

Mathematical Preliminaries

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- Sets
- Functions
- Relations
- Graphs
- Proof Techniques

SETS

A set is a collection of elements

$$A = \{1, 2, 3\}$$

$$B = \{\textit{train}, \textit{bus}, \textit{bicycle}, \textit{airplane}\}$$

We write

$$1 \in A$$

$$\textit{ship} \notin B$$

Set Representations

$$C = \{ a, b, c, d, e, f, g, h, i, j, k \}$$

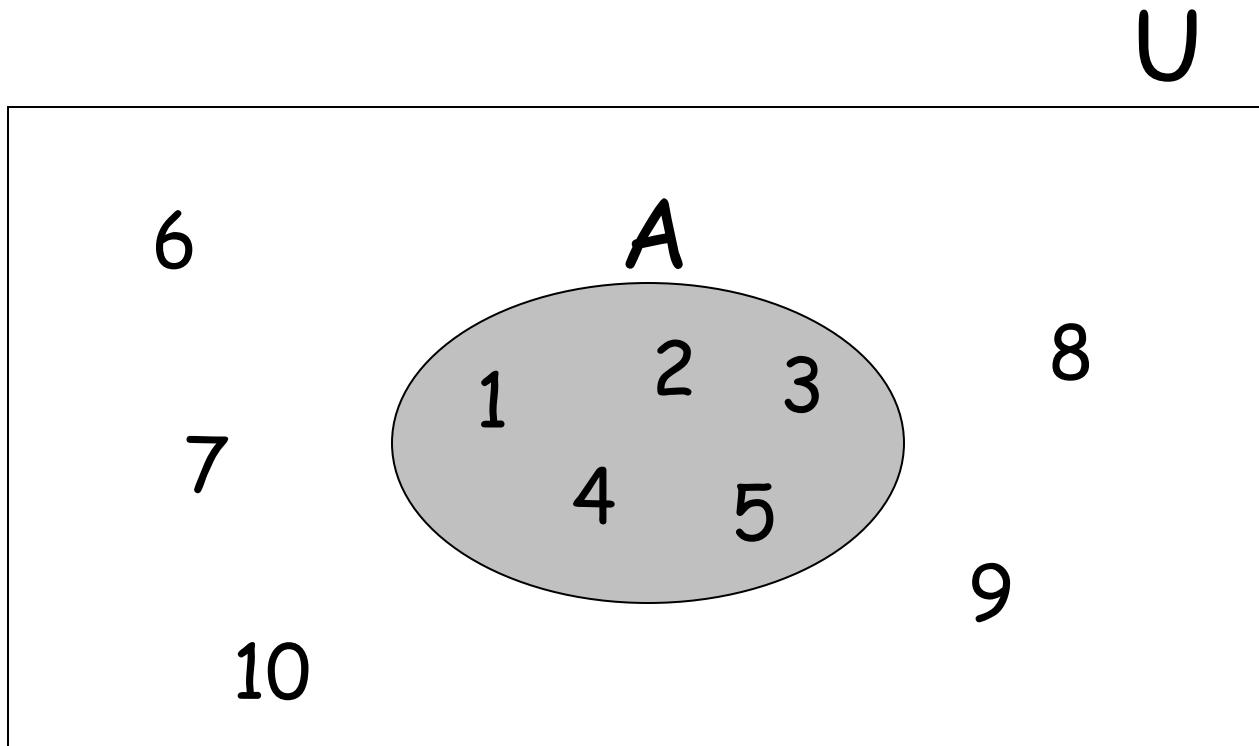
$$C = \{ a, b, \dots, k \} \longrightarrow \text{finite set}$$

