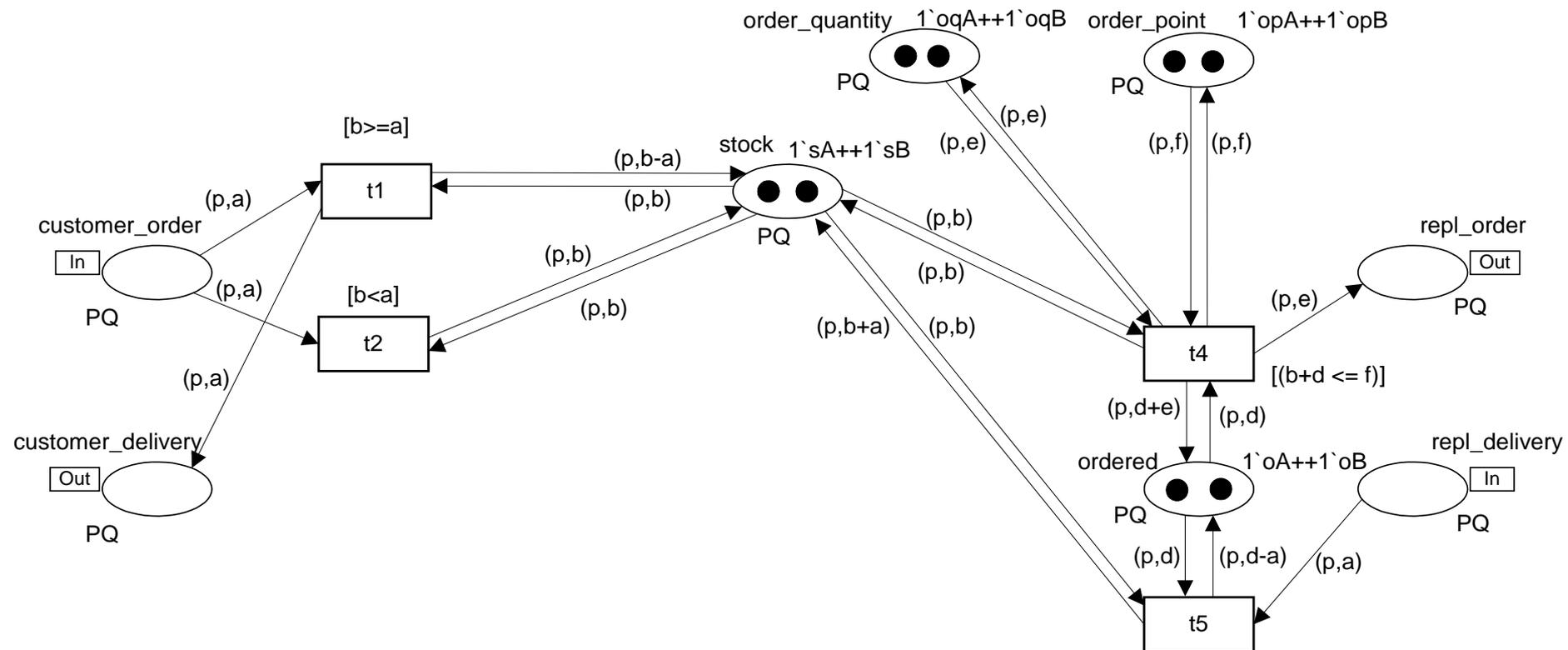


Modeling a supply chain



(For simplicity we do not add external configuration places.
Direction and circular structure may be confusing.)

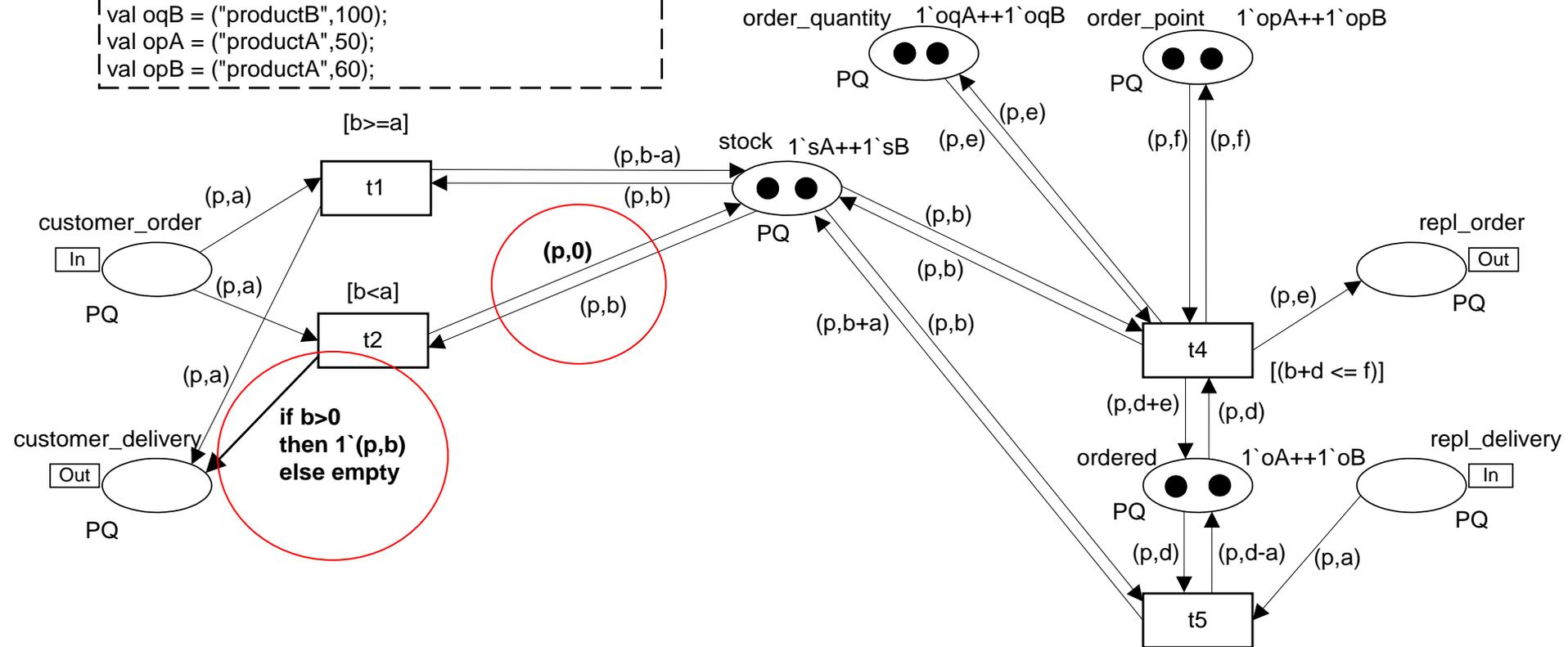
Exercise: Modify to allow for partial shipments



