Integer Programming Basic Concepts

Outline

- What is it?
- Problem formulations.
- Binary variables.
- Modeling nonlinear functions

What is it? [1]

- Mathematical programming is the branch of mathematics concerned with the theory and methods for solving problems on finding the extrema of functions on sets defined by linear and non-linear constraints (equalities and inequalities) in a finite-dimensional vector space.
- Mathematical programming is a branch of operations research, which comprises a wide class of control problems

Mathematical programming



What is it? [1]

- The problems of mathematical programming find applications in various areas of human activity where it is necessary to choose one of the possible ways of action, e.g. in solving numerous problems of control and planning of production processes as well as in problems of design and long-term planning.
- The term "mathematical programming" is connected with the fact that the goal of solving various problems is choosing programs of action.

Mathematical programming



What is it?

Linear Program (LP) max cx st. $Ax \le b$ $x \ge 0$







Max $x_1 + x_2$

s.t.

x1

 $x_1 + 2x_2 \le 2$

 $2x_1 + x_2 \le 2$

 $x_1 \ge 0, x_2 \ge 0$

Solution: (2/3, 2/3)

What is it? Integer Program (IP)

max cx st. $Ax \le b$ $x \in Z_{+}$ (positive integer)



Max $x_1 + x_2$

 $x_1 + 2x_2 \le 2$

 $2x_1 + x_2 \le 2$

 $x_1 \ge 0, x_2 \ge 0$ integer

Solution: (1, 0) and (0,1)

What is it? Integer Program (IP)



Max $5x_1 + 8x_2$

s.t.

 $x_1 + x_2 \le 6$

 $5x_1 + 9x_2 \le 45$

 $x_1 \ge 0, x_2 \ge 0$ integer

What is it? Mixed-Integer Program (MIP) max cx + hyst. $Ax + Gy \le b$ $x \ge 0, y \in Z_{\perp}$ (positive integer)



Problem Formulations

Knapsack Problem



Problem Formulations

Traveling Salesman Problem (TSP)

Problem Formulations

Lot-sizing problem



Logical constraints [2]

Constraints feasibility: When is the constraint satisfied?

 $f(x1,x2,...,xn) \le b$ $f(x1,x2,...,xn) - By \le b , y \in \{0,1\}$ $y=0 \Longrightarrow f(x1,x2,...,xn) \le b$ $y=1 \Longrightarrow f(x1,x2,...,xn) \le b + B$

Logical constraints [2]

• Alternative constraints

 $f1(x1,x2,...,xn) \le b1$

 $f2(x1,x2,...,xn) \le b2$

 $f(x1,x2,...,xn) - By1 \le b1$ $f(x1,x2,...,xn) - By2 \le b2$ $y1+y2 \le 1, y \in \{0,1\}$ $f(x1,x2,...,xn) - By \le b1$ $f(x1,x2,...,xn) - B(1-y) \le b2$ $y \in \{0,1\}$

Logical constraints [2]

• Conditional constraints:

 $f1(x1,x2,...,xn) > b1 \implies f2(x1,x2,...,xn) \le b2$

 \Leftrightarrow

 $f1(x1,x2,...,xn) \le b1 \lor f2(x1,x2,...,xn) \le b2$

Logical constraints [2]

Conditional constraints: k-fold alternatives:

We must satisfy at least k constraints from

$$\begin{split} fj(x1,x2,...,xn) &\leq bj \quad \text{for } j=1,...,p \\ fj(x1,x2,...,xn) - Bj(1-yj) &\leq bj \quad \text{for } j=1,...,p \\ y1+...+yj &\geq k \\ yj &\in \{0,1\}, \ j=1,...,p \end{split}$$

Logical constraints [2]

Compound alternatives

Region 1: $f1(x1,x2) - B1y1 \le b1$ $f_2(x_1,x_2) - B_2y_1 \le b_2$ $f_3(x_1,x_2) - B_3y_2 \le b_3$ Region 2: $f4(x1,x2) - B4y2 \le b4$ $f5(x1,x2) - B5y3 \le b5$ Region 3: $f6(x1,x2) - B6y3 \le b6$ $f7(x1,x2) - B7y3 \le b7$ $y1 + y2 + y3 \le 2$, $x1 \ge 0, x2 \ge 0,$ y1, y2, y3 binary.





- $\mathbf{x} = \delta \mathbf{1} + \delta \mathbf{2} + \delta \mathbf{3},$
- $0 \le \delta 1 \le 4$
- $0 \le \delta 2 \le 6$
- $0 \le \delta 3 \le 5$
- $Cost = 5\delta 1 + \delta 2 + 3\delta 3.$
- $x = 2 \Longrightarrow \delta 1 = \delta 3 = 0$
 - δ2 = 2, Cost=2



 $x = \delta 1 + \delta 2 + \delta 3$, $4w1 \le \delta 1 \le 4$, $6w2 \le \delta 2 \le 6w1$, $0 \leq \delta 3 \leq 5w2$. w1, w2 binary $Cost = 5\delta 1 + \delta 2 + 3\delta 3.$ $x = 2 \Longrightarrow \delta 1 = 2, w 1 = 0$ δ3=δ2 = 0, w2=0 Cost=10



- $\mathbf{x} = \delta \mathbf{1} + \delta \mathbf{2} + \delta \mathbf{3},$
- $4w1 \le \delta 1 \le 4$,
- $6w2 \le \delta 2 \le 6w1$,
- $0 \le \delta 3 \le 5w2$,
- w1, w2 binary

- w1 = 0 \Rightarrow w2 = 0 with 0 $\leq \delta 1 \leq 4$, $\delta 2 = 0$, and $\delta 3 = 0$.
- w1 = 1 e w2 = 0 with $\delta 1 = 4$, $0 \le \delta 2 \le 6$, and $\delta 3 = 0$.
- w1 = 1 e w2 = 1 with $\delta 1$ = 4, $\delta 2$ = 6, and $0 \le \delta 3 \le 5$.
- w1 = 0, w2 = 0 \Rightarrow 0 \leq x \leq 4 since δ 2 = δ 3 = 0.
- w1 = 1, w2 = 0 \Rightarrow 4 \leq x \leq 10 since δ 1 = 4 and δ 3 = 0.
- w1 = 1, w2 = 1 \Rightarrow 10 \leq x \leq 15 since δ 1 = 4 and δ 2 = 6.

 $L_j w_j \le \delta_j \le L j w_{j-1}$,

 $Cost = \overline{\delta}1 + 3\overline{\delta}2 + 6\overline{\delta}3$ s.t. $0 \le \delta 1 \le 4$ $0 \le \delta 2 \le 6$, $0 \le \delta 3 \le 5$.





Referências

[1] Enciclopedia of Mathematics,

https://www.encyclopediaofmath.org/index.php/Mathematical_programming.

[2] MIT - Integer Programming,

http://web.mit.edu/15.053/www/AMP-Chapter-09.pdf.