$$S = \{ 2, 4, 6, \dots \} \longrightarrow \text{infinite set}$$

$$S = \{ j : j > 0, \text{ and } j = 2k \text{ for some } k > 0 \}$$

$$S = \{ j : j \text{ is nonnegative and even} \}$$

$$A = \{ 1, 2, 3, 4, 5 \}$$



Universal Set: all possible elements

$$U = \{ 1, \dots, 10 \}$$

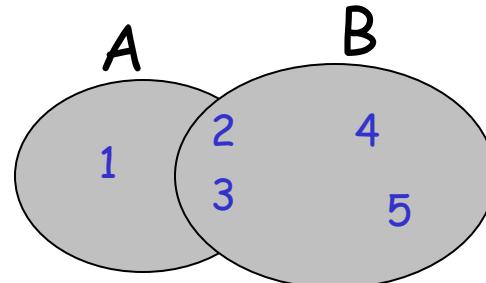
Set Operations

$$A = \{ 1, 2, 3 \}$$

$$B = \{ 2, 3, 4, 5 \}$$

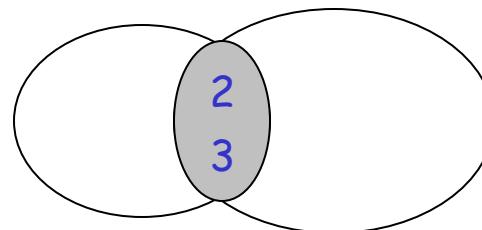
- Union

$$A \cup B = \{ 1, 2, 3, 4, 5 \}$$



- Intersection

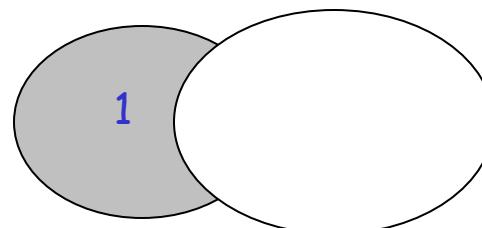
$$A \cap B = \{ 2, 3 \}$$



- Difference

$$A - B = \{ 1 \}$$

$$B - A = \{ 4, 5 \}$$

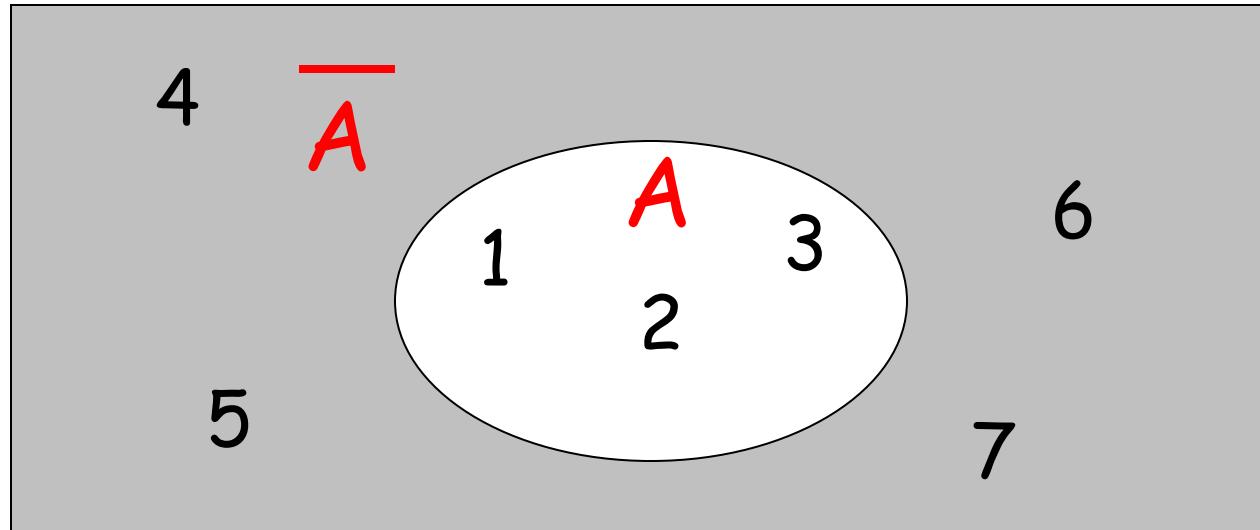


Venn diagrams

- Complement

Universal set = {1, ..., 7}

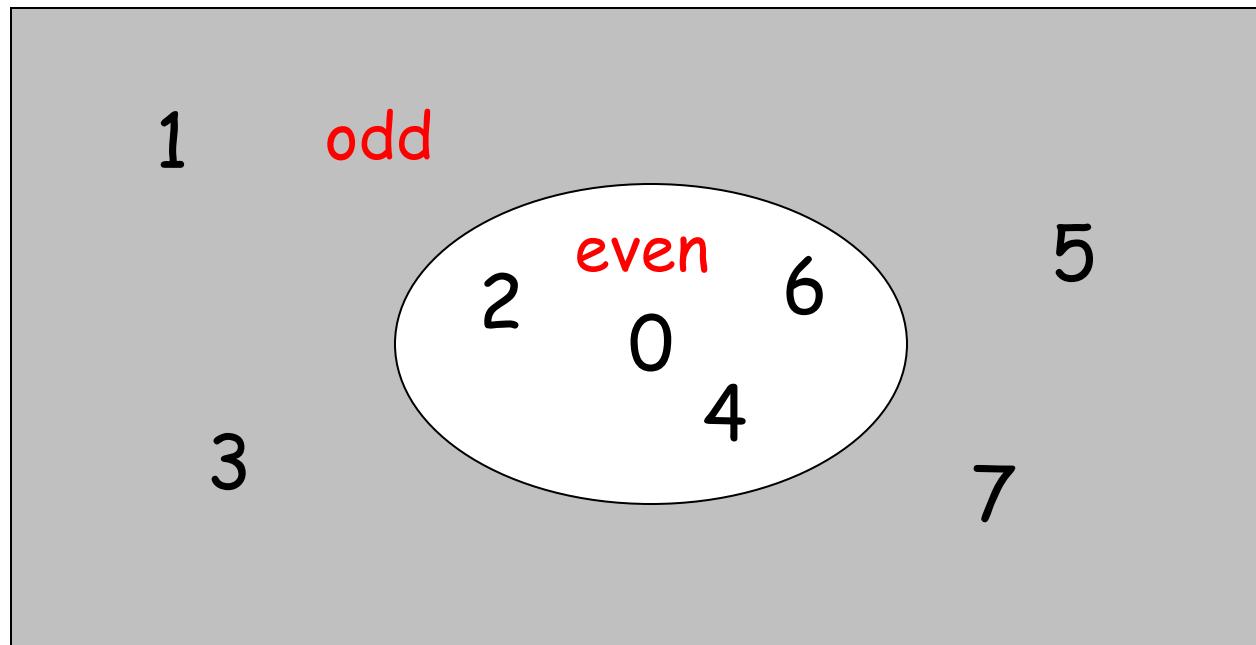
$$A = \{1, 2, 3\} \rightarrow \overline{A} = \{4, 5, 6, 7\}$$



$$\overline{\overline{A}} = A$$

$$\{ \text{even integers} \} = \{ \text{odd integers} \}$$

Integers



DeMorgan's Laws

$$\overline{A \cup B} = \overline{A} \cap \overline{B}$$

$$\overline{A \cap B} = \overline{A} \cup \overline{B}$$

Empty, Null Set: \emptyset

$$\emptyset = \{ \}$$

$$S \cup \emptyset = S$$

$$S \cap \emptyset = \emptyset$$

$\overline{\emptyset}$ = Universal Set

$$S - \emptyset = S$$

$$\emptyset - S = \emptyset$$

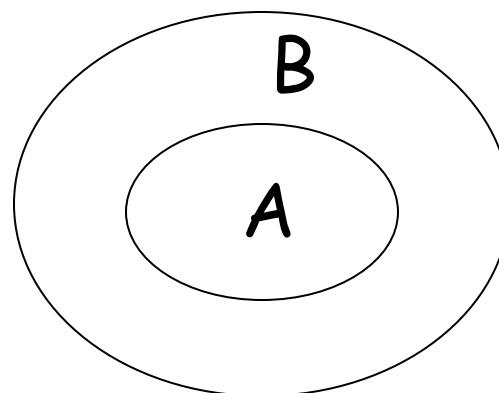
Subset

$$A = \{ 1, 2, 3 \}$$

$$B = \{ 1, 2, 3, 4, 5 \}$$

$$A \subseteq B$$

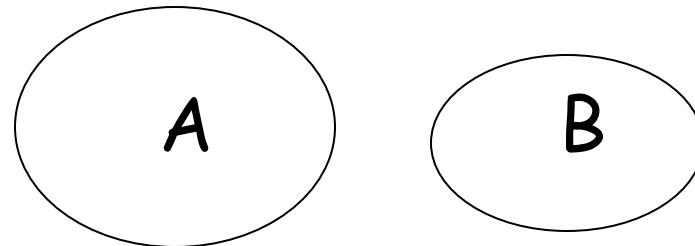
Proper Subset: $A \subset B$



Disjoint Sets

$$A = \{ 1, 2, 3 \} \quad B = \{ 5, 6 \}$$

$$A \cap B = \emptyset$$



Set Cardinality

- For finite sets

$$A = \{ 2, 5, 7 \}$$

$$|A| = 3$$

(set size)

Powersets

A powerset is a set of sets

$$S = \{ a, b, c \}$$

Powerset of S = the set of all the subsets of S

$$2^S = \{ \emptyset, \{a\}, \{b\}, \{c\}, \{a, b\}, \{a, c\}, \{b, c\}, \{a, b, c\} \}$$

Observation: $|2^S| = 2^{|S|}$ ($8 = 2^3$)

Cartesian Product

$$A = \{ 2, 4 \}$$

$$B = \{ 2, 3, 5 \}$$

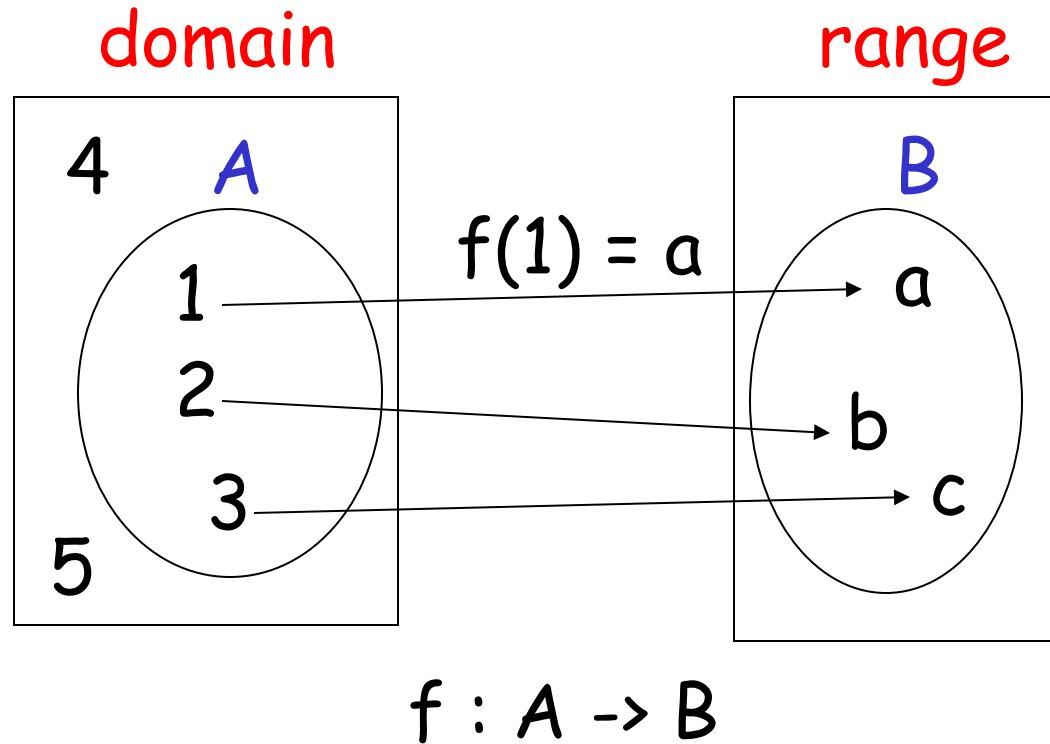
$$A \times B = \{ (2, 2), (2, 3), (2, 5), (4, 2), (4, 3), (4, 5) \}$$