Solution

```

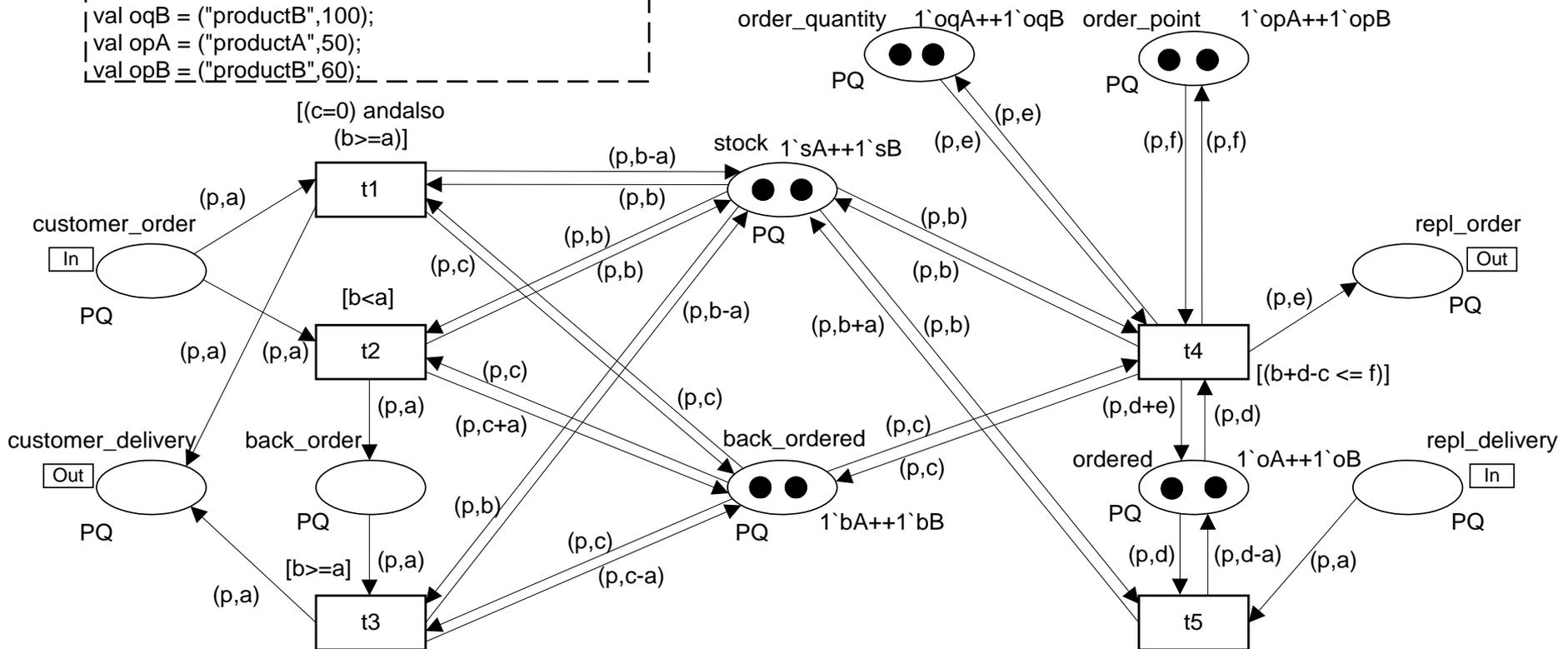
color Product = string;
color Quantity = int;
color PQ = product Product * Quantity;
var p: Product;
var a,b,c,d,e,f: Quantity
val sA = ("productA",0);
val sB = ("productB",0);
val oA = ("productA",0);
val oB = ("productB",0);
val oqA = ("productA",150);
val oqB = ("productB",100);
val opA = ("productA",50);
val opB = ("productA",60);
    
```



Exercise: Modify to allow for partial shipments

```

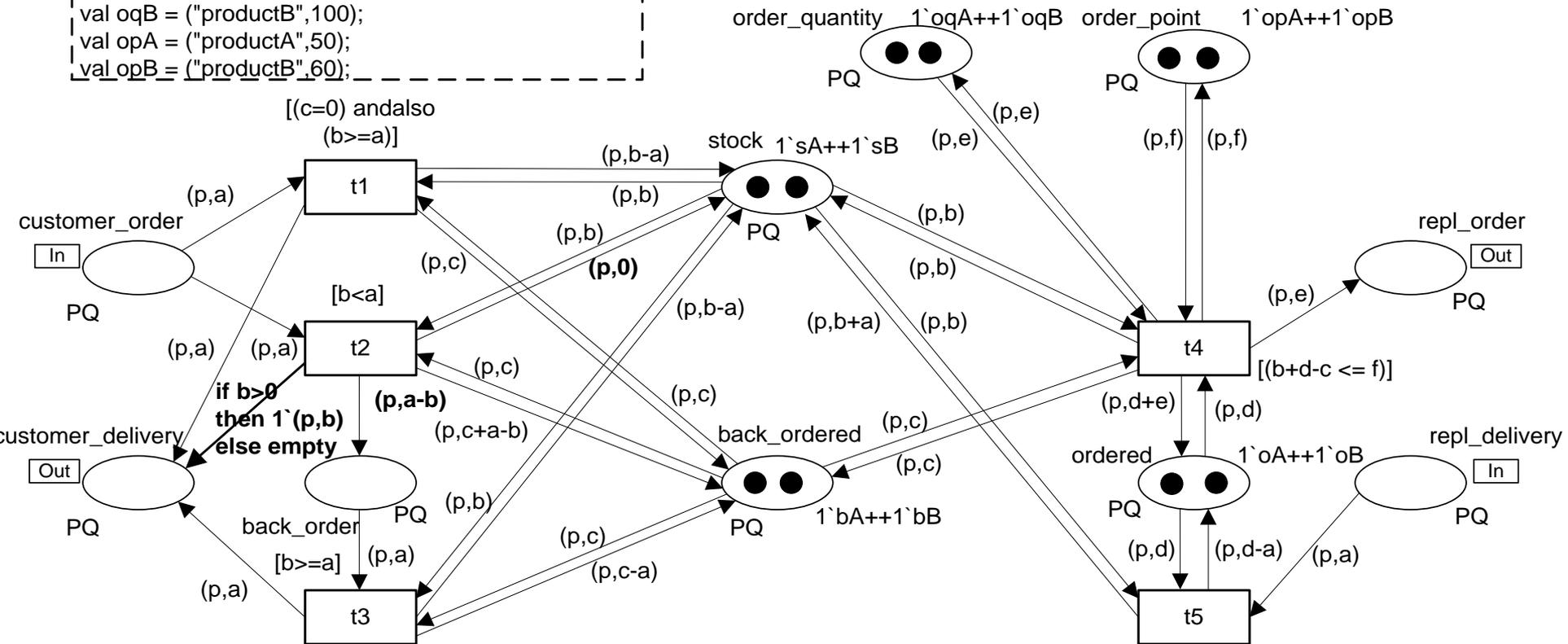
color Product = string;
color Quantity = int;
color PQ = product Product * Quantity;
var p: Product;
var a,b,c,d,e,f: Quantity
val sA = ("productA",0);
val sB = ("productB",0);
val oA = ("productA",0);
val oB = ("productB",0);
val bA = ("productA",0);
val bB = ("productB",0);
val oqA = ("productA",150);
val oqB = ("productB",100);
val opA = ("productA",50);
val opB = ("productB",60);
    
```



Solution

```

color Product = string;
color Quantity = int;
color PQ = product Product * Quantity;
var p: Product;
var a,b,c,d,e,f: Quantity
val sA = ("productA",0);
val sB = ("productB",0);
val oA = ("productA",0);
val oB = ("productB",0);
val bA = ("productA",0);
val bB = ("productB",0);
val oqA = ("productA",150);
val oqB = ("productB",100);
val opA = ("productA",50);
val opB = ("productB",60);
    
```