$$|A \times B| = |A| |B|$$

Generalizes to more than two sets

$$A \times B \times \dots \times Z$$

FUNCTIONS



If $A = \text{domain}$

then f is a total function

otherwise f is a partial function

RELATIONS

$R = \{(x_1, y_1), (x_2, y_2), (x_3, y_3), \dots\}$

$x_i R y_i$

e. g. if $R = >$: $2 > 1, 3 > 2, 3 > 1$

Equivalence Relations

- Reflexive: $x R x$
- Symmetric: $x R y \xrightarrow{\hspace{1cm}} y R x$
- Transitive: $x R y$ and $y R z \xrightarrow{\hspace{1cm}} x R z$

Example: $R = '='$

- $x = x$
- $x = y \xrightarrow{\hspace{1cm}} y = x$
- $x = y$ and $y = z \xrightarrow{\hspace{1cm}} x = z$

Equivalence Classes

For equivalence relation R

equivalence class of $x = \{y : x R y\}$

Example:

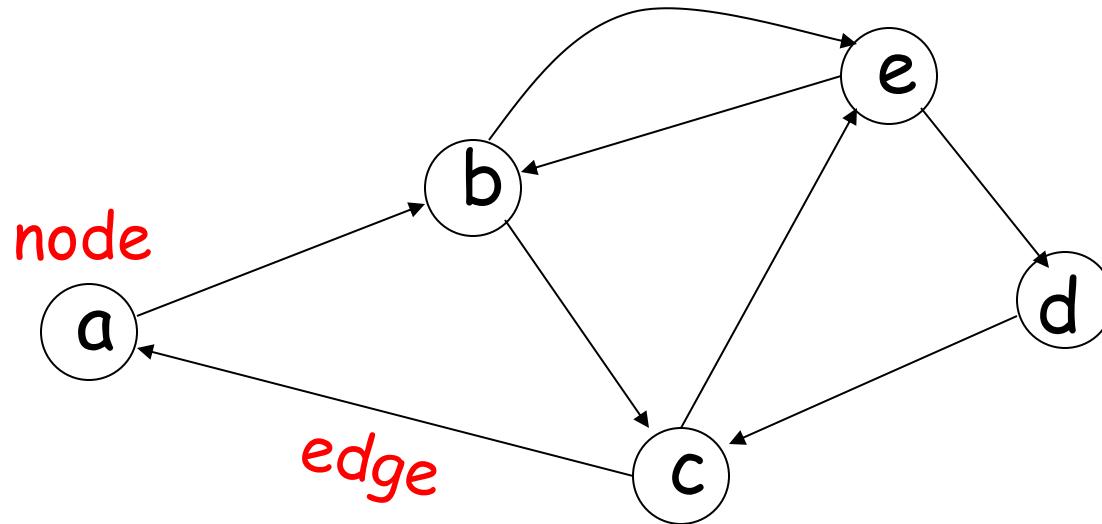
$$R = \{ (1, 1), (2, 2), (1, 2), (2, 1), (3, 3), (4, 4), (3, 4), (4, 3) \}$$

Equivalence class of $1 = \{1, 2\}$

Equivalence class of $3 = \{3, 4\}$

GRAPHS

A directed graph



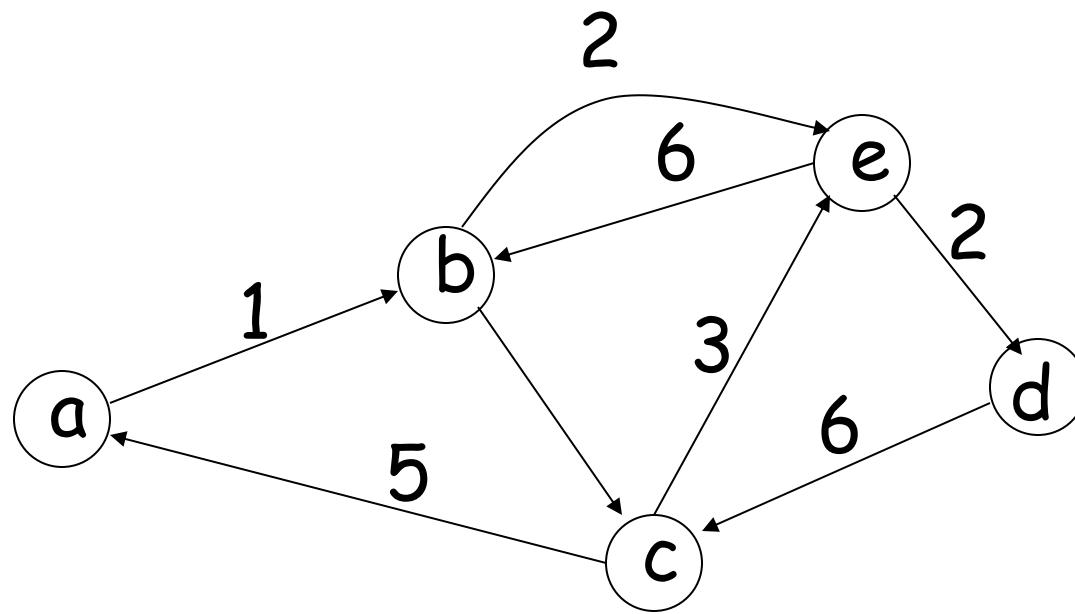
- Nodes (Vertices)