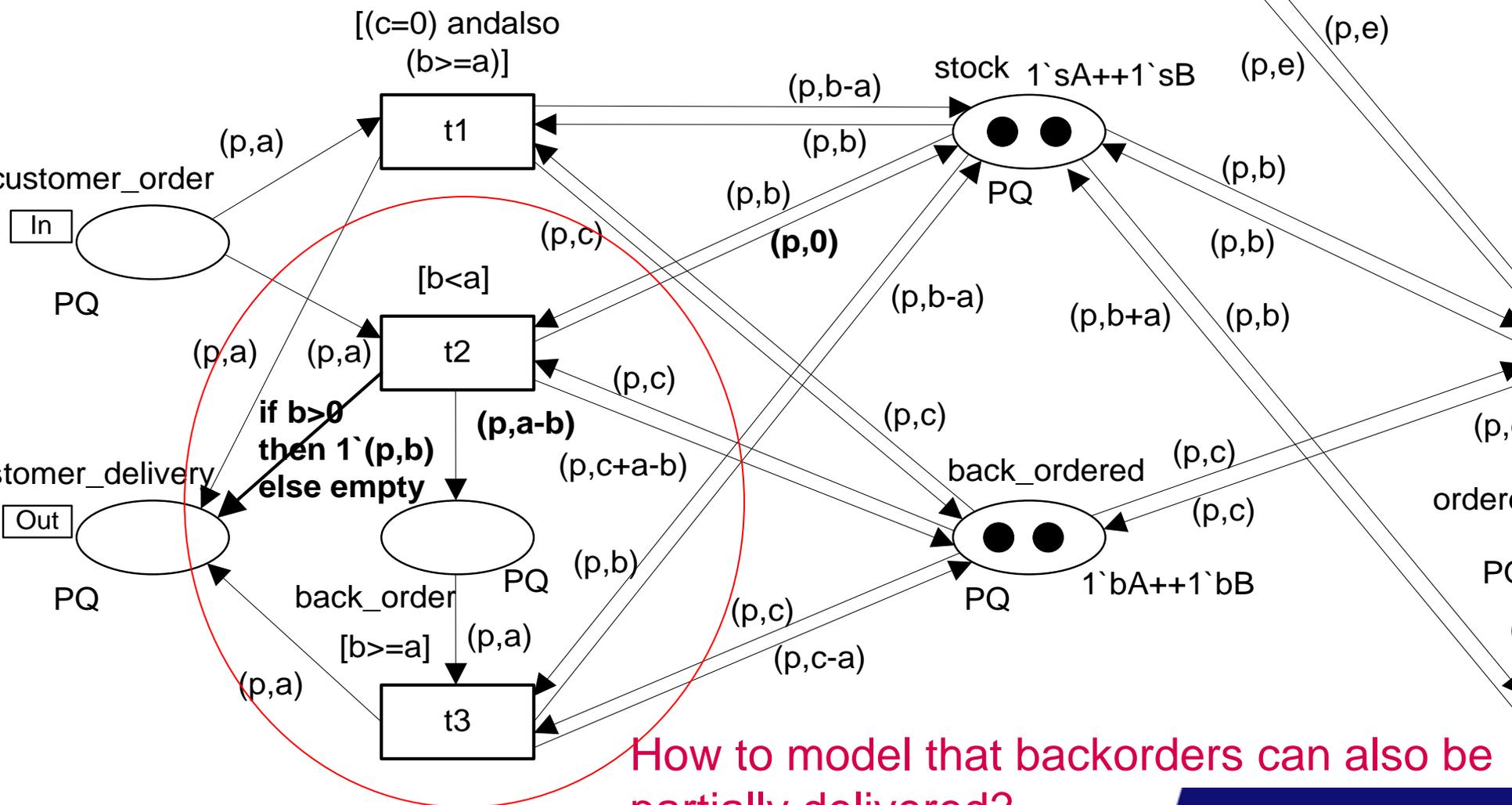


```

val bB = ("productB",0);
val bA = ("productA",0);
val bB = ("productB",0);
val oqA = ("productA",150);
val oqB = ("productB",100);
val opA = ("productA",50);
val opB = ("productB",60);

```

order_quantity 1`oqA++1`oqB order

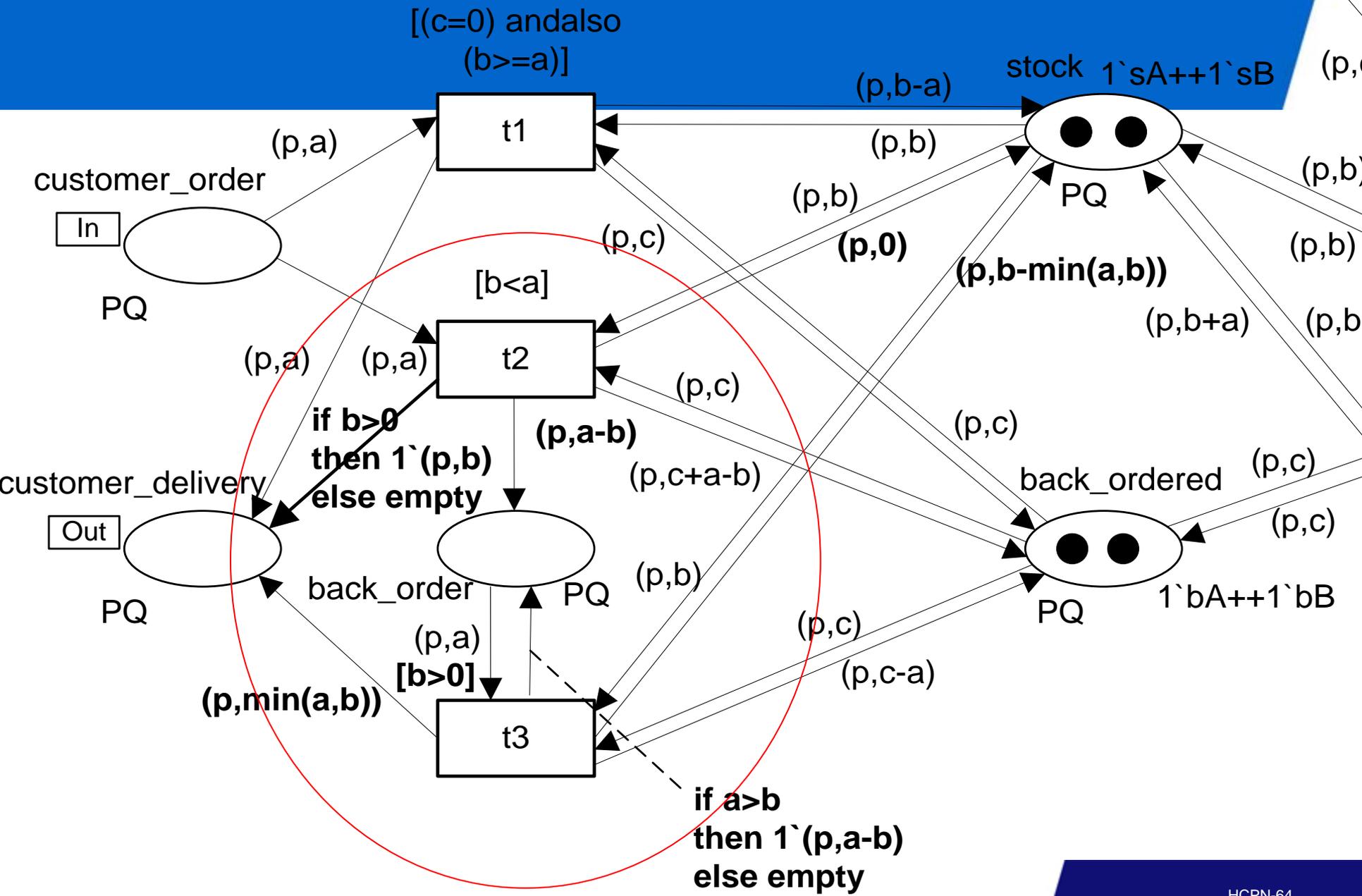


How to model that backorders can also be partially delivered?


```

var a,b,c,a,b,n Quantity
fun min(x:int,y:int) = if x<y the x else y;

```



After studying Chapter 6 one should be able to:

- Flatten a hierarchical CPN model.
- Design a hierarchical CPN model for scratch.
- Modify a hierarchical CPN model.