$$V = \{ a, b, c, d, e \}$$

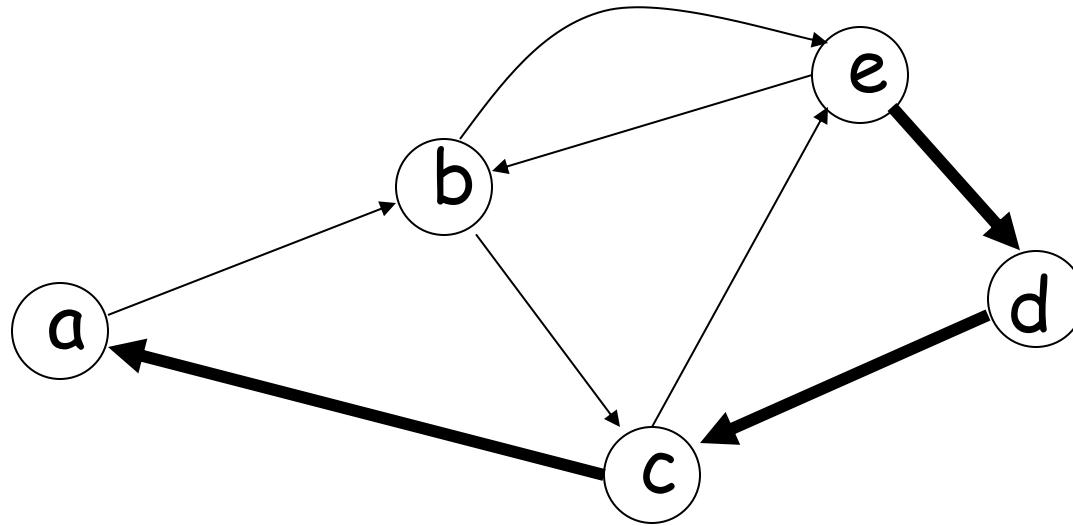
- Edges

$$E = \{ (a,b), (b,c), (b,e), (c,a), (c,e), (d,c), (e,b), (e,d) \}$$

Labeled Graph



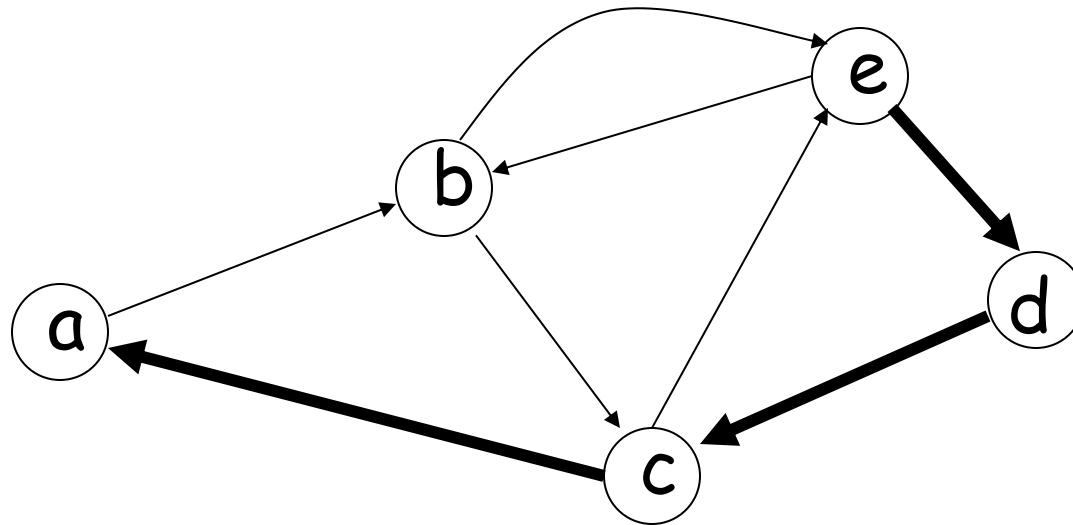
Walk



Walk is a sequence of adjacent edges

$(e, d), (d, c), (c, a)$

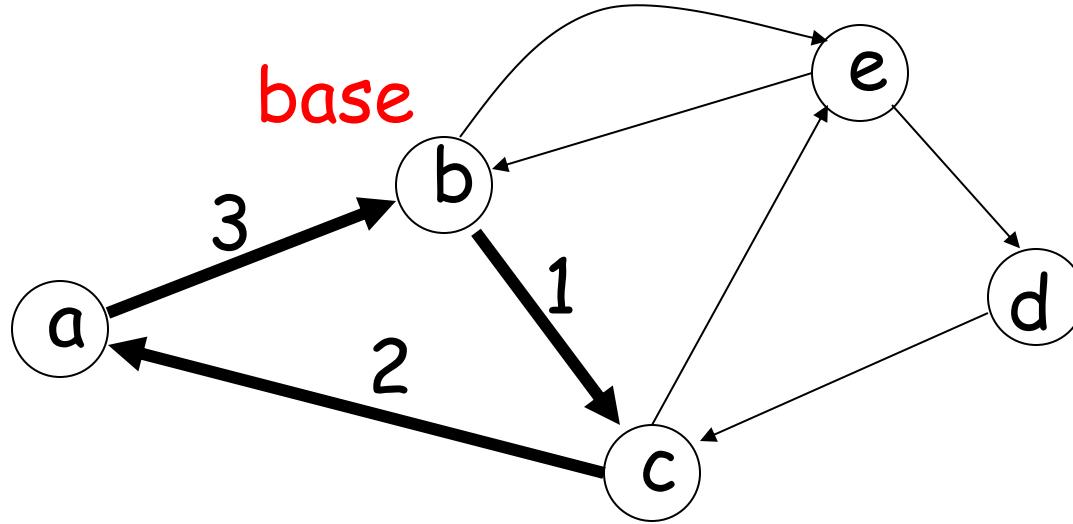
Path



Path is a walk where no edge is repeated

Simple path: no node is repeated

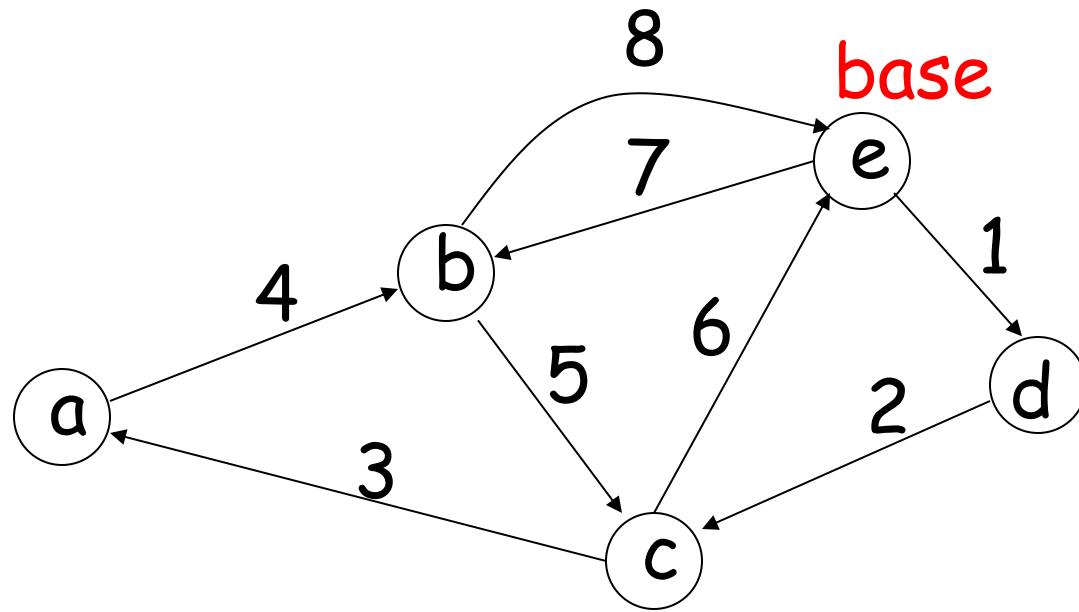
Cycle



Cycle: a walk from a node (base) to itself

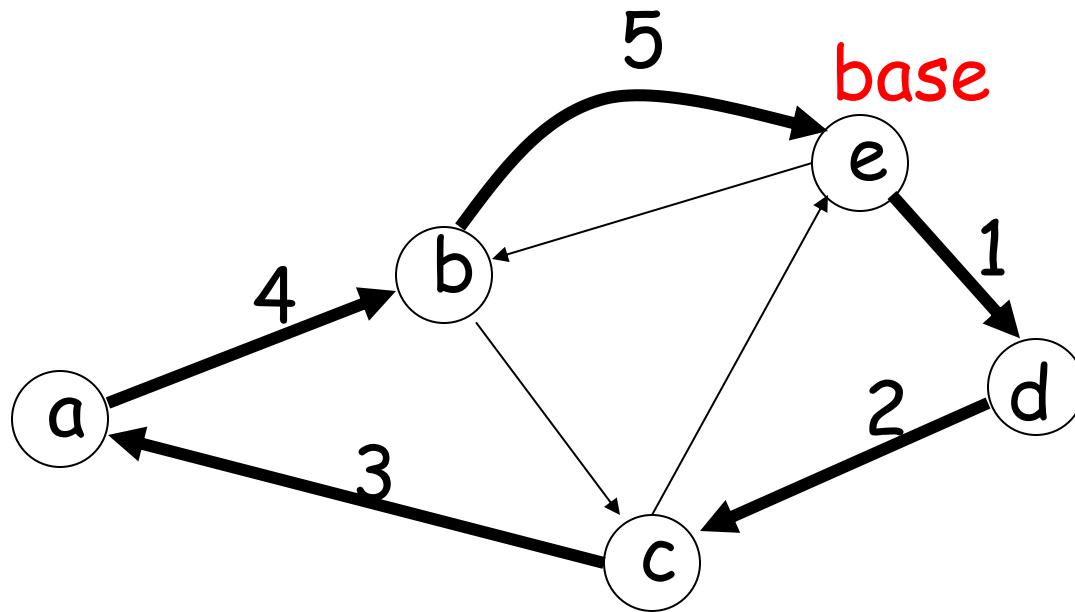
Simple cycle: only the base node is repeated