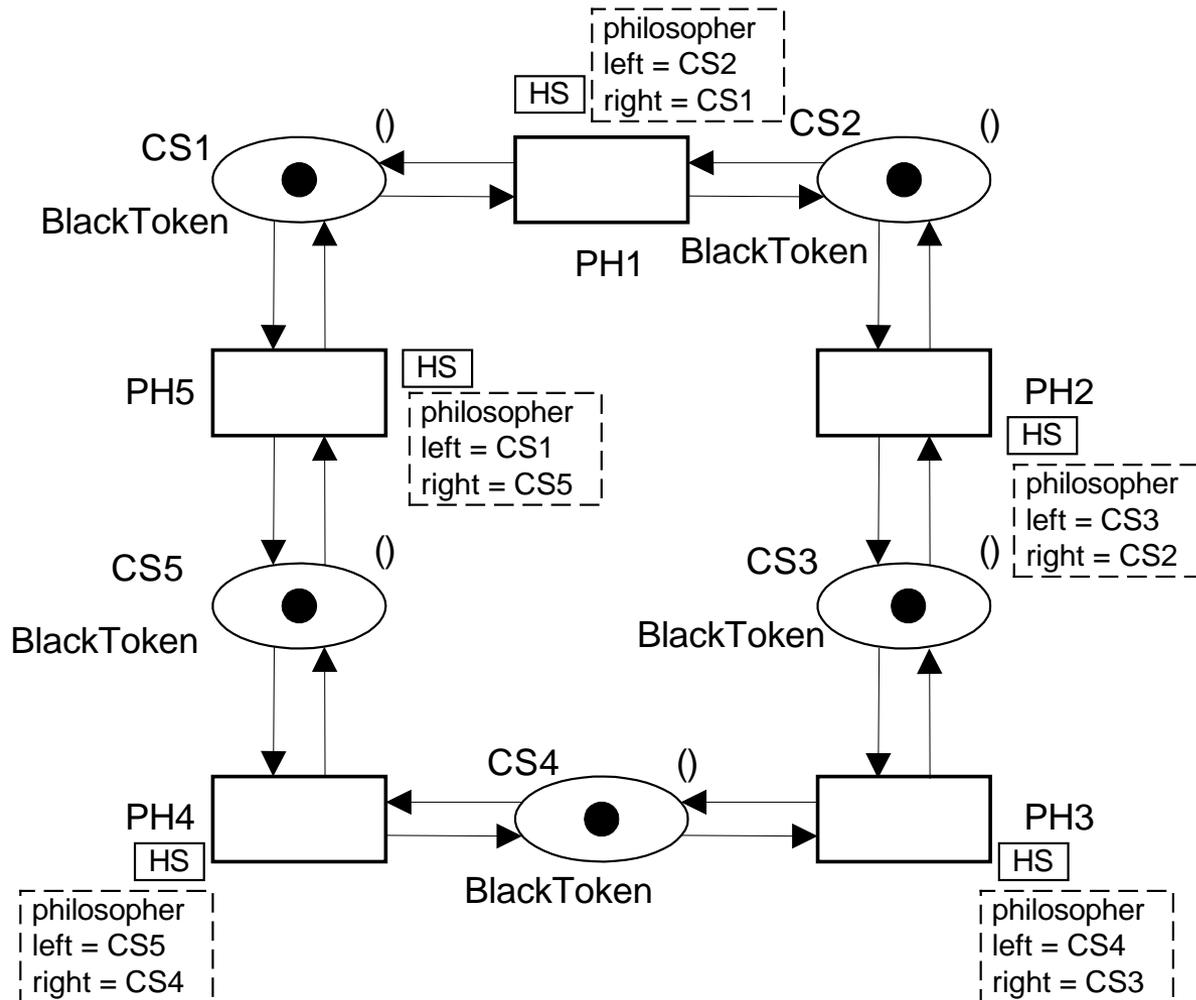


Exercise: Five Chinese philosophers

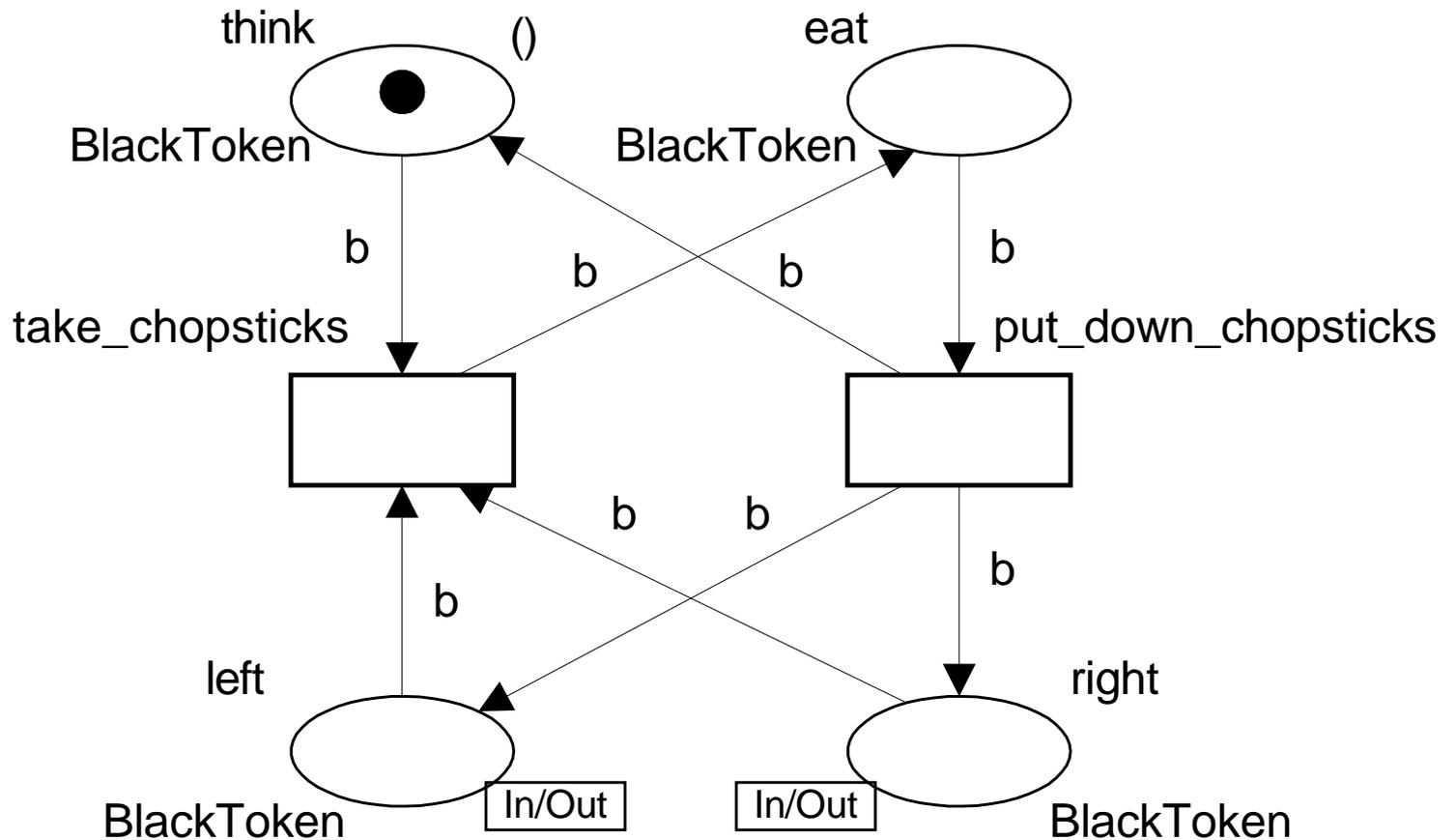
- **Make a hierarchical CPN model of five Chinese philosophers alternating between states thinking and eating. To eat two chopsticks are needed. In total there are five chopsticks. The philosophers are sitting in a circle, and need to complete for chopsticks with their direct neighbors (left and right). Assume that both chopsticks need to be taken at the same time.
Model this using a hierarchical CPN model. Make sure to model the behavior of a philosopher only once and just use the color set BlackToken of type unit.**
- **Change the model such that philosophers can take one chopstick at a time but avoid deadlocks and a fixed ordering of philosophers.**
- **Flatten the hierarchical CPN model.**

Solution: Top-level page

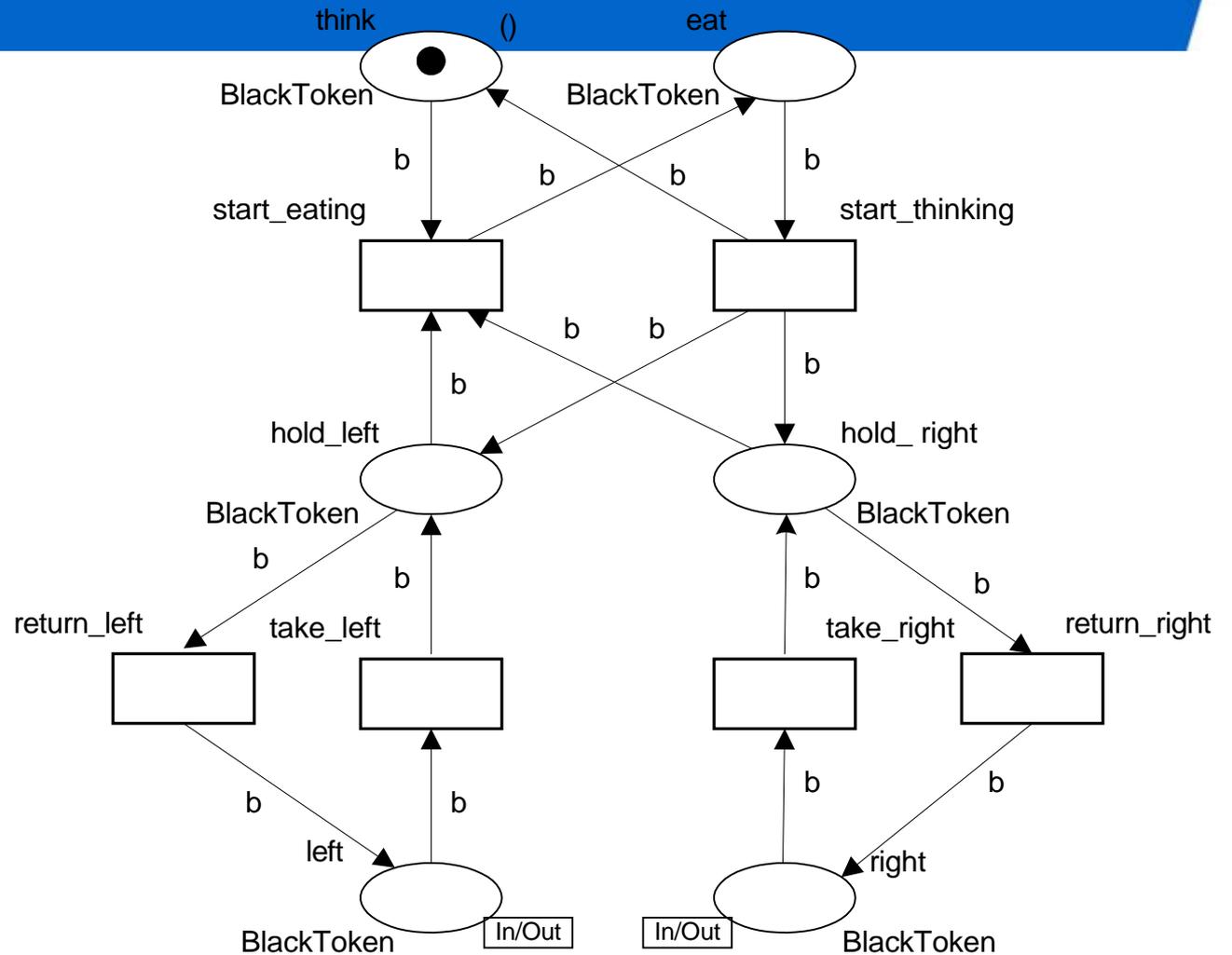
```
color BlackToken = unit;  
var b:BackToken
```



Page philosopher



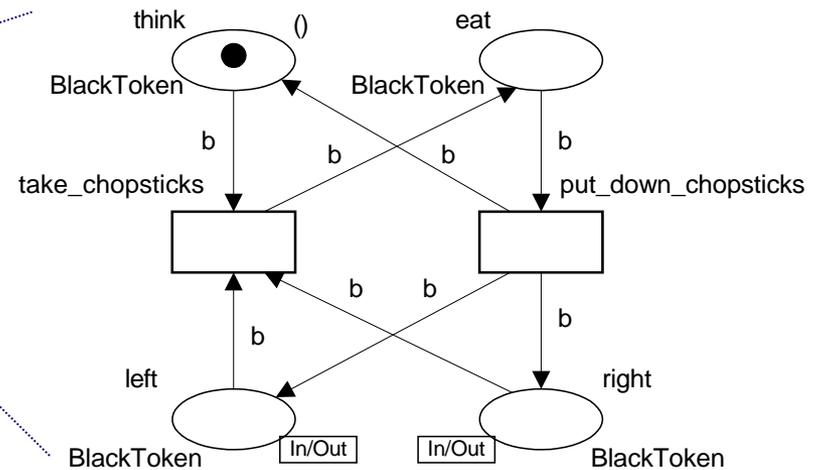
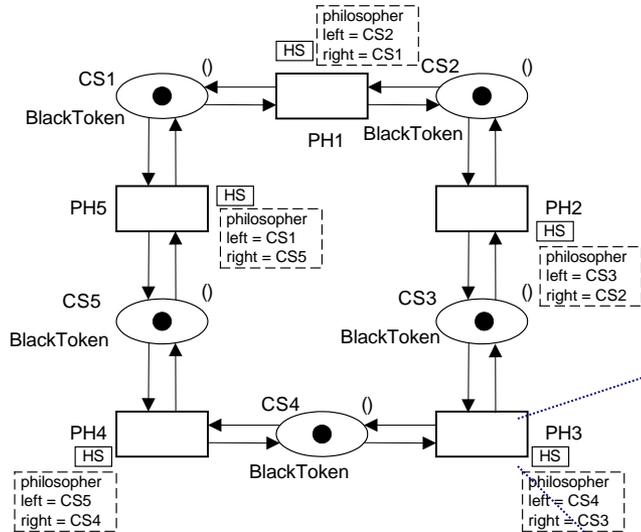
Alternative page