Euler Tour



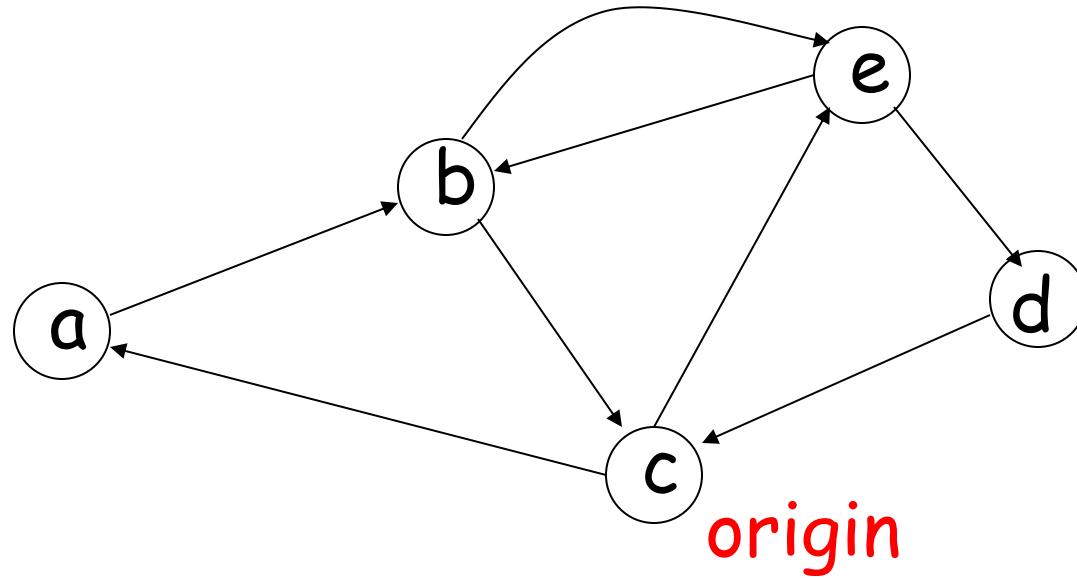
A cycle that contains each edge once

Hamiltonian Cycle

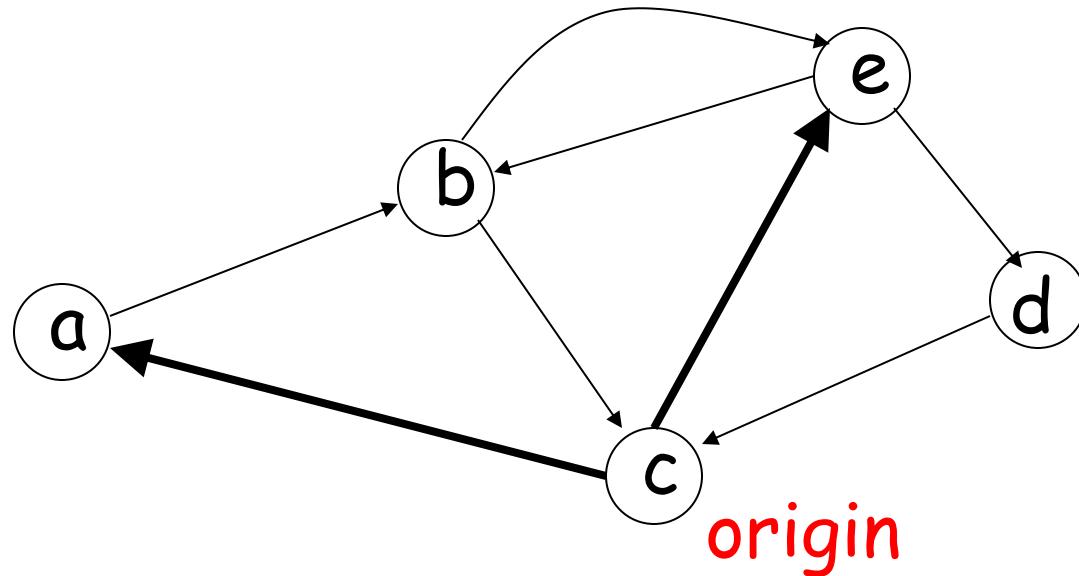


A simple cycle that contains all nodes

Finding All Simple Paths



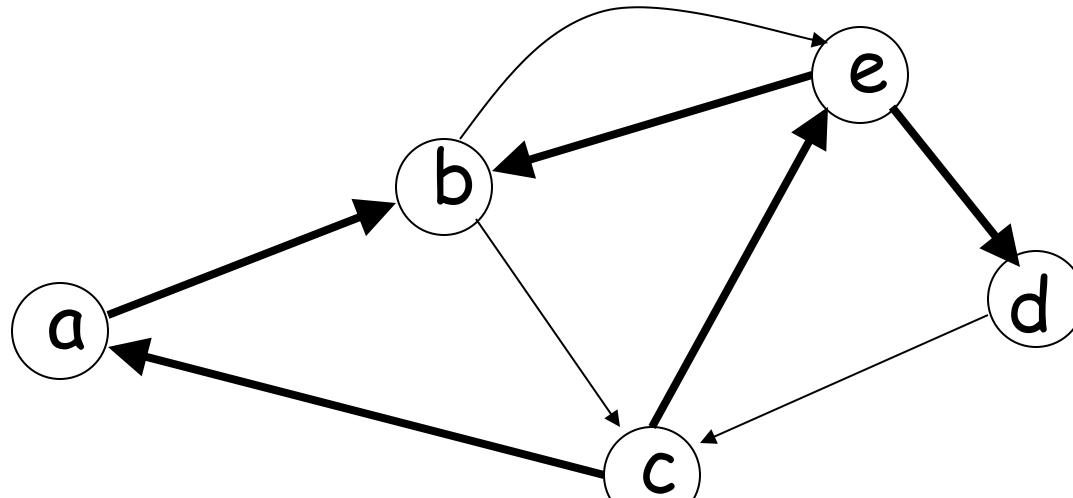
Step 1



(c, a)

(c, e)

Step 2



(c, a)

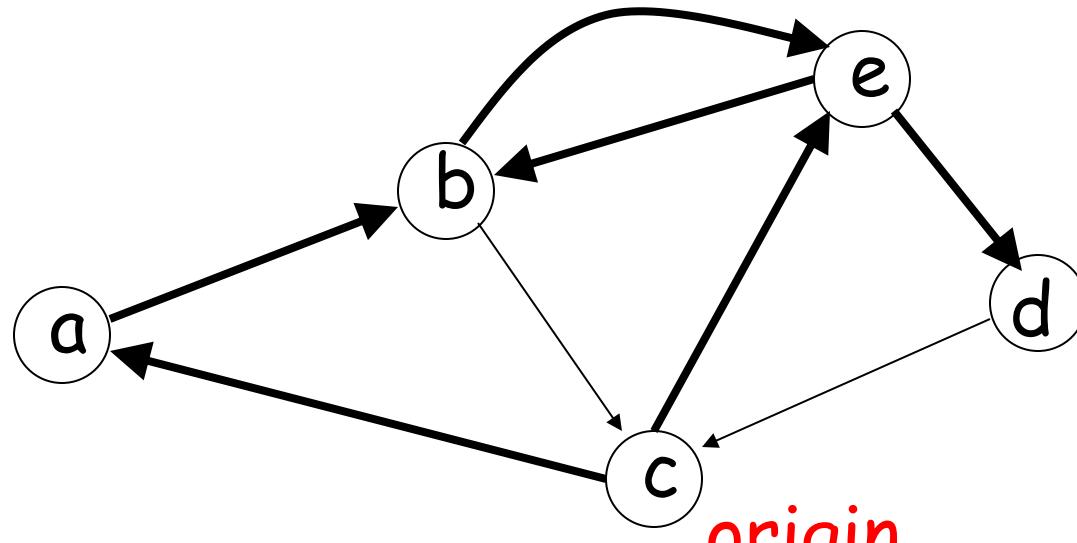
(c, a), (a, b)