Flat model is obtained by replacing substitution transitions by subpages

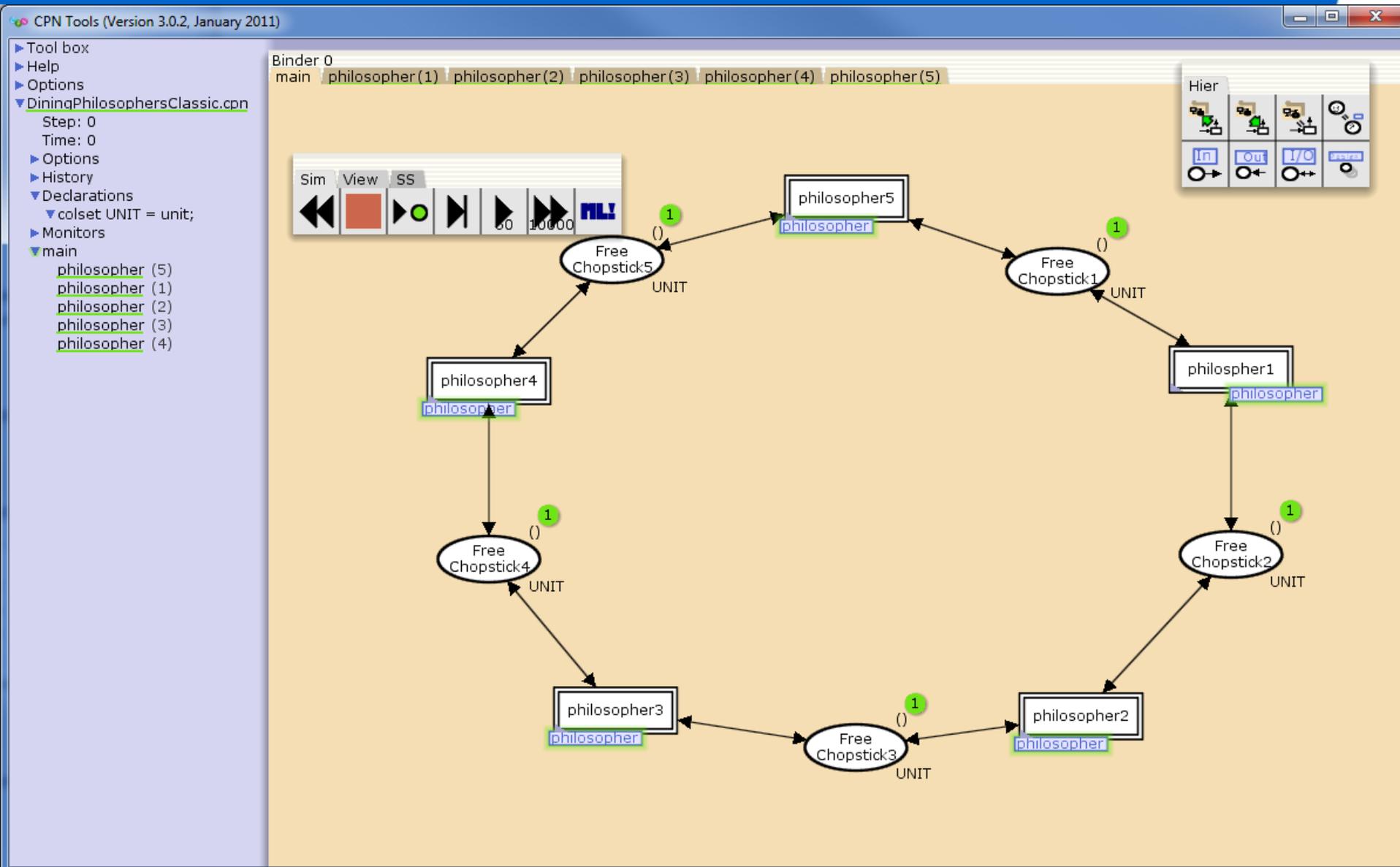
```

color BlackToken = unit;
var b:BlackToken
    
```



Repeat 5 times...

In CPN Tools



- ▶ Tool box
- ▶ Help
- ▶ Options
- ▶ DiningPhilosophersClassic.cpn
 - Step: 0
 - Time: 0
 - ▶ Options
 - ▶ History
 - ▶ Declarations
 - colset UNIT = unit;
 - ▶ Monitors
 - ▶ main
 - philosopher (5)
 - philosopher (1)
 - philosopher (2)
 - philosopher (3)
 - philosopher (4)

Binder 0
main philosopher (1) philosopher (2) philosopher (3) philosopher (4) philosopher (5)

Sim View SS



Hier