(c, e)

(c, e), (e, b)

(c, e), (e, d)

Step 3



(c, a)

(c, a), (a, b)

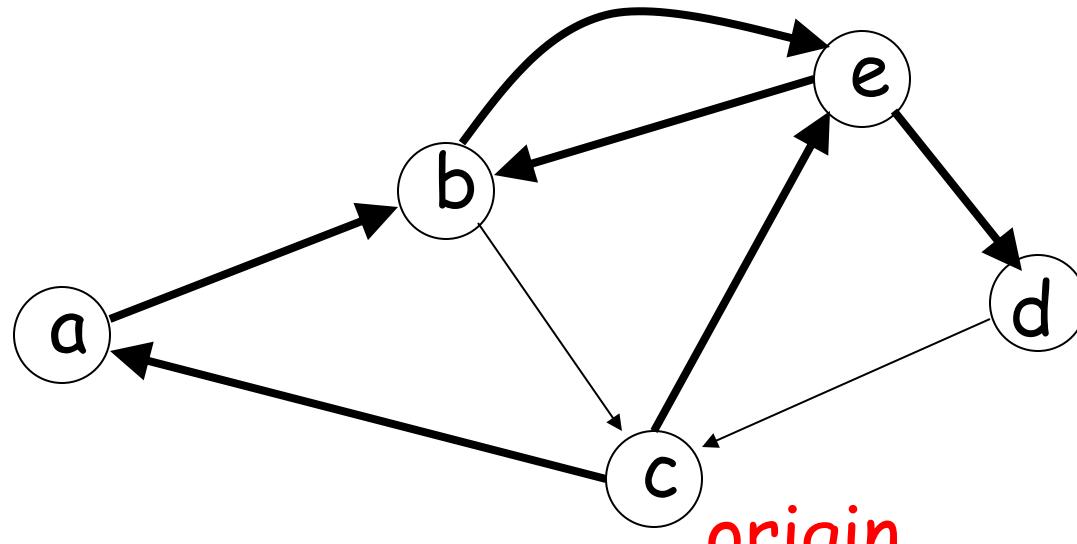
(c, a), (a, b), (b, e)

(c, e)

(c, e), (e, b)

(c, e), (e, d)

Step 4



(c, a)

(c, a), (a, b)

(c, a), (a, b), (b, e)

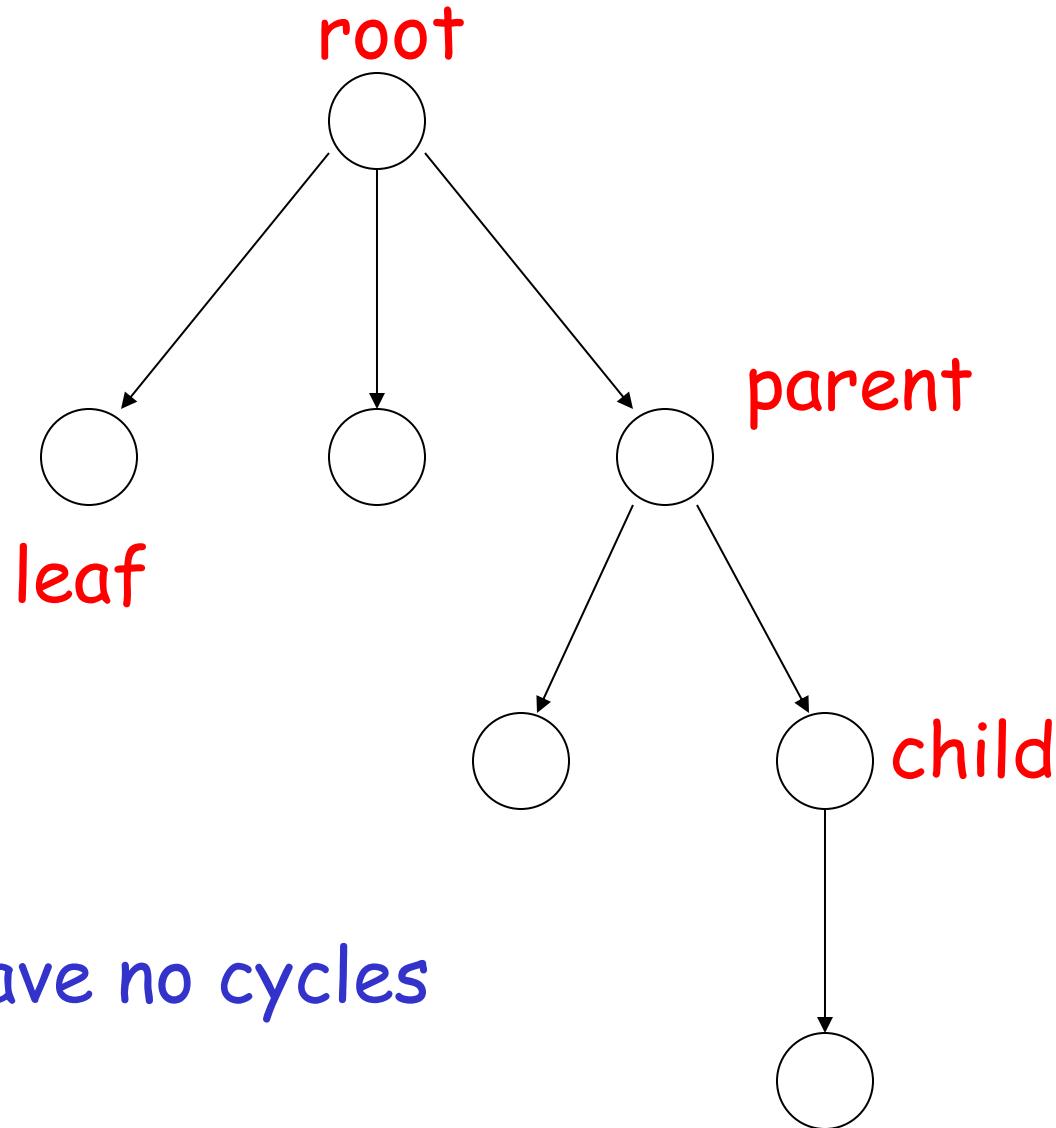
(c, a), (a, b), (b, e), (e, d)

(c, e)

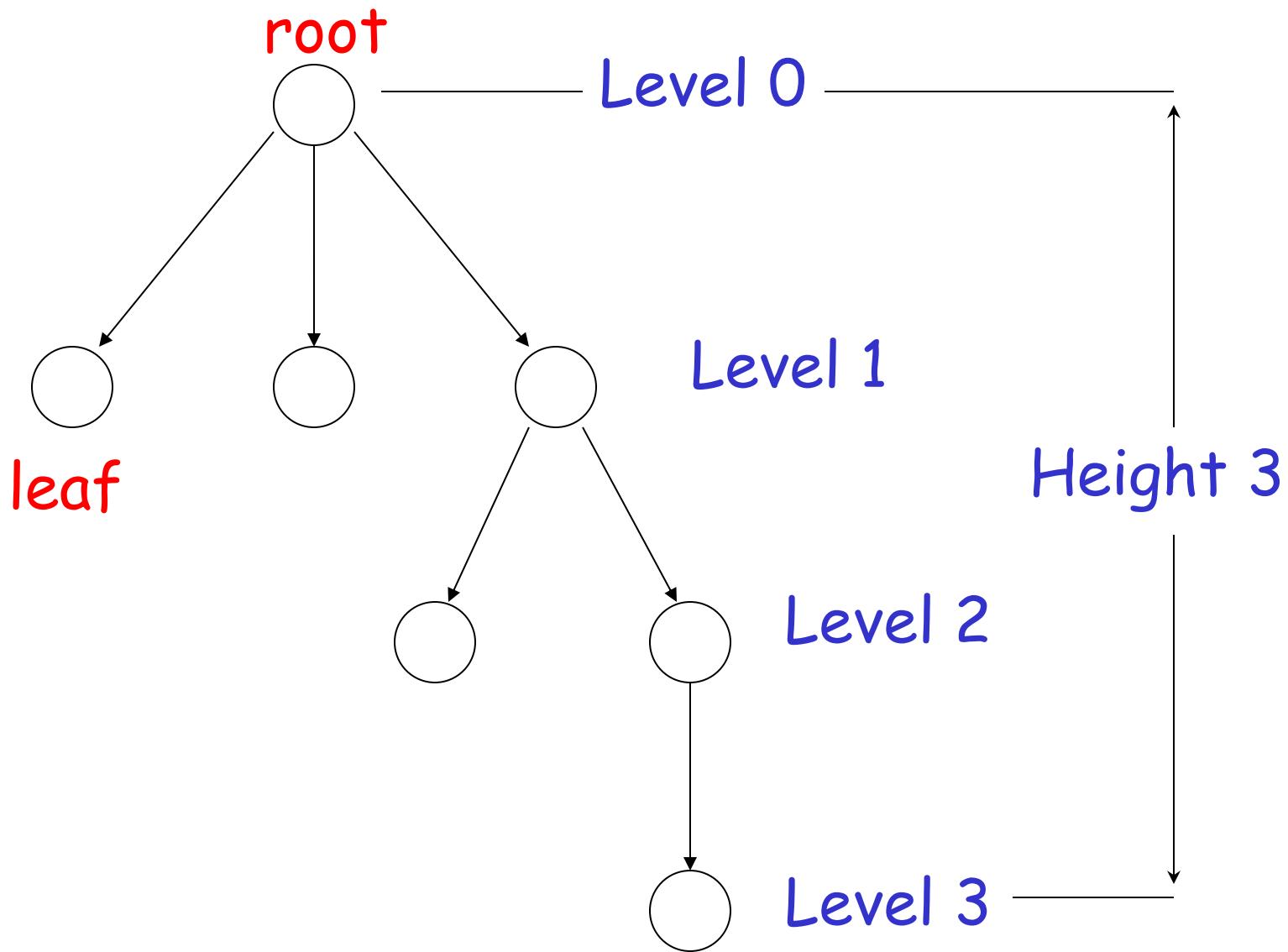
(c, e), (e, b)

(c, e), (e, d)

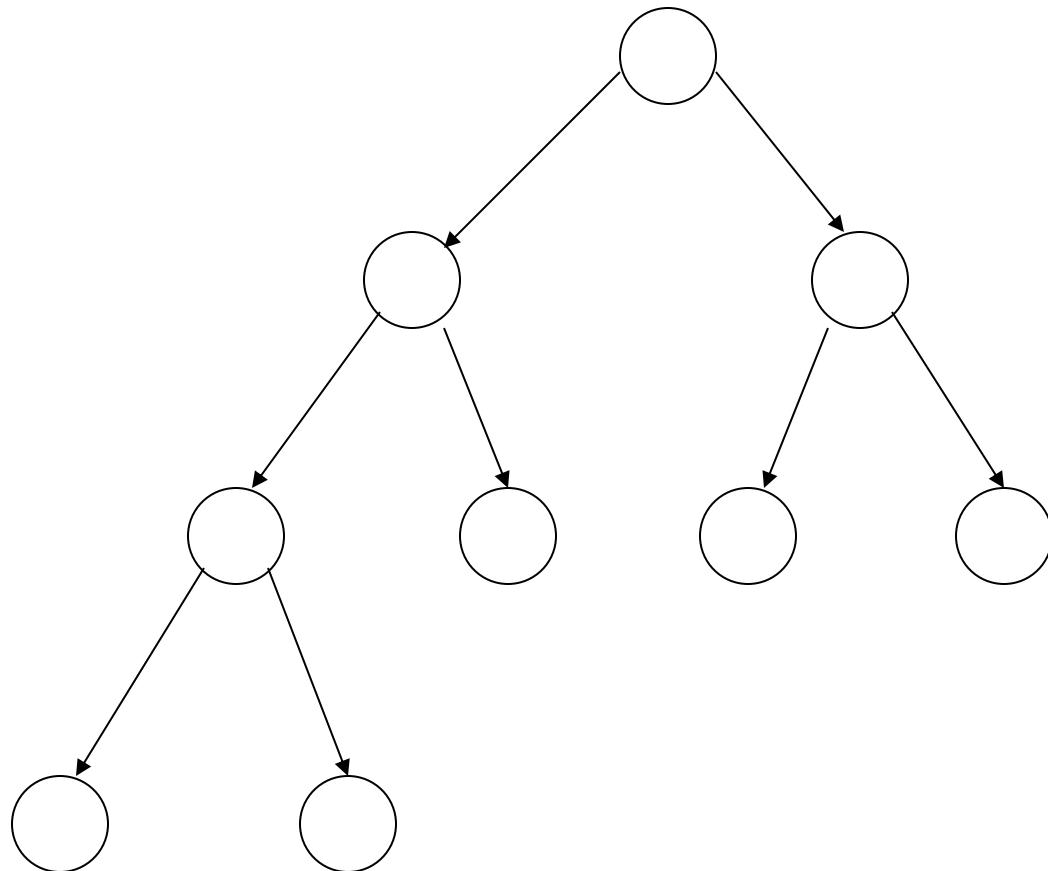
Trees



Trees have no cycles



Binary Trees



PROOF TECHNIQUES

- Proof by induction
- Proof by contradiction

Induction

We have statements P_1, P_2, P_3, \dots

If we know

- for some b that P_1, P_2, \dots, P_b are true
- for any $k \geq b$ that

P_1, P_2, \dots, P_k imply P_{k+1}

Then

Every P_i is true

Proof by Induction

- Inductive basis

Find P_1, P_2, \dots, P_b which are true

- Inductive hypothesis

Let's assume P_1, P_2, \dots, P_k are true,
for any $k \geq b$

- Inductive step

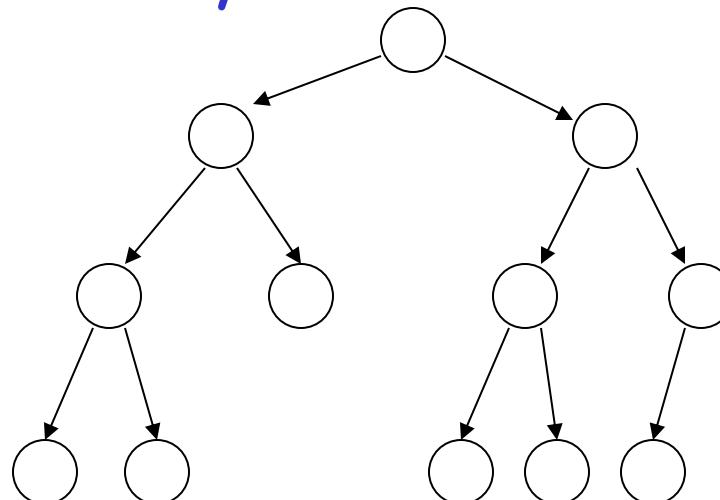
Show that P_{k+1} is true

Example

Theorem: A binary tree of height n
has at most 2^n leaves.

Proof by induction:

let $L(i)$ be the maximum number of
leaves of any subtree at height i



We want to show: $L(i) \leq 2^i$

- Inductive basis

$L(0) = 1$ (the root node)



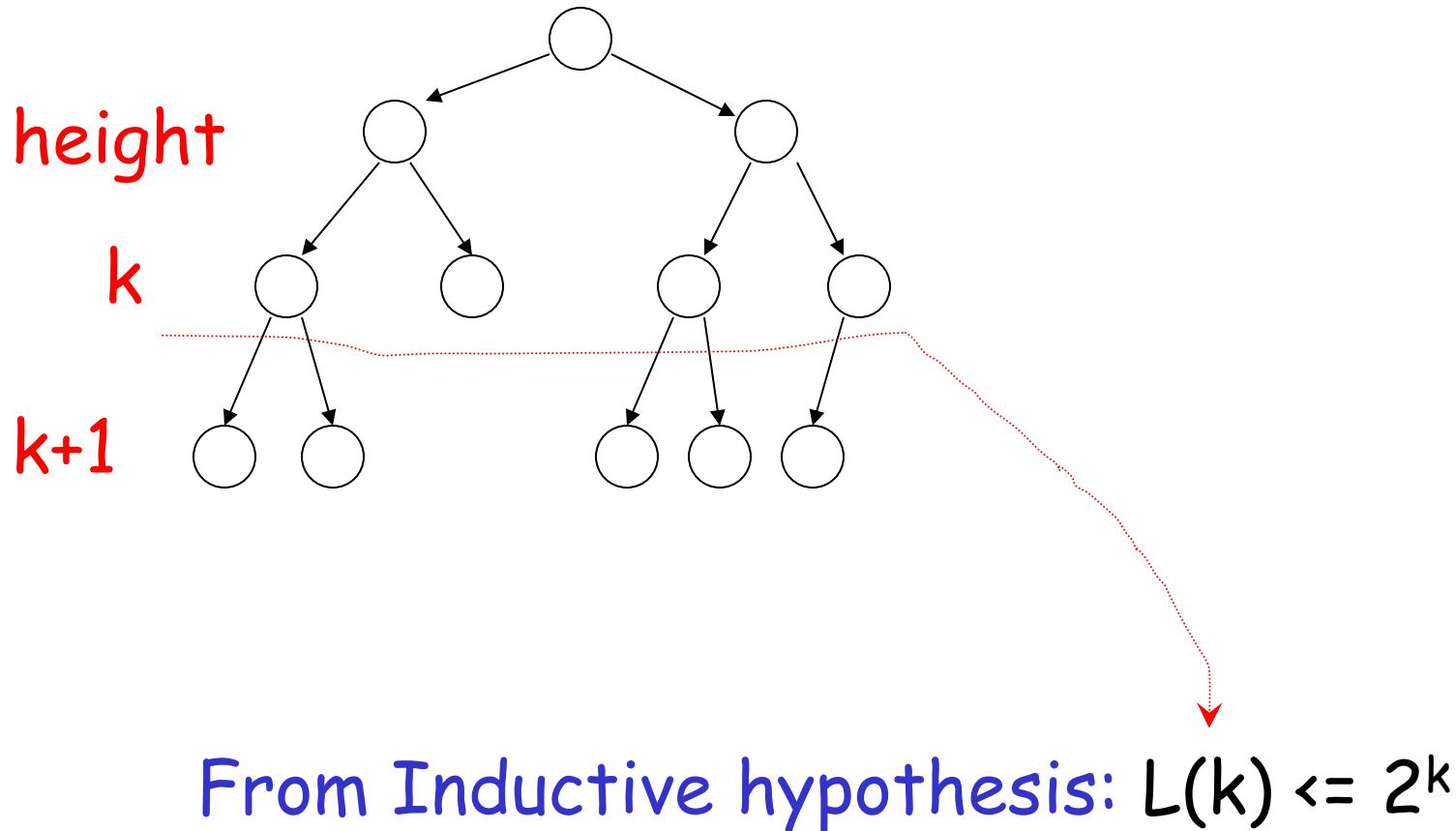
- Inductive hypothesis

Let's assume $L(i) \leq 2^i$ for all $i = 0, 1, \dots, k$

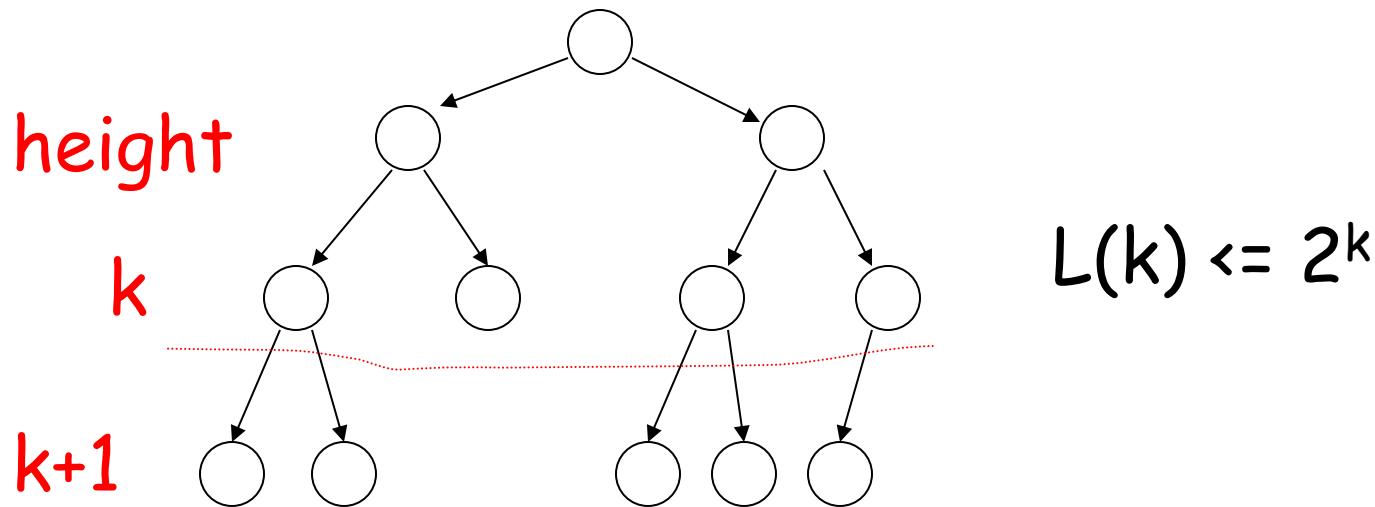
- Induction step

we need to show that $L(k + 1) \leq 2^{k+1}$

Induction Step



Induction Step



$$L(k+1) \leq 2 * L(k) \leq 2 * 2^k = 2^{k+1}$$

(we add at most two nodes for every leaf of level k)

Remark

Recursion is another thing

Example of recursive function:

$$f(n) = f(n-1) + f(n-2)$$

$$f(0) = 1, \quad f(1) = 1$$

Proof by Contradiction

We want to prove that a statement P is true

- we assume that P is false
- then we arrive at an incorrect conclusion
- therefore, statement P must be true

Example

Theorem: $\sqrt{2}$ is not rational

Proof:

Assume by contradiction that it is rational

$$\sqrt{2} = n/m$$

n and m have no common factors

We will show that this is impossible

$$\sqrt{2} = n/m \longrightarrow 2m^2 = n^2$$

Therefore, n^2 is even \longrightarrow n is even
 $n = 2k$

$$2m^2 = 4k^2 \longrightarrow m^2 = 2k^2 \longrightarrow m \text{ is even}$$
$$m = 2p$$

Thus, m and n have common factor 2

Contradiction!